

Oct. 6, 1970

M. A. GOETZ
AUTOMATIC SYSTEM FOR CONSTRUCTING
AND RECORDING DISPLAY CHARTS

3,533,086

Filed Dec. 24, 1968

38 Sheets-Sheet 1

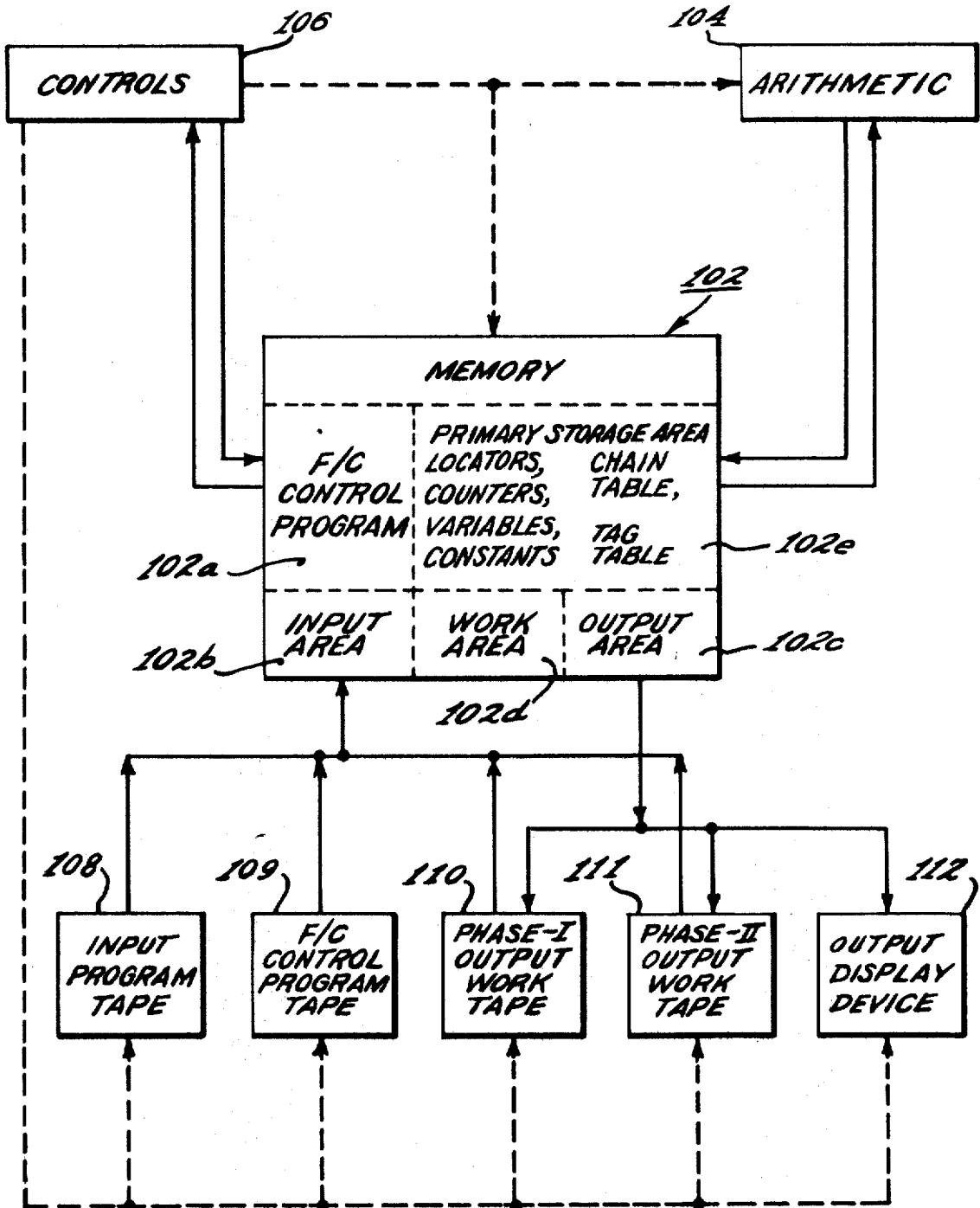


Fig. 1.

INVENTOR.
Martin A. Goetz
BY
Millman and Jacobs
ATTORNEYS.

Oct. 6, 1970

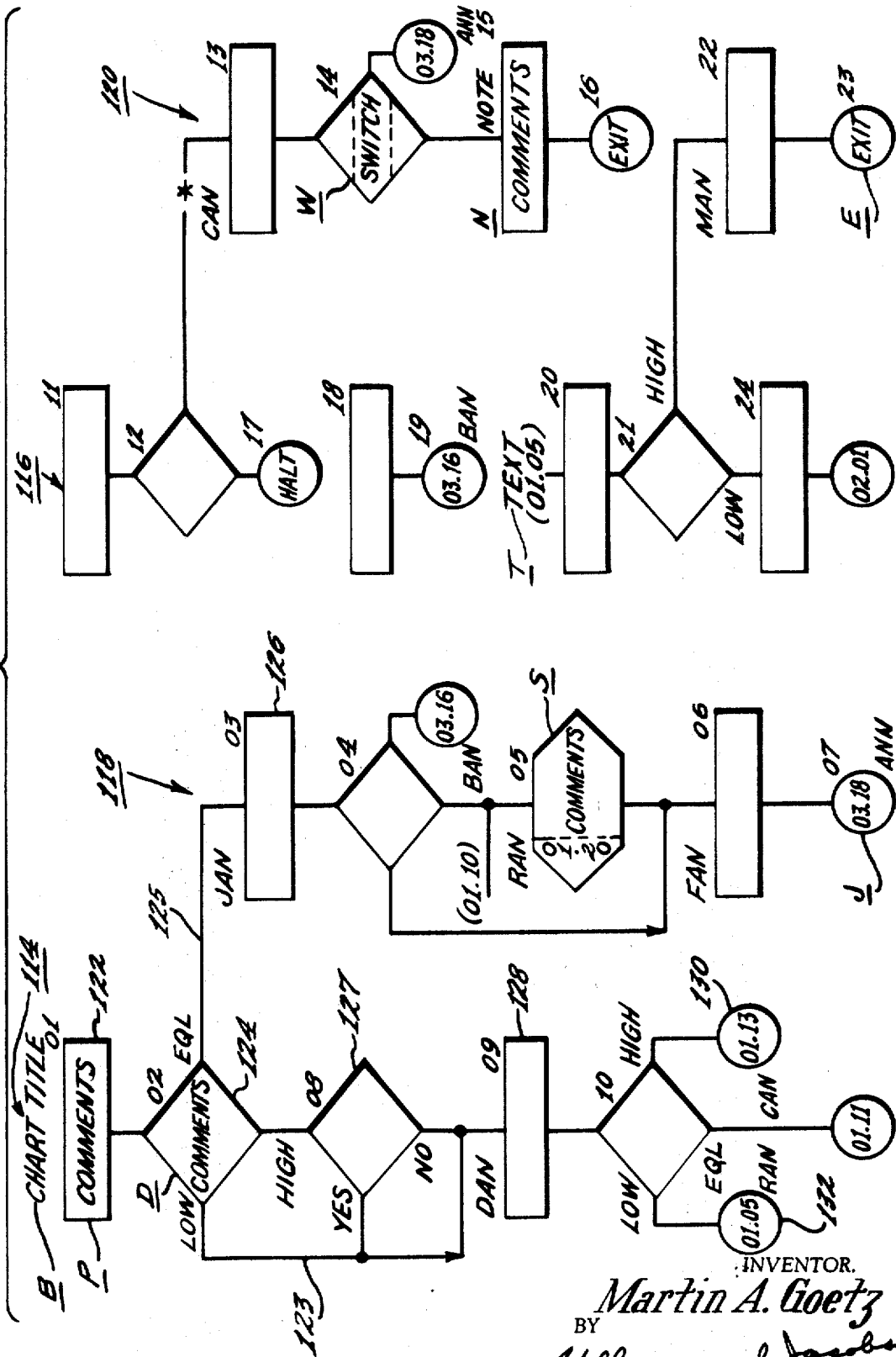
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Fig. 2.



INVENTOR.
Martin A. Goetz
BY *Millman and Jacobs*
ATTORNEYS.

Oct. 6, 1970

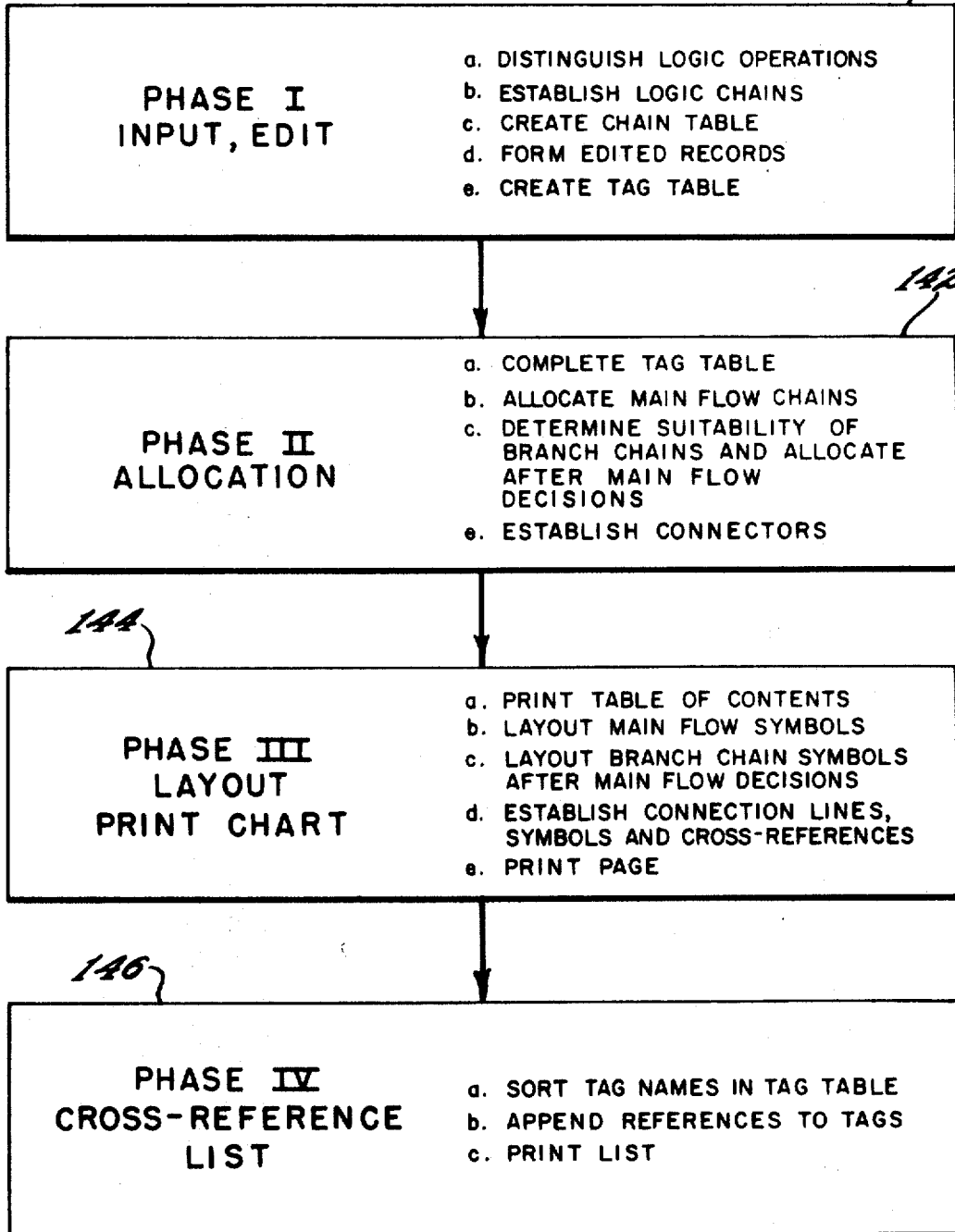
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Fig. 3.



INVENTOR.

Martin A. Goetz

BY

William and Jacob

ATTORNEYS.

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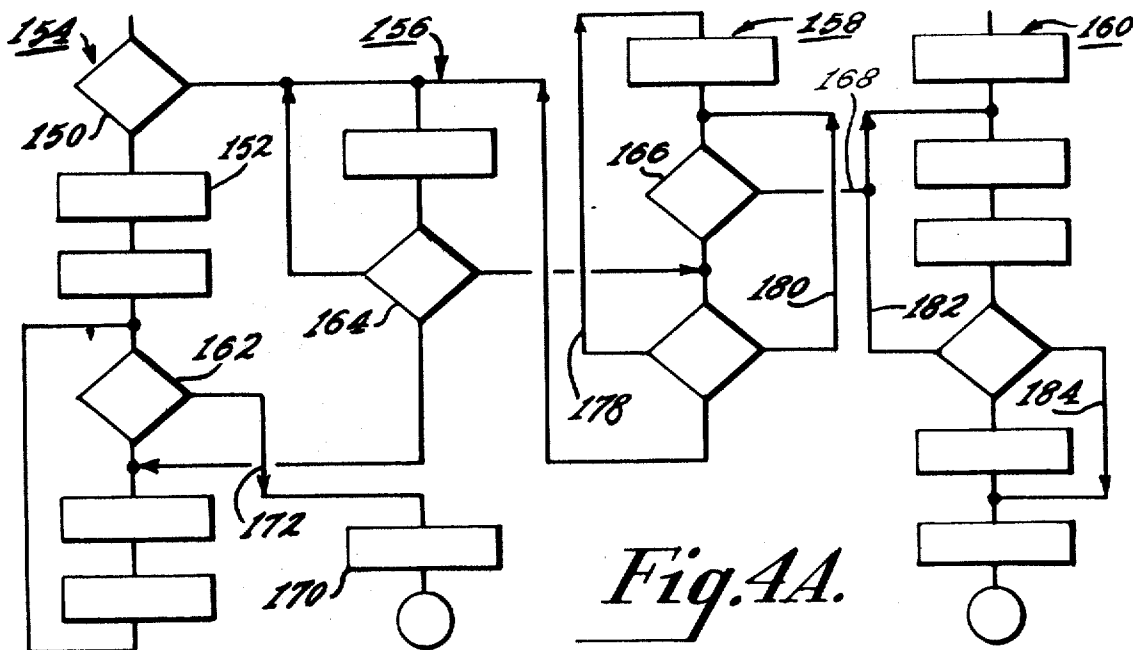


Fig. 4A.

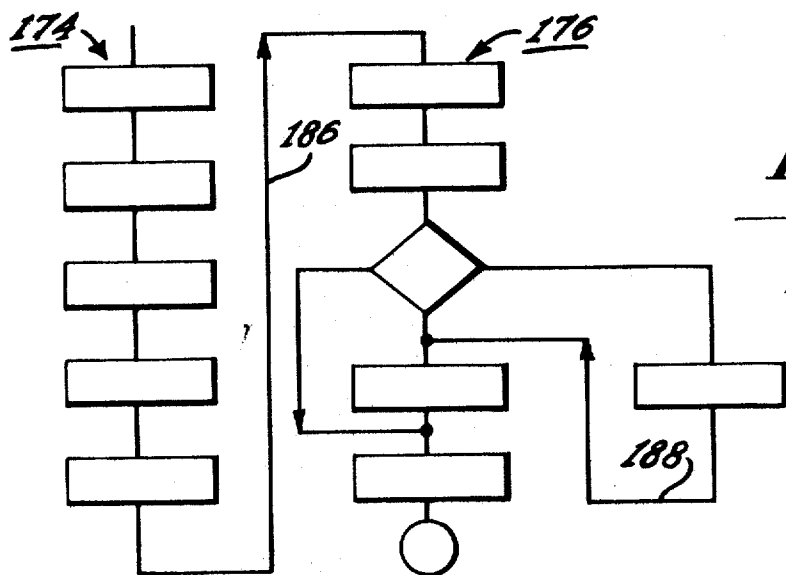


Fig. 4B.

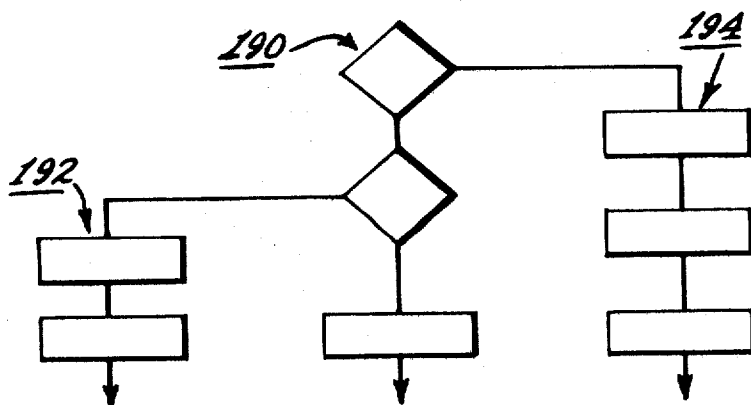


Fig. 5.

INVENTOR.
Martin A. Goetz
BY
Williamson and Jacobs
ATTORNEYS.

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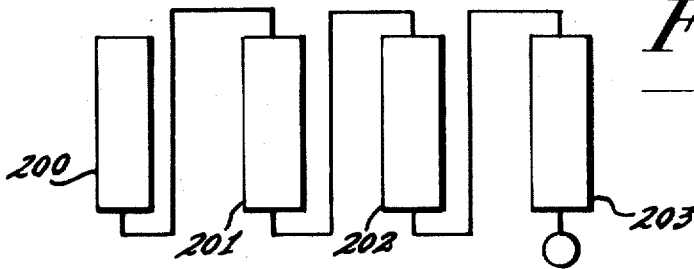


Fig. 6A.

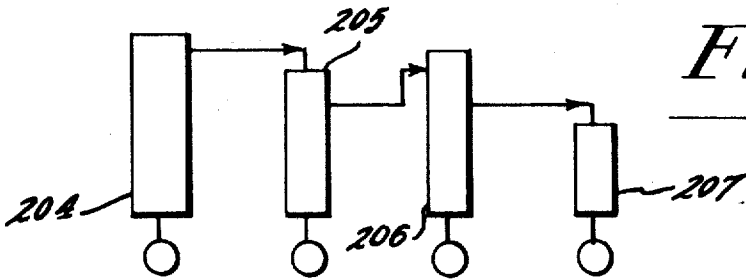


Fig. 6B.

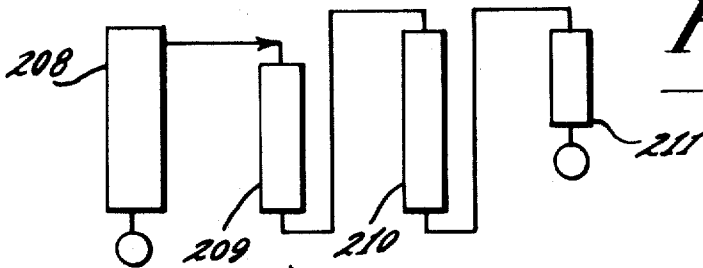


Fig. 6C.

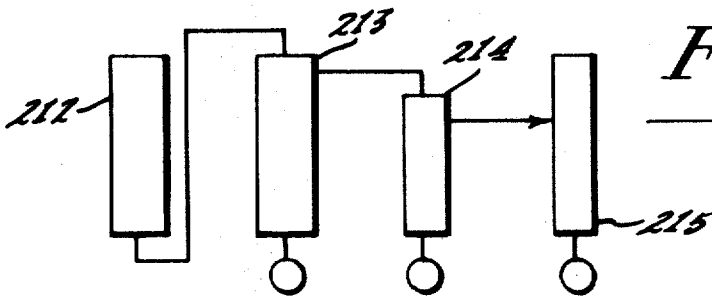


Fig. 6D.

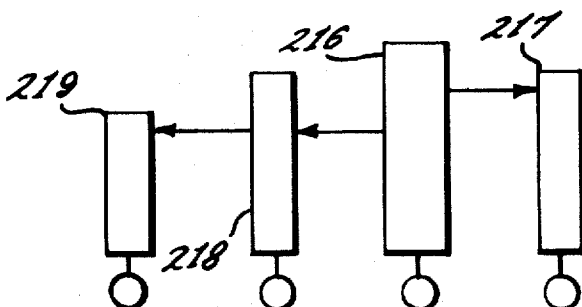


Fig. 6E.

INVENTOR.

Martin A. Goetz

BY

Willson and Jacobs

ATTORNEYS.

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Fig. 7A.

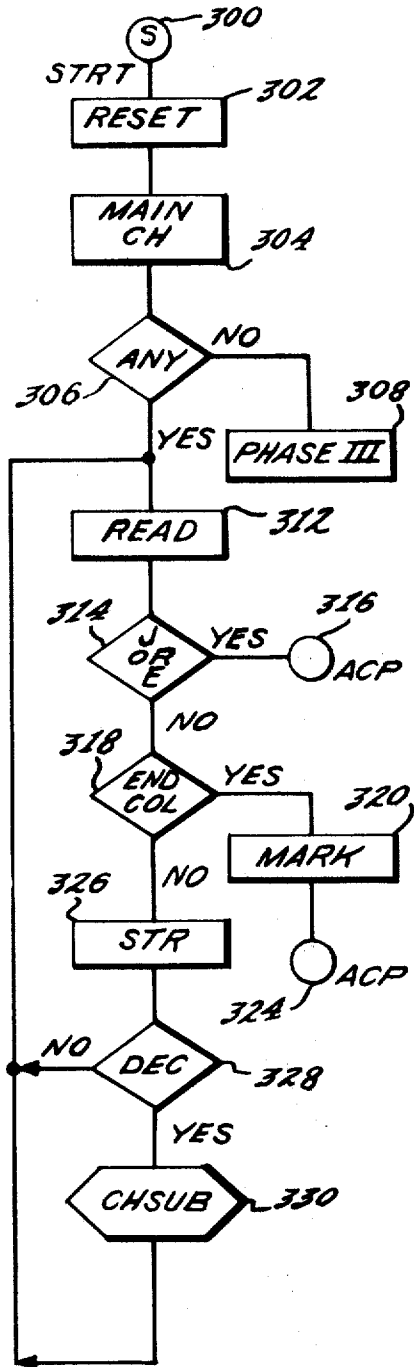
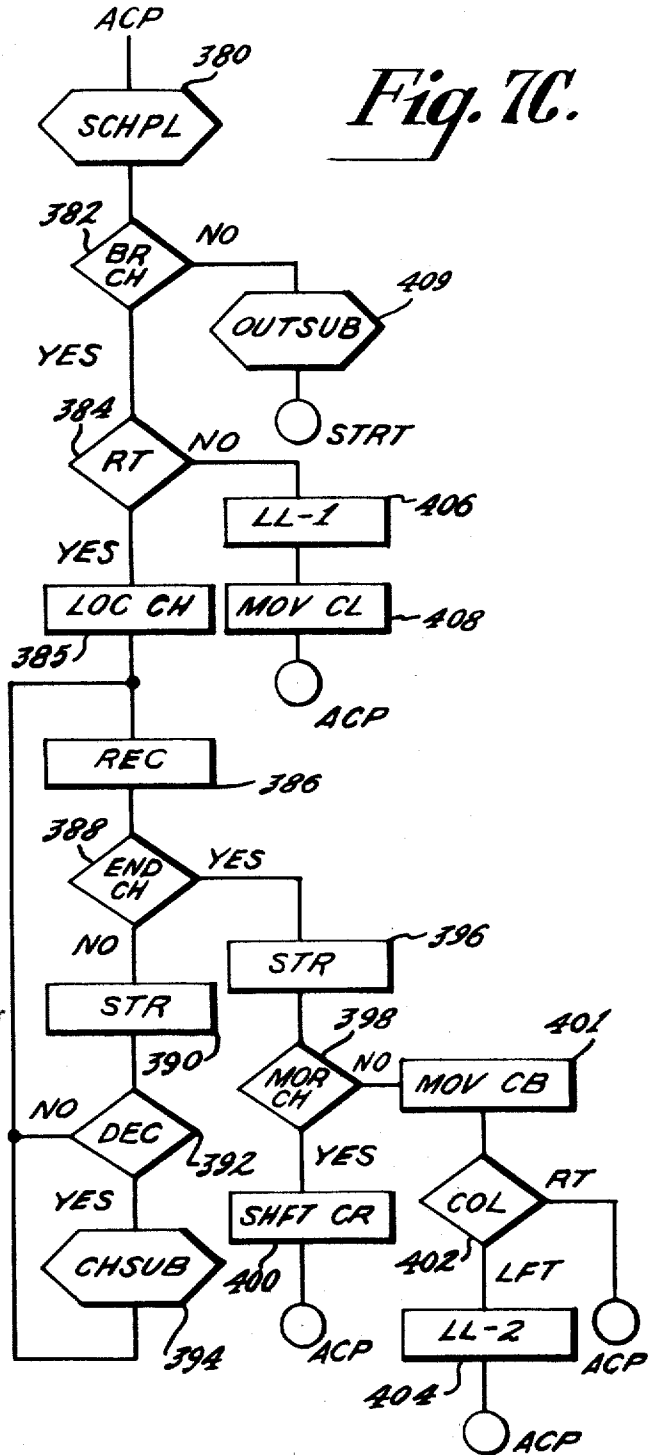


Fig. 7C.



INVENTOR.

Martin A. Goetz

BY

Millman and Jacobs

ATTORNEYS.

Oct. 6, 1970

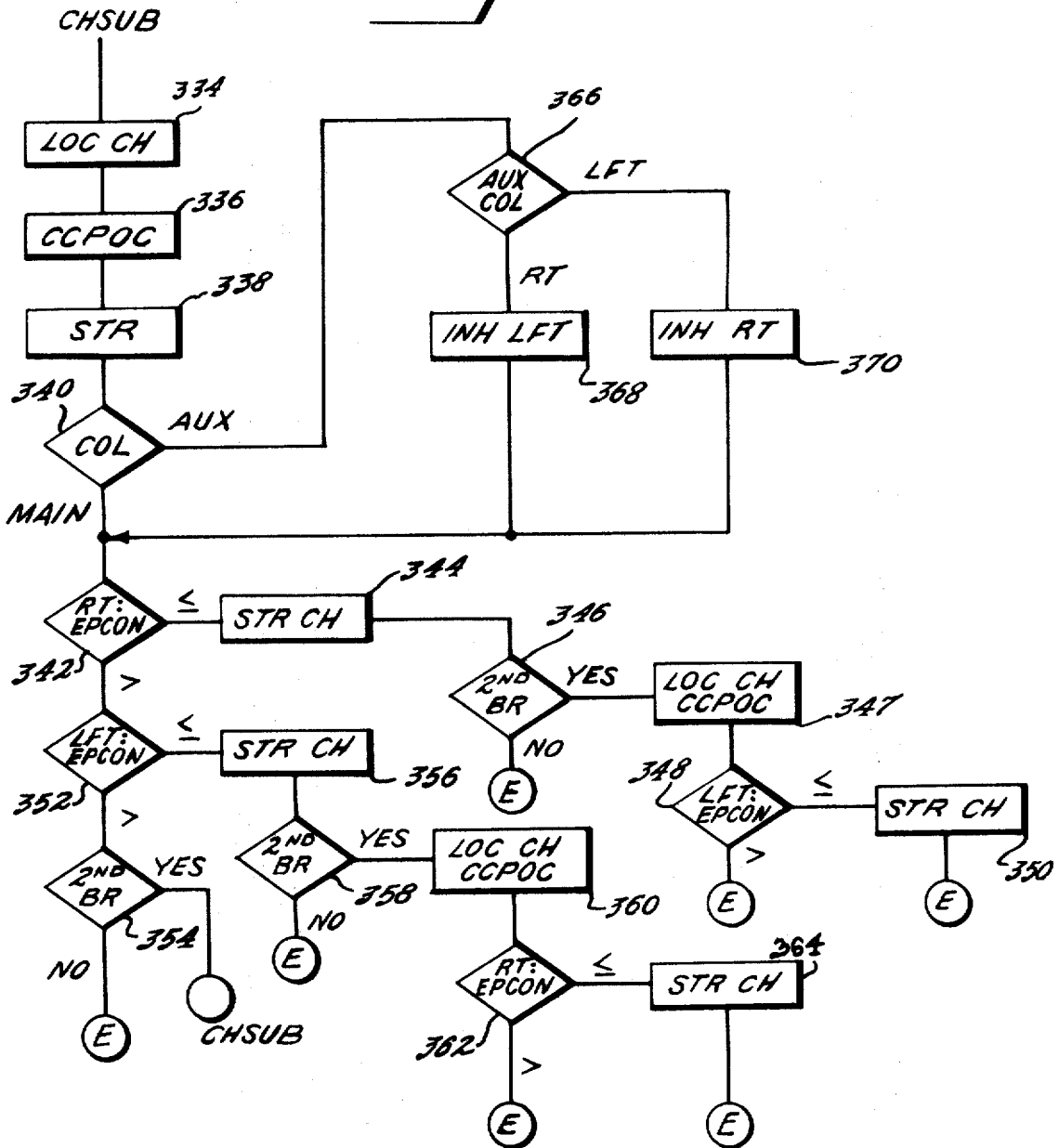
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Fig. 7B.



INVENTOR.
Martin A. Goetz
 BY
William and Jacobs
 ATTORNEYS.

Oct. 6, 1970

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Fig. 7D.

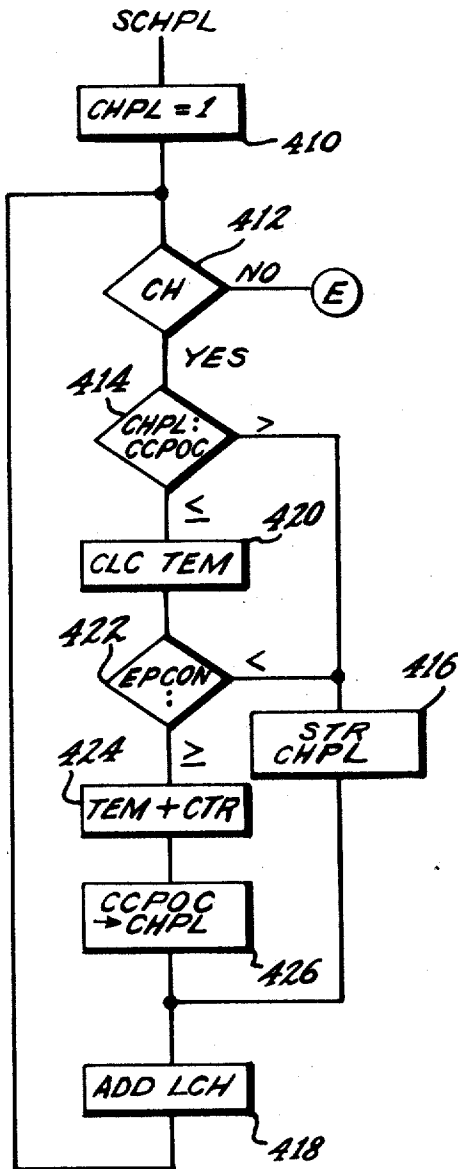
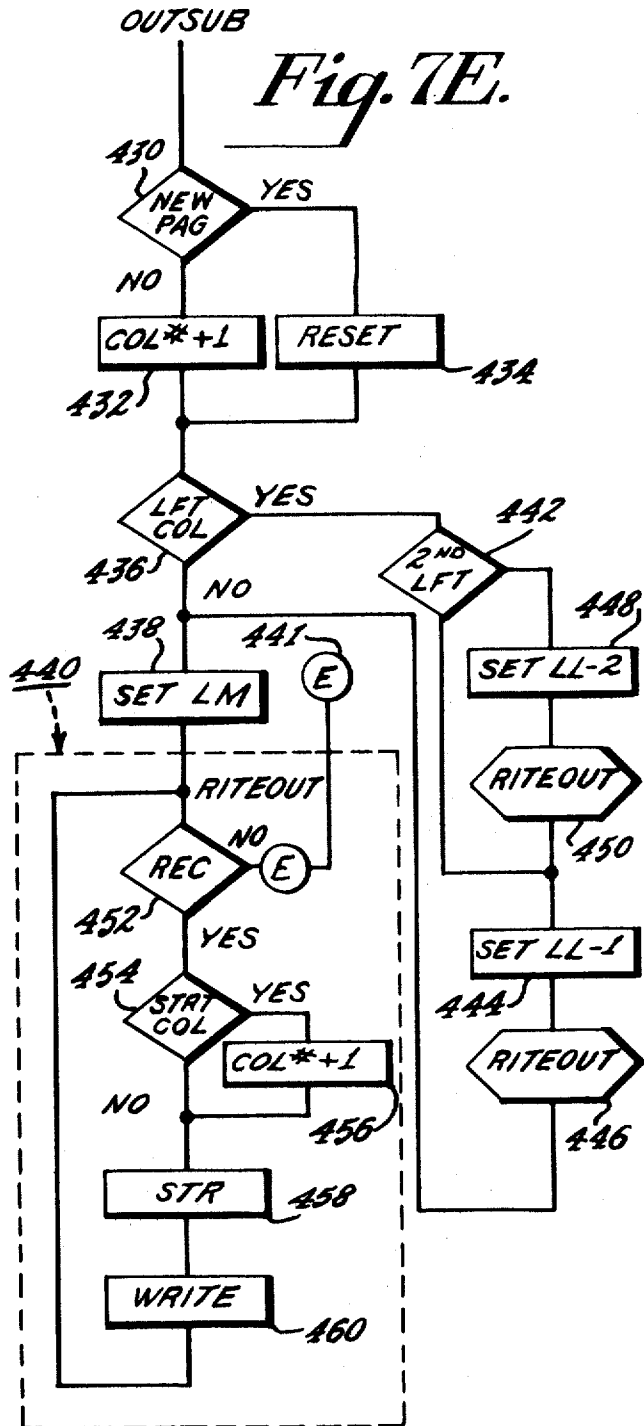


Fig. 7E.



INVENTOR.
Martin A. Goetz
 BY
William and Jacobs
 ATTORNEYS.

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Fig. 8/1

PROGRAM: AUTOFLOW

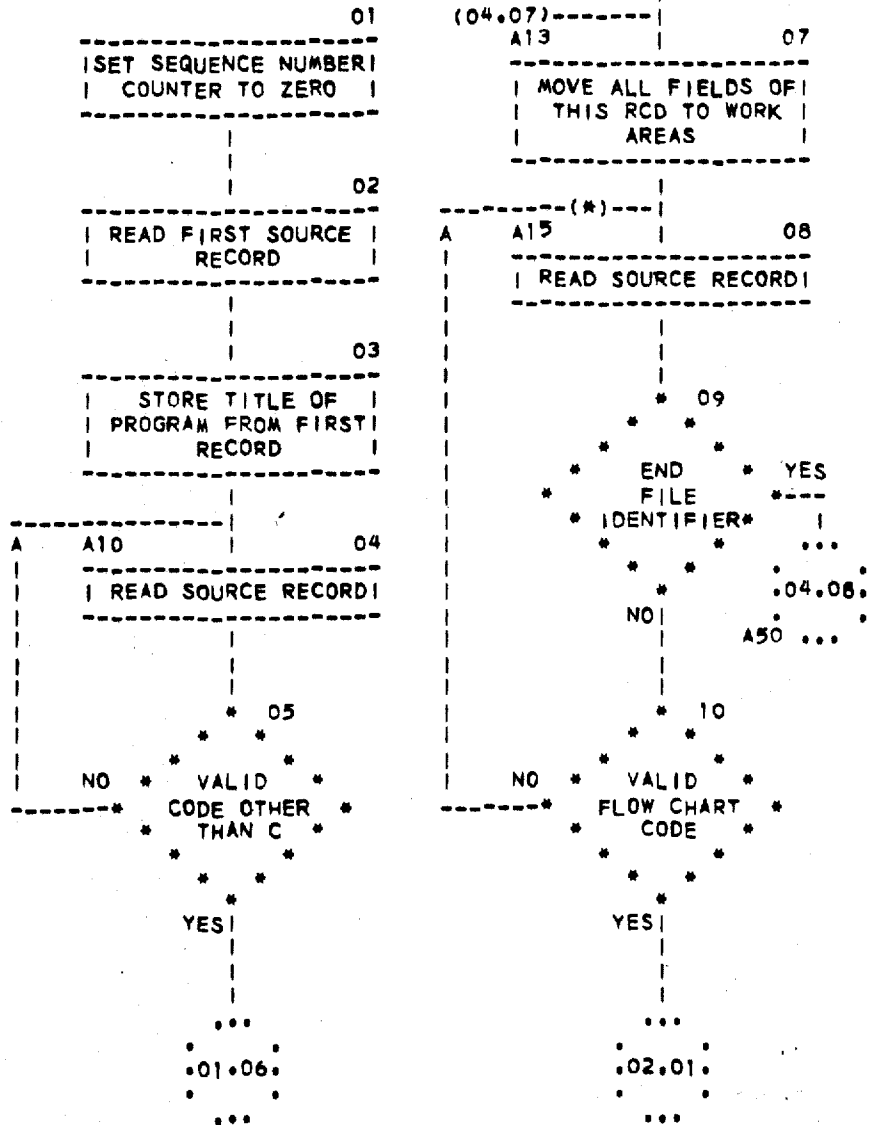
PAGE 1

CHART: PASS ONE - INPUT EDIT

FUNCTIONS OF PASS ONE ARE TO
EDIT SOURCE INFORMATION AND
CONSTRUCT TAG AND CHAIN
TABLES

NOTE 06

* INPUT AREA NOW*
* HAS RCD TO BE *
* PROCESSED *



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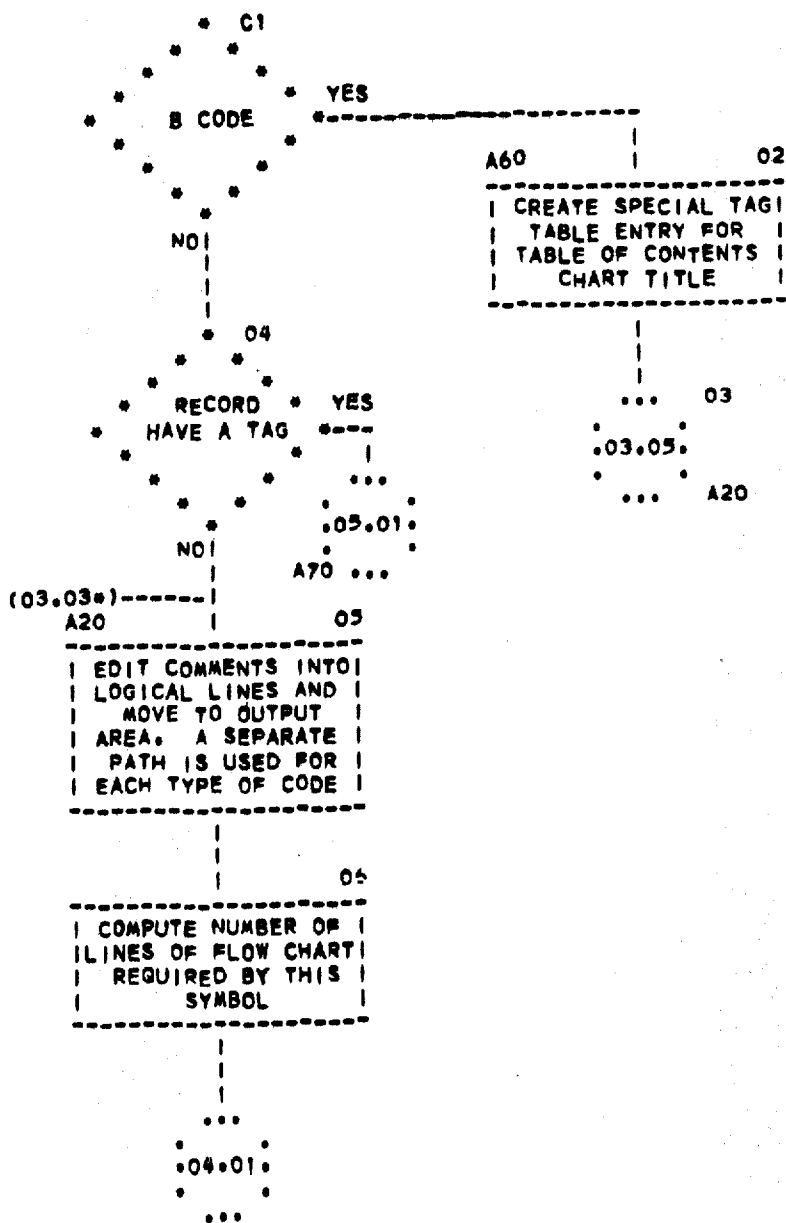
38 Sheets-Sheet 11

Fig. 8/3

PROGRAM: AUTOFLOW

PAGE 3

CHART: PASS ONE - INPUT EDIT



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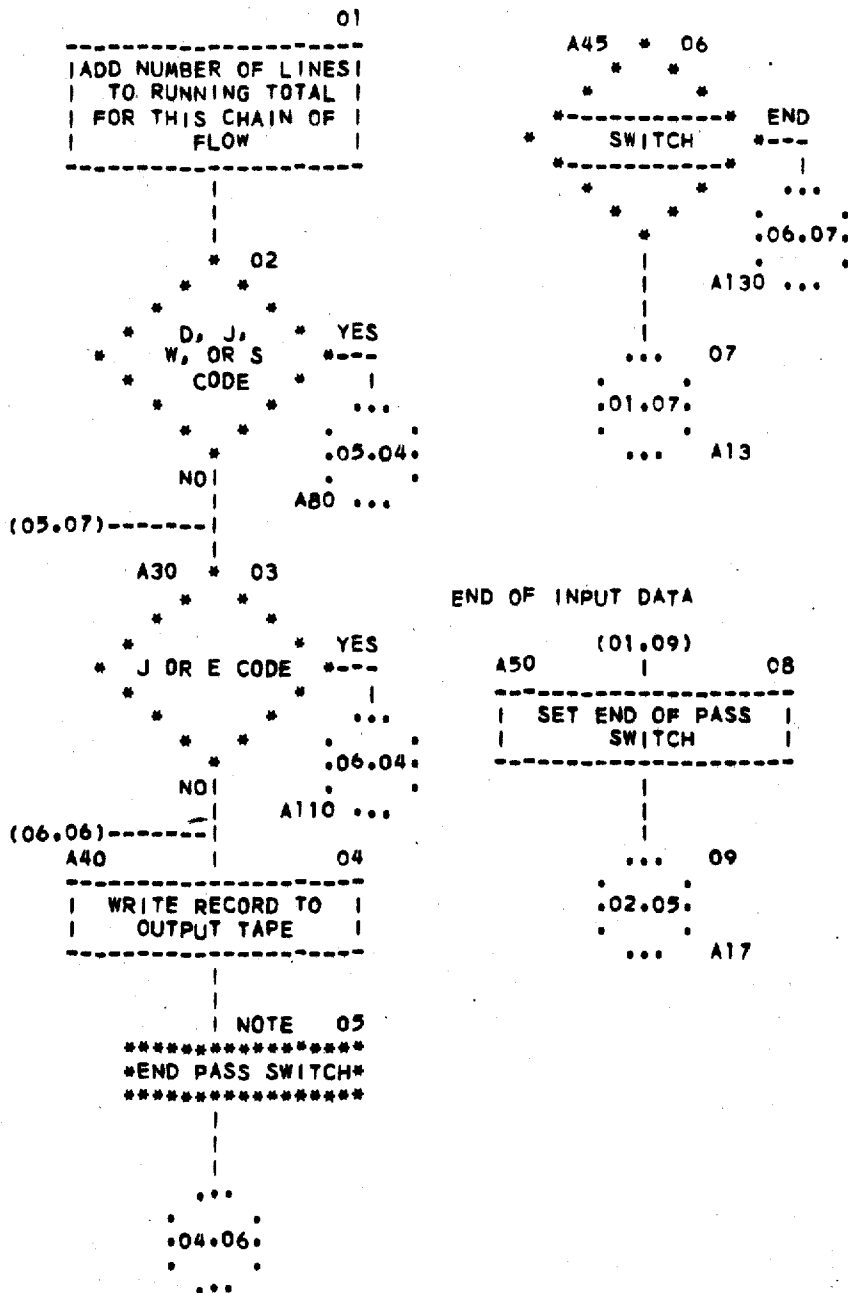
58 Sheets-Sheet 12

Fig. 8/4

PROGRAM: AUTOFLOW

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CHART: PASS ONE - INPUT EDIT



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Fig. 8/5

PROGRAM: AUTOFLOW

PAGE 5

CHART: PASS ONE - INPUT EDIT

INPUT RECORD HAS A TAG

D, J, W, OR S CODE

(03.04)
A70 | 01

| CREATE TAG TABLE |
| ENTRY CONTAINING |
| TAG, SEQUENCE |
| NUMBER, AND |
| ASSEMBLY LINE |
| NUMBER OF THIS |
RECORD

02

| MOVE TAG TO TAG |
| FIELD OF OUTPUT |
RECORD

... 03

.03.05.

... A20

(04.02)

A80 * 04

| * DOES * |
| * COMMENTS * NO |
| * CONTAIN * --- |
| * ASTERISK * |
* FIELD * ...

* * .06.01.

YES |

A90 ...

05

| EXTRACT |
| DESTINATIONS AND |
| LABELS FROM |
ASTERISK FIELD

(06.03)

A100

06

| MOVE LABELS AND |
| DESTINATIONS TO |
OUTPUT

... 07

.04.03.

... A30

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38 Sheets-Sheet 14

Fig. 8/6

PROGRAM: AUTOFLOW

PAGE 6

CHART: PASS ONE - INPUT EDIT

A90 (05.04) 01

| PICK UP
| DESTINATIONS FROM
SOURCE OPERANDS

02

| PICK UP LABELS FROM
| STORED TABLE, BASED
ON LABEL CODE

... 03

.05.06.

... A100

A130 (04.06) 07

| REWIND TAPES, CALL
IN PASS 2 SEGMENT

... 08

.07.01.

... 800

A110 (04.03) 04

| CREATE NEW ENTRY IN
| CHAIN TABLE FOR
| CHAIN ENDING WITH
CURRENT RECORD

05

| PLACE IN CHAIN
| TABLE ENTRY THE
| SEQUENCE NUMBER OF
| THE CURRENT RECORD
| AND NUMBER OF LINES
REQUIRED ON CHART

... 06

.04.04.

... A40

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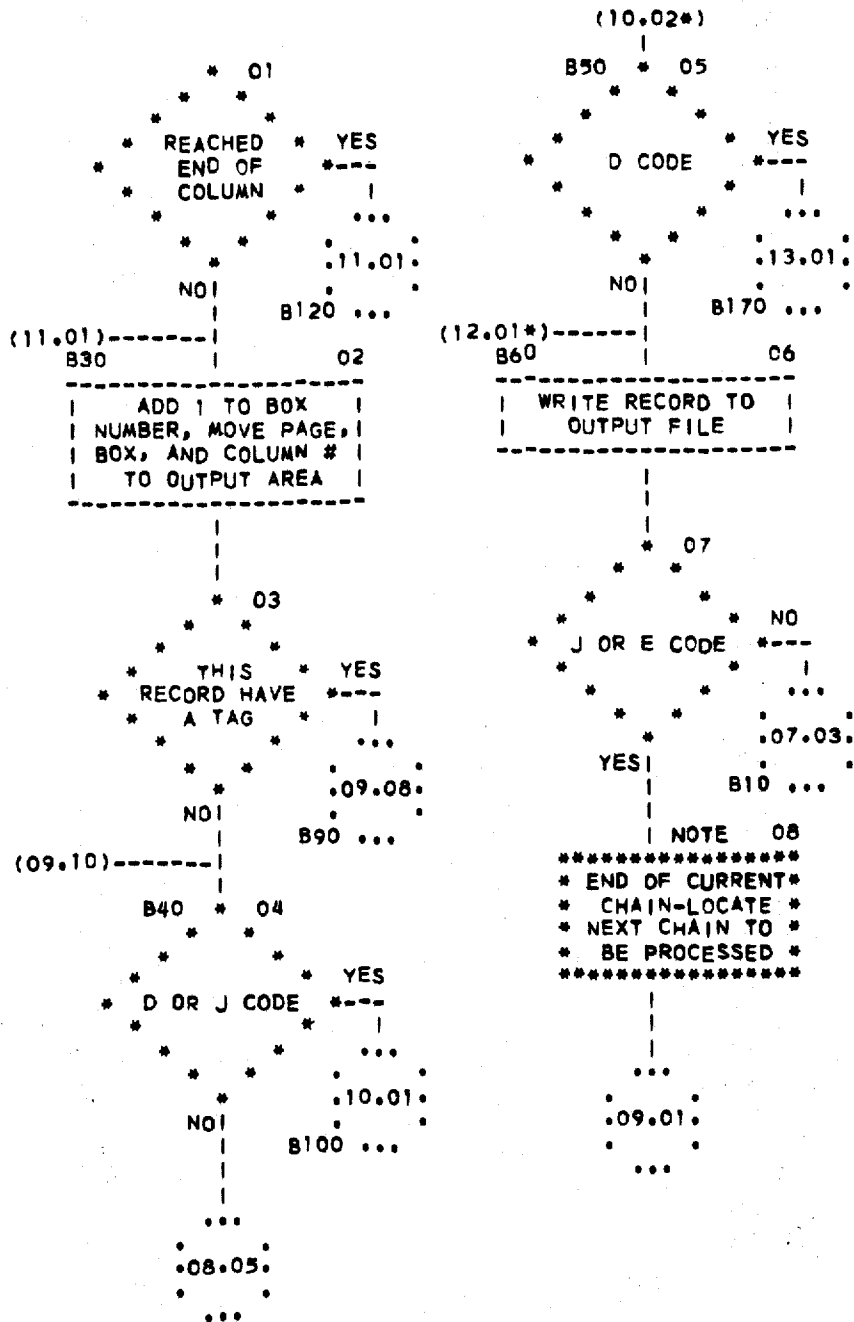
38 Sheets-Sheet 16

Fig. 8/8

PROGRAM: AUTOFLOW

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CHART: PASS TWO - ALLOCATION



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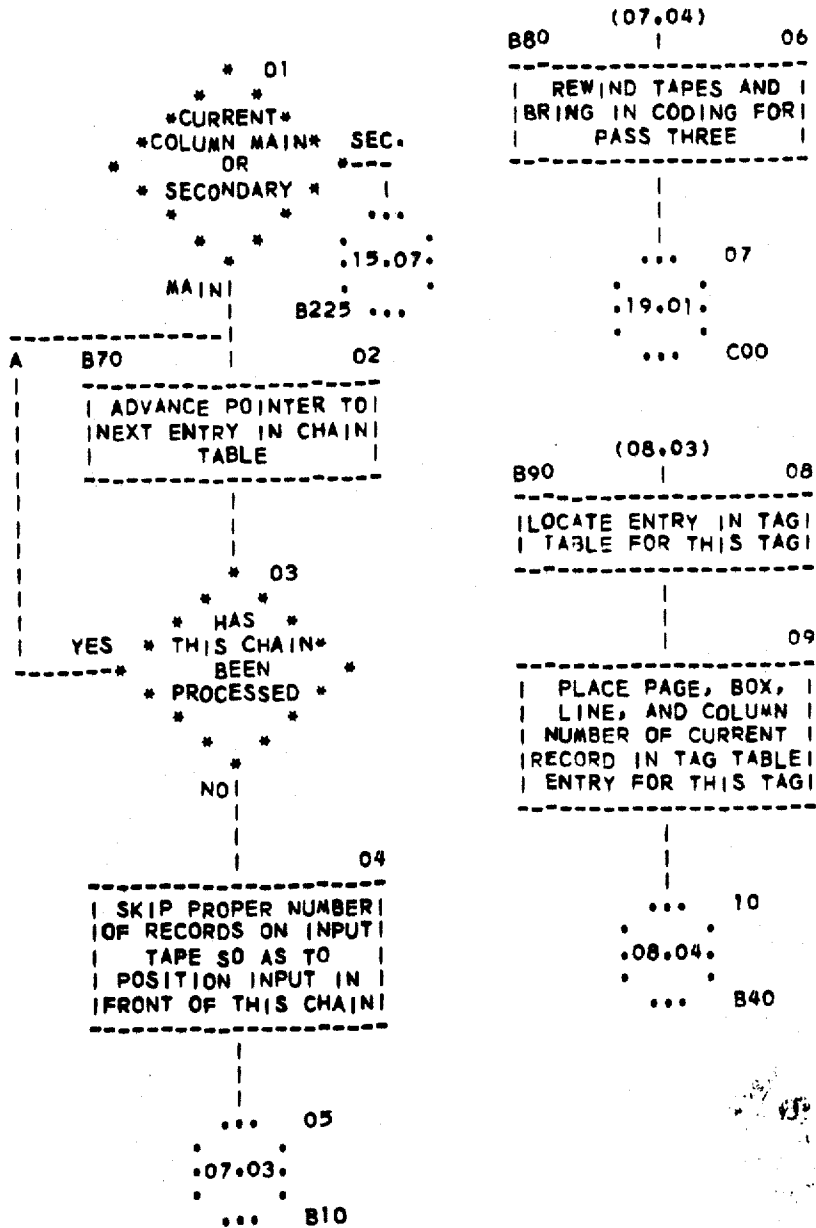
38 Sheets-Sheet 17

Fig. 8/9

PROGRAM: AUTOFLOW

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CHART: PASS TWO - ALLOCATION



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Fig. 8/10

PROGRAM: AUTOFLOW

PAGE 10

CHART: PASS TWO - ALLOCATION

CHECK FOR "FROM" CONNECTOR

B100 (08.04) 01

LOOK UP DESTINATION
TAGS IN TAG TABLE

* 02

* TAG * NO
* FOUND IN *
* TABLE * |

...
.08.05.
... B50 ...

YES |

* 03

* FIRST * NO
* REFERENCE *
* TO THIS TAG * |

YES |

06

PUT PAGE AND BOX
NUMBER OF CURRENT
RECORD IN CROSS
REFERENCE FIELD OF
TAG TABLE ENTRY

B110 04

SET SIGNAL IN TAG
TABLE ENTRY TO
INDICATE THAT MORE
THAN ONE REFERENCE
EXISTS

... 05

.08.05.

... B50

... 07

.08.05.

... B50

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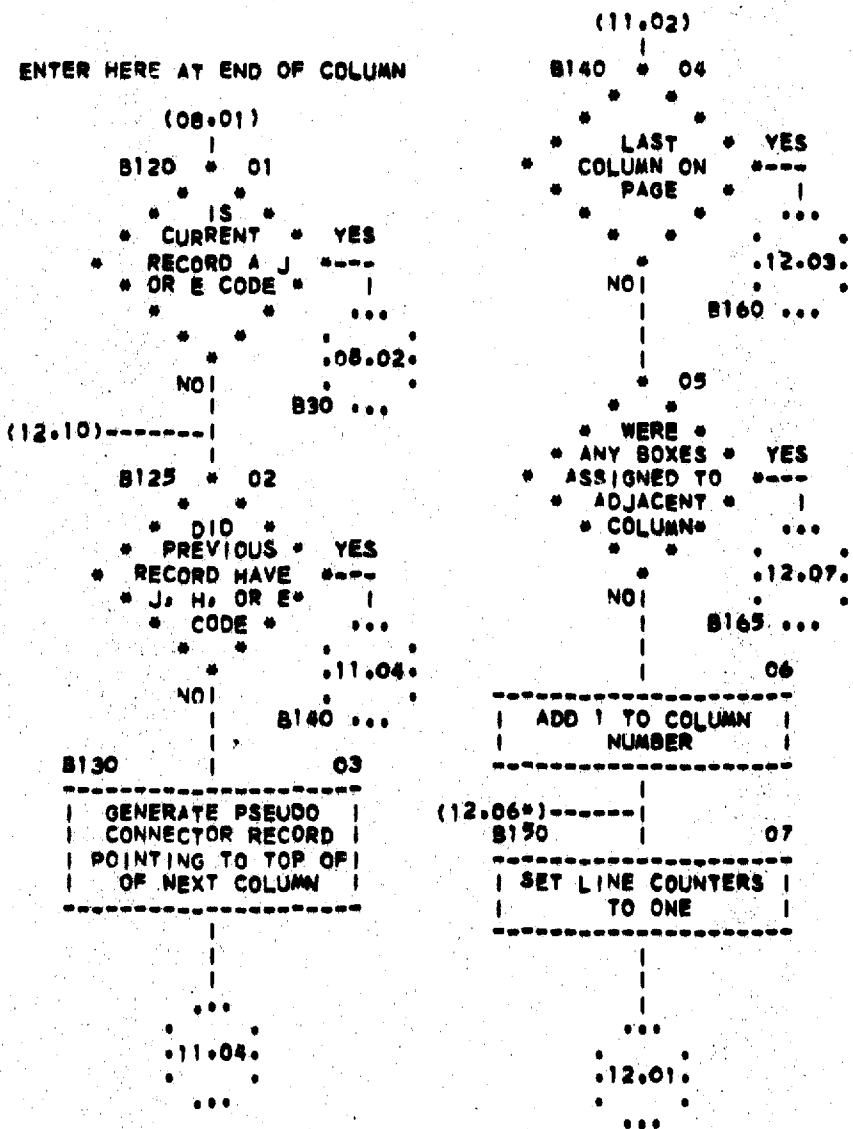
38 Sheets-Sheet 19

Fig. 8/II

PROGRAM: AUTOFLOW

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CHART: PASS TWO - ALLOCATION



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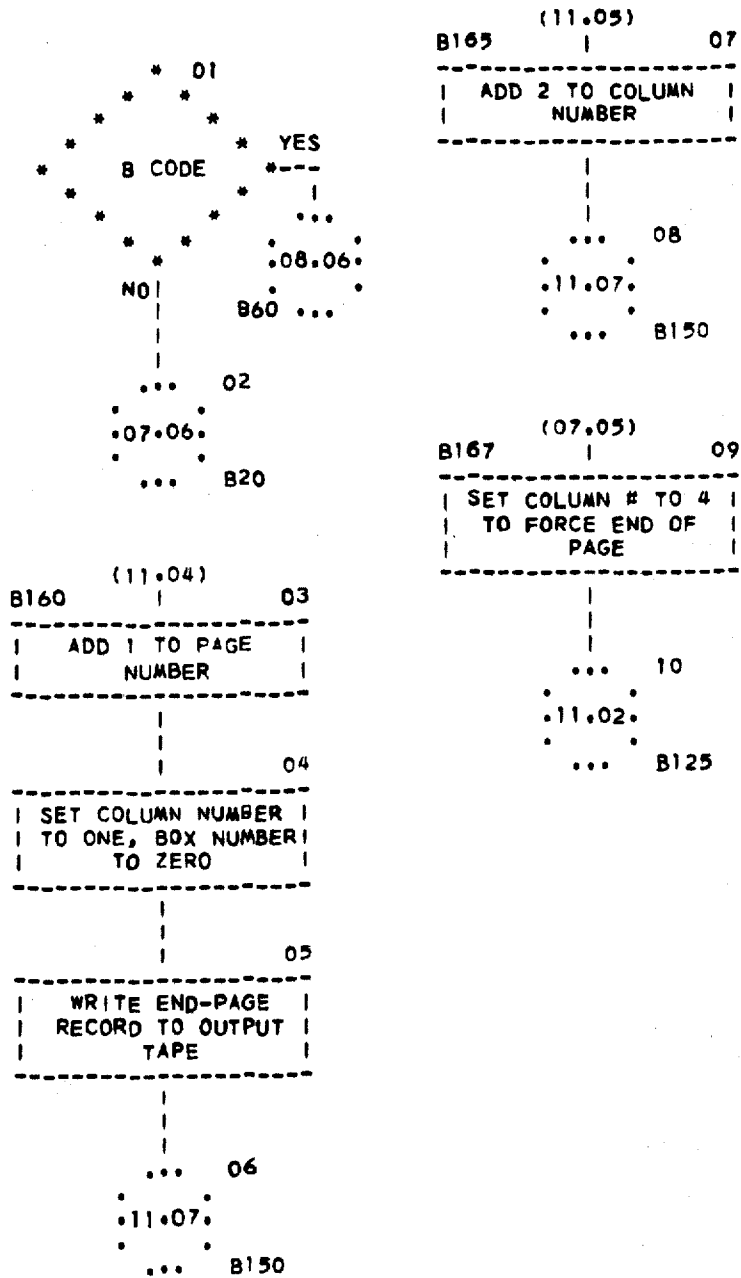
38 Sheets-Sheet 20

Fig. 8/12

PROGRAM: AUTOFLOW

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CHART: PASS TWO - ALLOCATION



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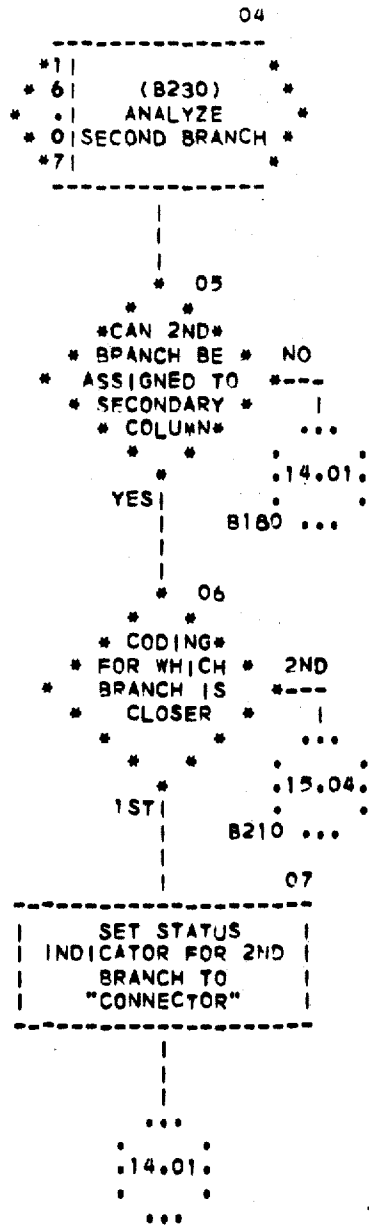
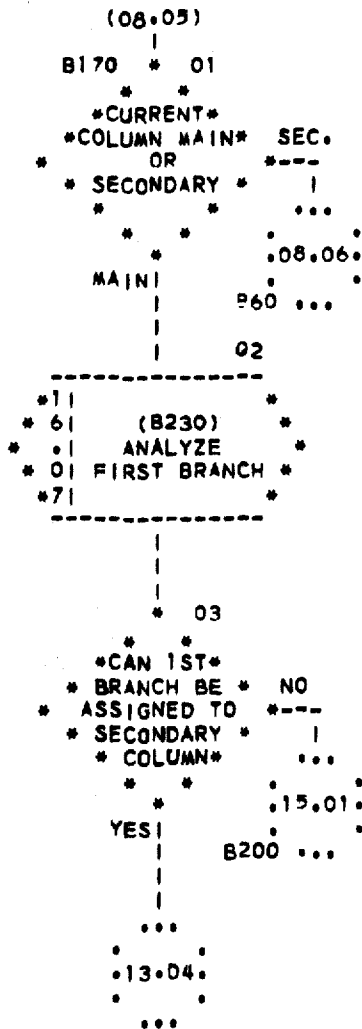
Fig. 8/13

PROGRAM: AUTOFLOW

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CHART: PASS TWO - ALLOCATION

ANALYZE DECISION RECORD TO
DETERMINE HOW TO SHOW ITS
BRANCHES ON THE CHART. A
"STATUS INDICATOR" FOR EACH
BRANCH IS SET TO EITHER
"SECONDARY COLUMN,"
"CONNECTOR," OR "UNKNOWN."
IF "UNKNOWN," PASS 3 WILL
DECIDE ON A LINE OR
CONNECTOR



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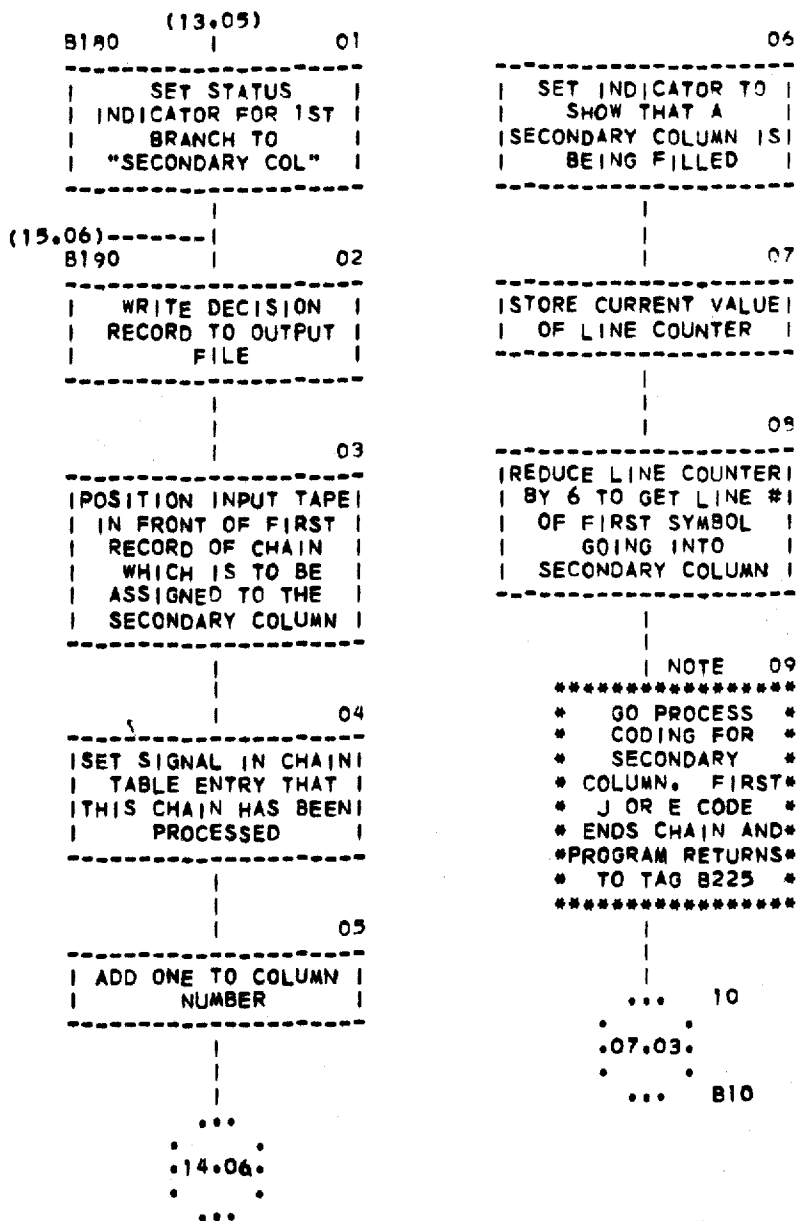
38 Sheets-Sheet 22

Fig. 8/14

PROGRAM: AUTOFLOW

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CHART: PASS TWO - ALLOCATION



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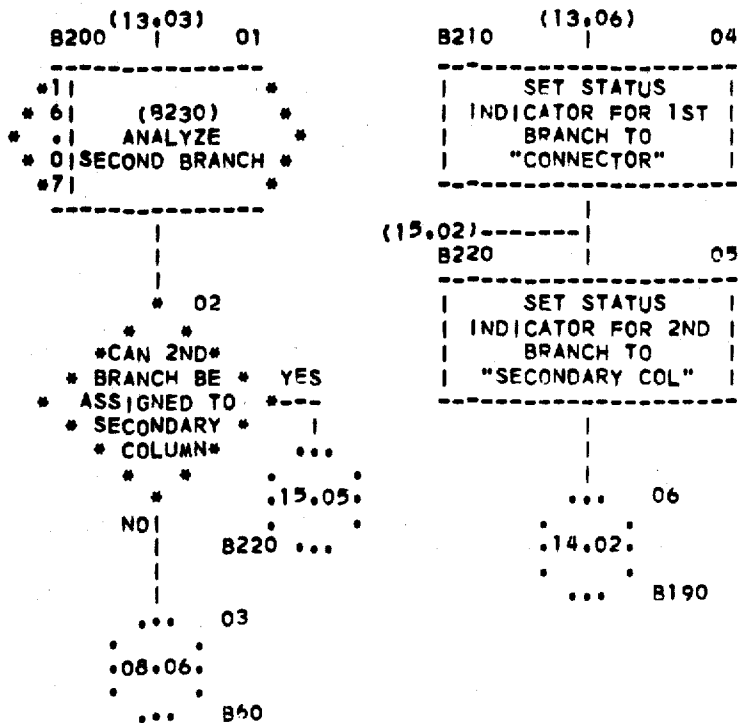
38 Sheets-Sheet 25

Fig. 8/15

PROGRAM: AUTOFLOW

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CHART: PASS TWO - ALLOCATION



RETURN AFTER PROCESSING
CHAIN FOR SECONDARY COLUMN

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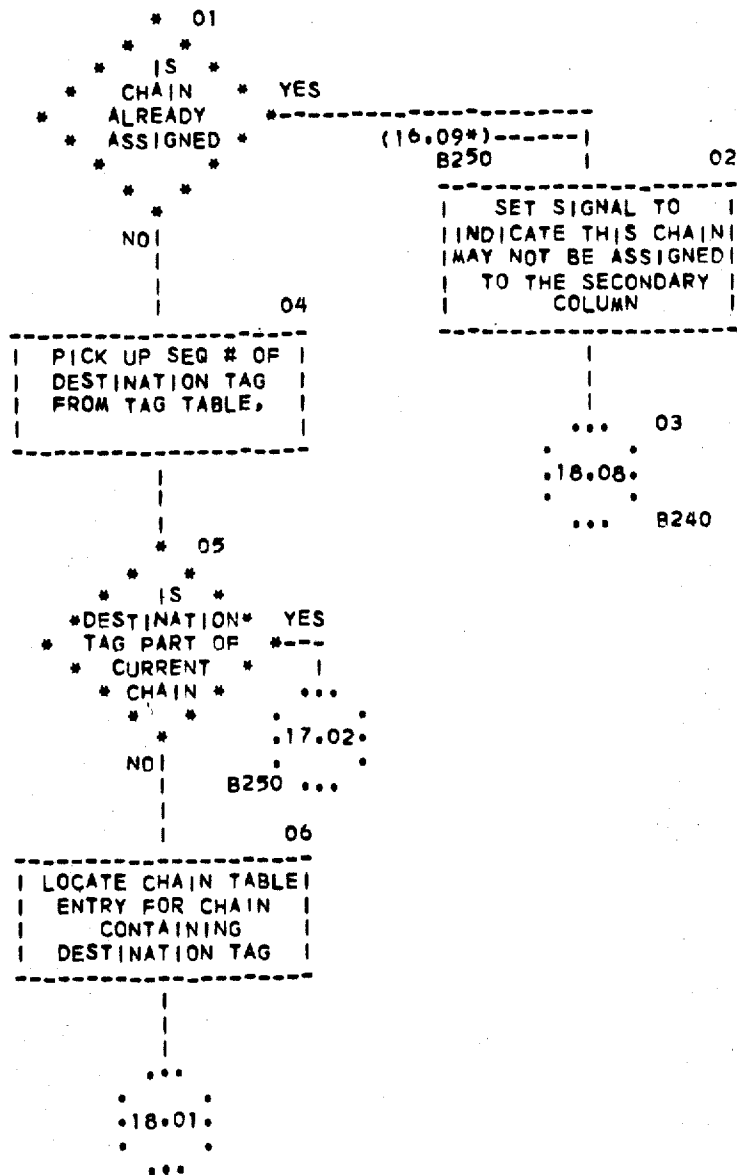
38 Sheets-Sheet 25

Fig. 8/17

PROGRAM: AUTOFLOW

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CHART: PASS TWO - ALLOCATION



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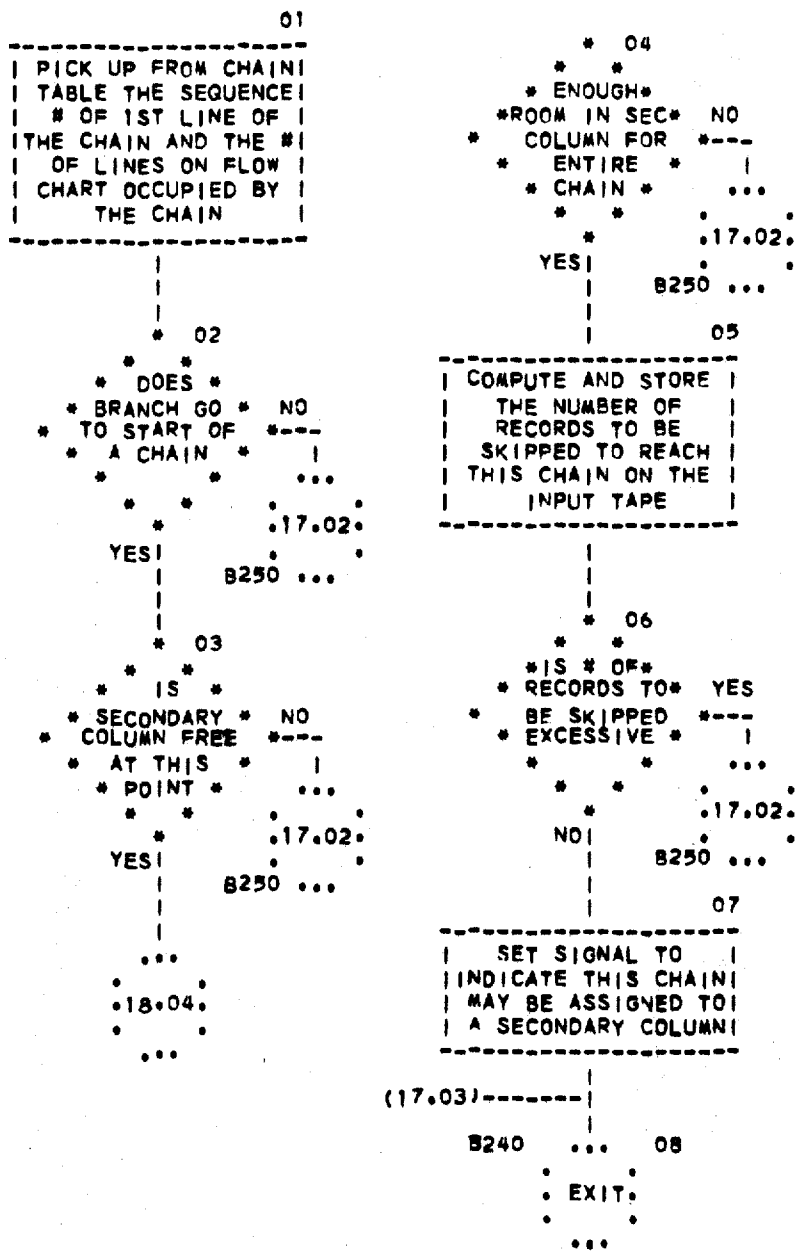
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Fig. 8/18

PROGRAM: AUTOFLOW

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CHART: PASS TWO - ALLOCATION



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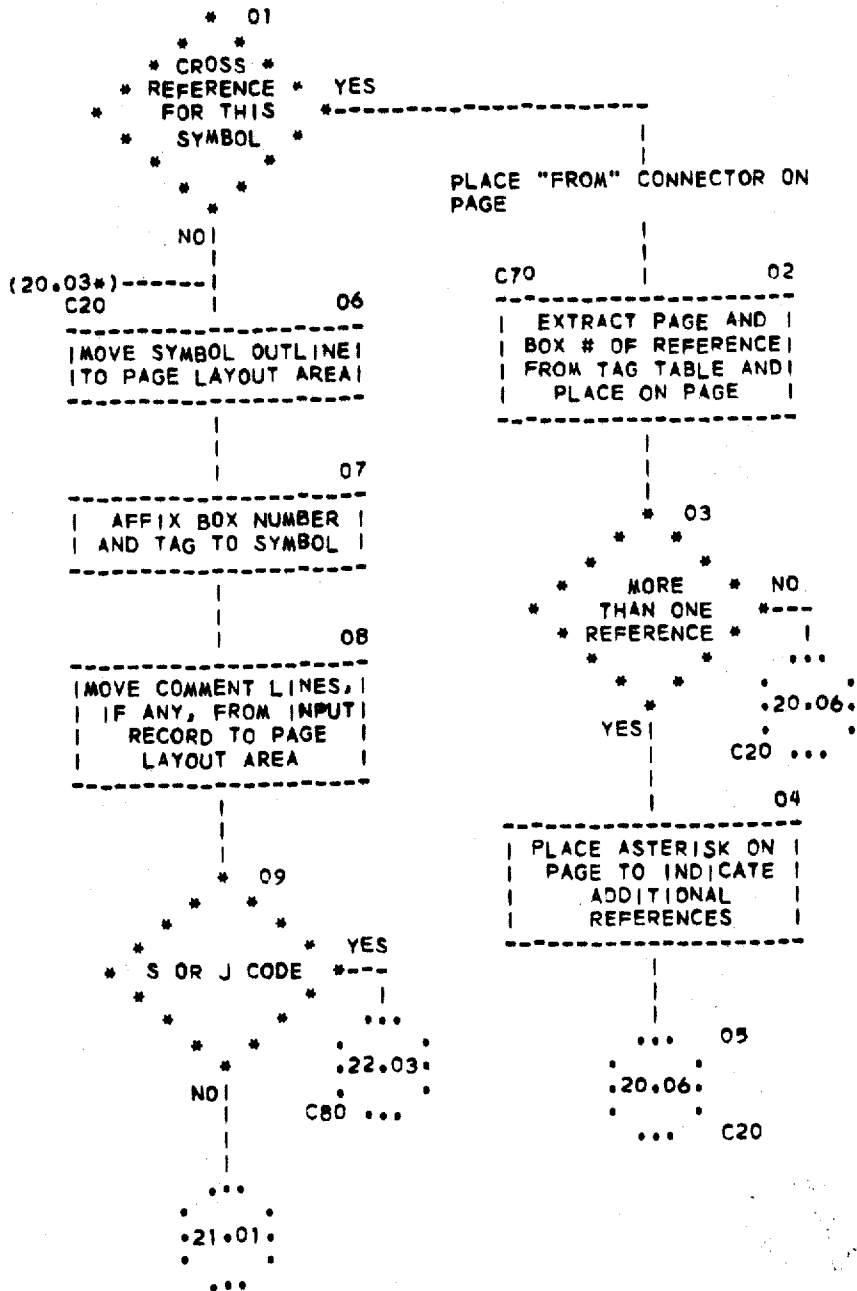
38 Sheets-Sheet 28

Fig. 8/20

PROGRAM: AUTOFLOW

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CHART: PASS THREE - LAYOUT



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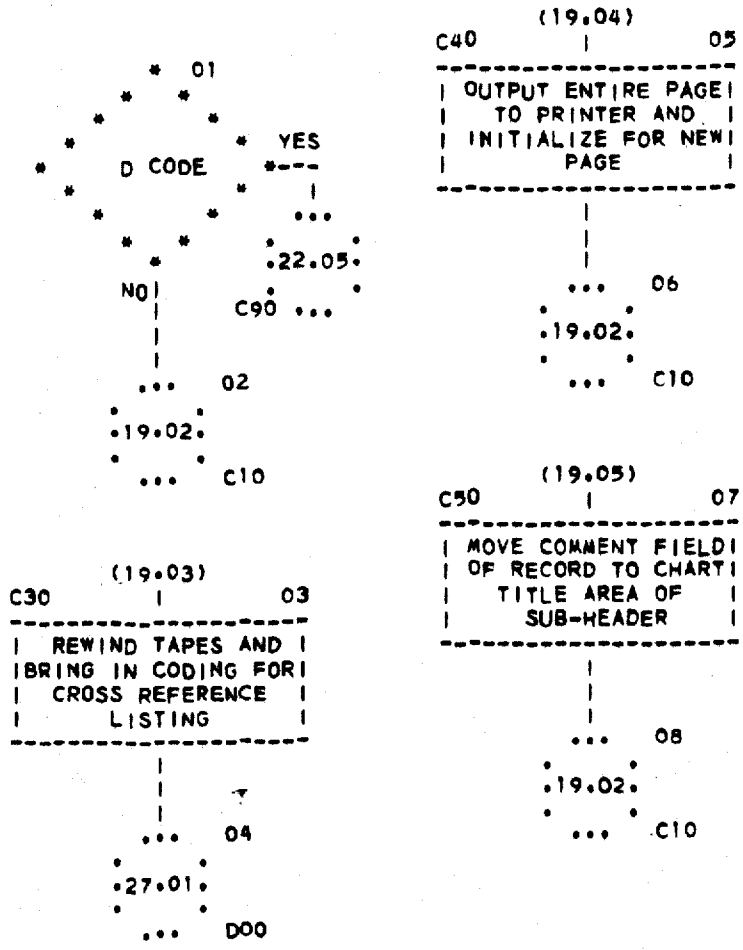
38 Sheets-Sheet 29

Fig. 8/21

PROGRAM: AUTOFLOW

PAGE 21

CHART: PASS THREE - LAYOUT



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Fig. 8/22

PROGRAM: AUTOFLOW

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CHART: PASS THREE - LAYOUT

```

(19.07)
C60 | 01
-----|
| PLACE PSEUDO |
| CONNECTOR SYMBOL ON |
| PAGE - PAGE AND BOX |
| # OF DESTINATION IS |
| CONTAINED IN INPUT |
| RECORD |
-----|

```

```

|
| ... 02
| .19.02.
| ... C10

```

```

(20.09)
C80 | 03
-----|
| EXTRACT PAGE AND |
| BOX # OF |
| DESTINATION FROM |
| TAG TABLE AND PLACE |
| INSIDE SYMBOL |
-----|

```

```

|
| ... 04
| .19.02.
| ... C10

```

PROCESSING OF DECISION RECORDS

EACH DECISION RECORD HAS A MAXIMUM OF TWO BRANCHES WHICH MUST BE ANALYZED FOR SHOWING CONNECTION.

```

(21.01)
C90 | 05
-----|
| SET LEFT AND RIGHT |
| SIDE SIGNALS TO |
| "OFF" |
-----|

```

```

(24.05)
C100 | 06
-----|
| LOOK UP BRANCH TAG |
| IN TAG TABLE |
-----|

```

```

|
| * 07
| * *
| * TAG * NO
| * FOUND IN *
| * TABLE *
| * *
| * * .25.01.
| YES | C130 ...
|
| ...
| .23.01.
| ...

```


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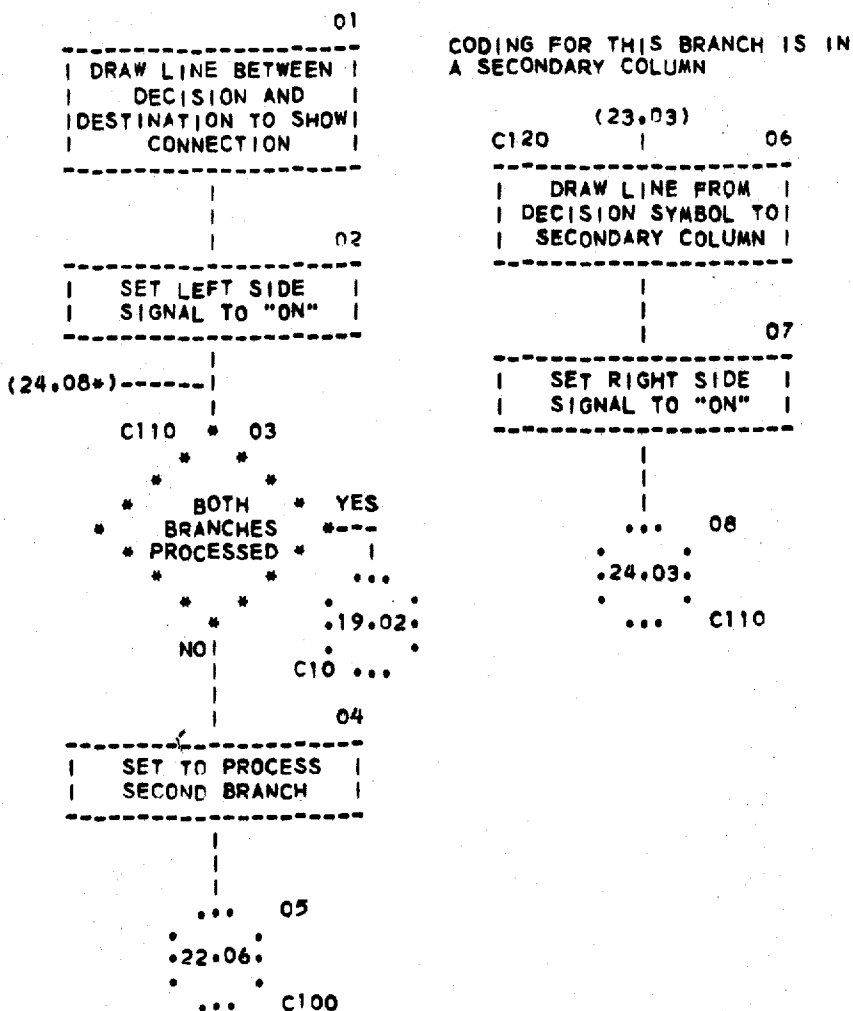
38 Sheets-Sheet 32

Fig. 8/24

PROGRAM: AUTOFLOW

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CHART: PASS THREE - LAYOUT



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38 Sheets-Sheet 33

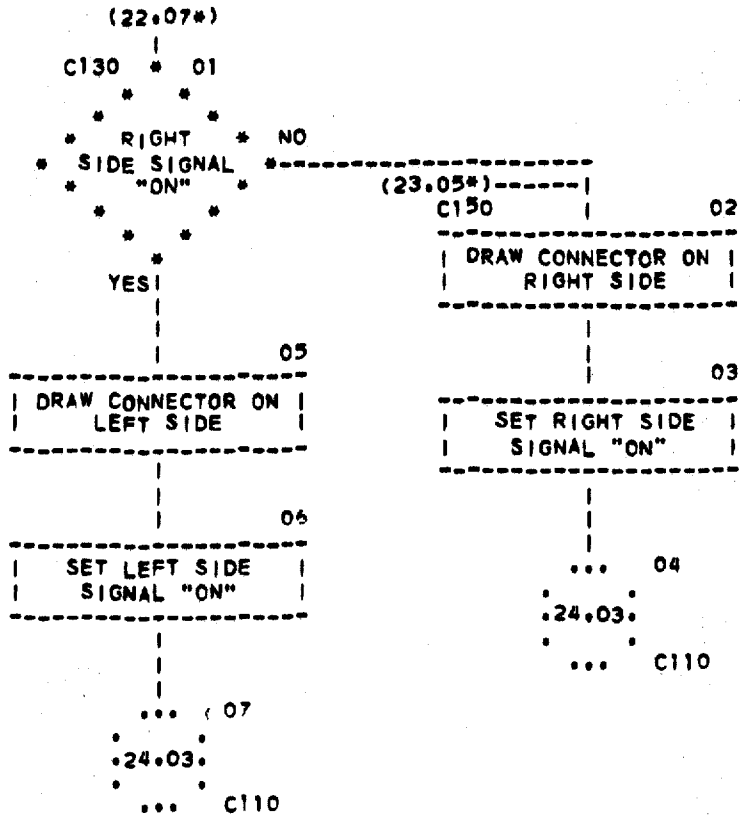
Fig. 8/25

PROGRAM: AUTOFLOW

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CHART: PASS THREE - LAYOUT

USE A CONNECTOR TO SHOW
CONNECTION



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Fig. 8/26

PROGRAM: AUTOFLOW

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CHART: PASS THREE - LAYOUT

INTERFERING LINE

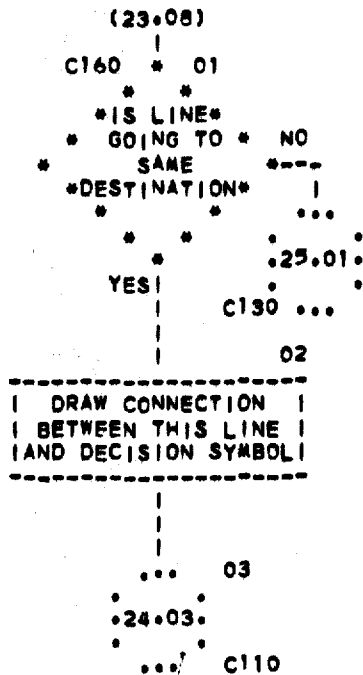
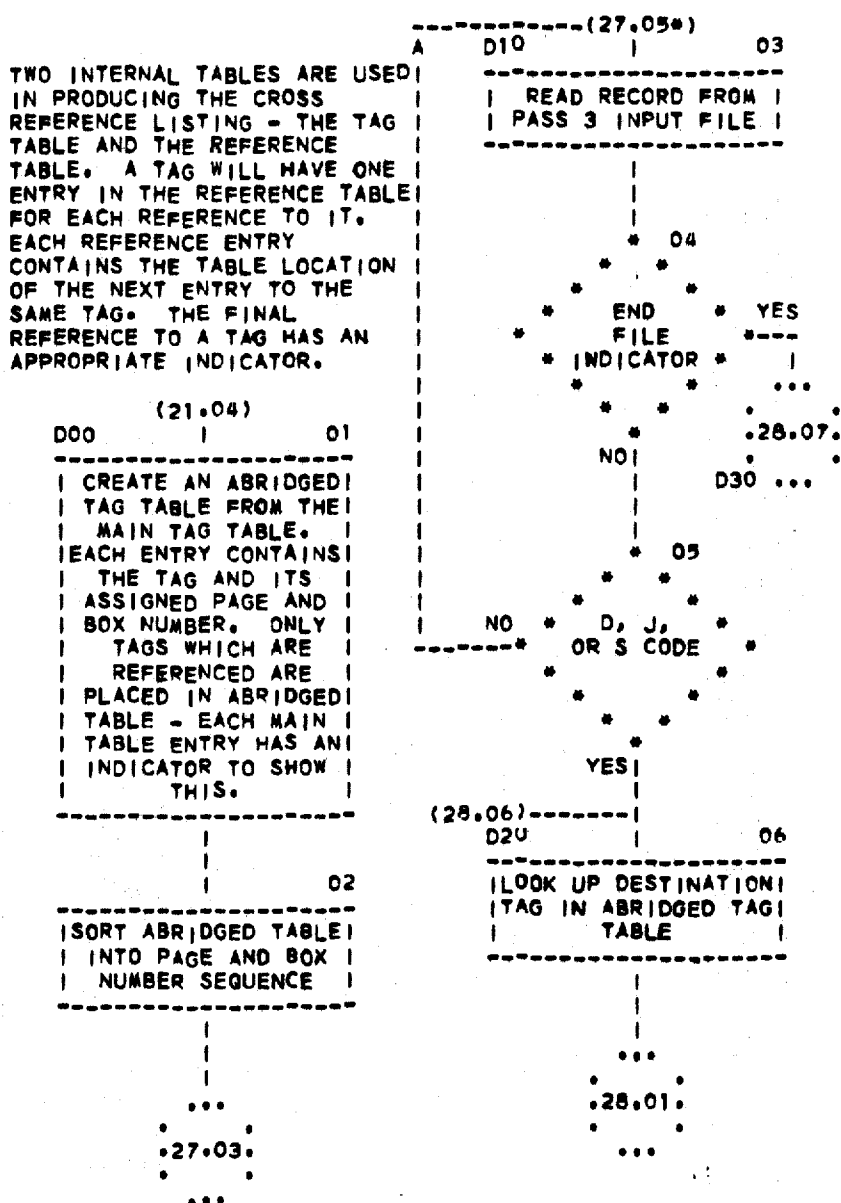


Fig. 8/27

PROGRAM: AUTOFLOW

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CHART: CROSS-REFERENCE LIST



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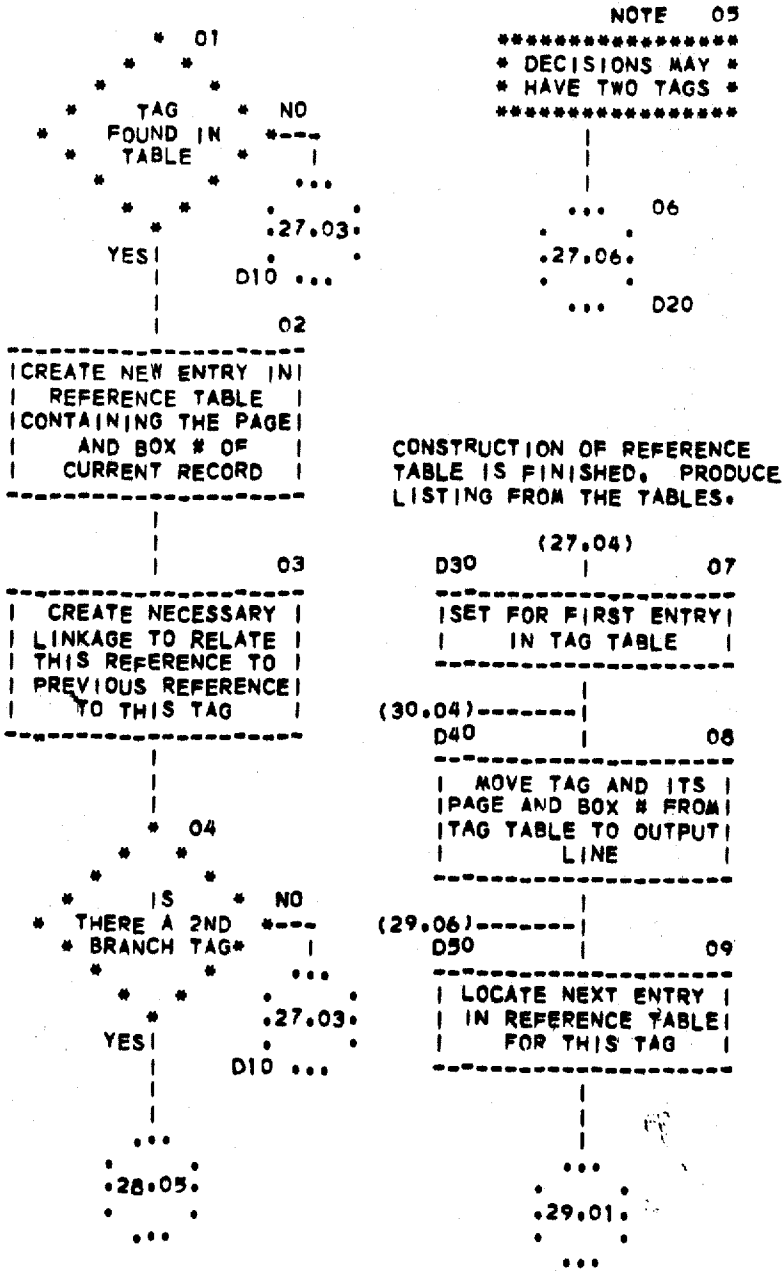
38 Sheets-Sheet 36

Fig. 8/28

PROGRAM: AUTOFLOW

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CHART: CROSS-REFERENCE LIST



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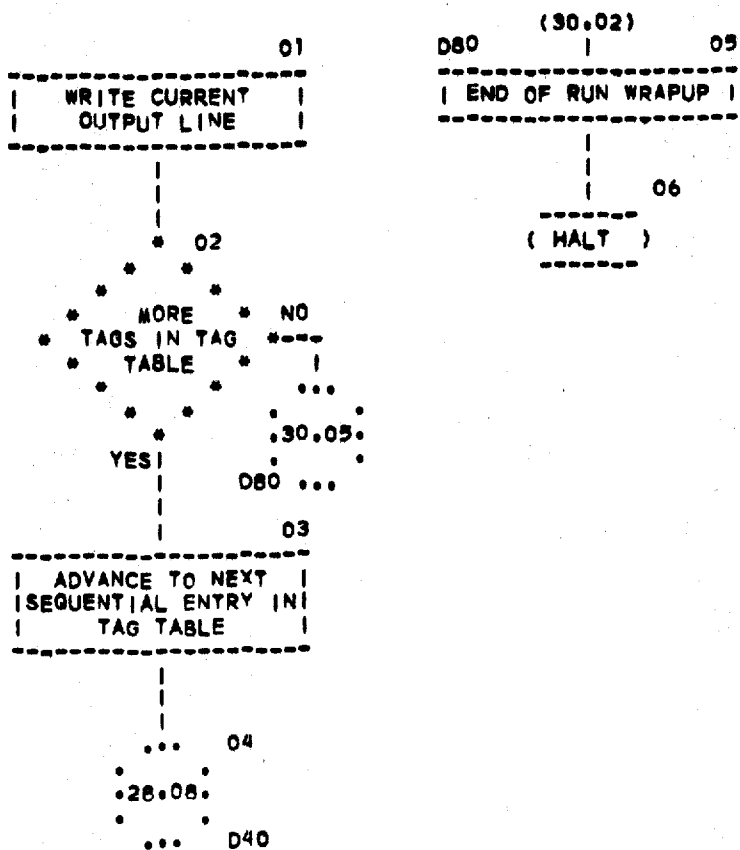
38 Sheets-Sheet 38

Fig. 8/30

PROGRAM: AUTOFLOW

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CHART: CROSS-REFERENCE LIST



1

2

3,533,086

AUTOMATIC SYSTEM FOR CONSTRUCTING AND RECORDING DISPLAY CHARTS

Martin A. Goetz, Princeton, N.J., assignor to Applied Data Research, Inc., a corporation of New Jersey
Continuation-in-part of application Ser. No. 512,113, Dec. 7, 1965. This application Dec. 24, 1968, Ser. No. 786,782

Int. Cl. G06f 9/06

U.S. Cl. 340—172.5

45 Claims

ABSTRACT OF THE DISCLOSURE

A data processor system for automatically making two-dimensional flow charts forms chain sequences of the flow chart symbols and allocates the symbol chains in parent and branch and sub-branch sequences as clusters to successive flow chart pages.

BACKGROUND OF THE INVENTION

This invention relates to a system for automatically constructing and recording display charts and particularly flow charts representative of control systems for digital computers.

This application is a continuation-in-part of copending application Ser. No. 512,113, filed Dec. 7, 1965, now abandoned.

Computer programs that are used to control the sequential operations of digital computers are made up of sequences of hundreds or thousands of computer instructions or commands which have complex interrelationships. The relationships of these instruction sequences, whether presented in machine coding or in machine dependent or independent languages, are difficult to interpret, even when they are read by skilled programmers. For that reason the program is generally presented in the form of a flow chart, which graphically presents the logic flow of machine operation and enables the programmer and users of the program to more readily interpret and understand the program. When a programmer constructs a new program he may develop rough sketches of a flow chart prior to implementation of the program, but commonly, such sketches are an inadequate description of the final program that is implemented, which may incorporate numerous changes and revisions. Moreover, a draftsman is needed to convert the sketches to suitable drawings, and the drawings, in turn, should be checked to ensure that no errors have been made in the transcription. Due to the tediousness of making a good flow chart, the pressure of other duties, and changes in personnel, the flow chart documentation of a program by the programmer is often incomplete and inaccurate. Yet, without reliable flow chart documentation, skilled personnel who were not involved in the original design of the program have great difficulty in learning and understanding its construction and operation, and in developing modifications and variations of the program, as circumstances often require. In addition, as a program is updated or revised, procedures are needed for readily updating the flow chart documentation.

SUMMARY OF THE INVENTION

Accordingly, it is among the objects of this invention to provide a new and improved data processing system for automatically producing flow chart documentation of a computer program.

Another object is to provide a new and improved automatic flow chart documentation system for computer programs that automatically produces from a computer program a flow chart which is an accurate and informative graphical representation of the program.

Another object is to provide a new and improved automatic flow chart documentation system for computer programs that relieves the programmer of documentation chores and makes it possible to obtain documentation immediately upon the program being constructed.

Another object is to provide a new and improved automatic flow chart documentation system for computer programs which assists a programmer in debugging the programs that he constructs and in revising the program as may be required.

Another object is to provide a new and improved flow chart documentation system for computer programs by means of generally available digital computers.

Another object is to provide a new and improved method of operating digital computers to produce flow chart documentation of computer programs.

Another object is to provide a new and improved computer programming system for operating stored program computers to produce automatically flow chart documentation of other computer programs.

In accordance with an embodiment of this invention a computer program is provided for operating a stored program digital computer to perform the flow chart documentation of other computer programs. The computer program to be documented is in the form of groups or combinations of digital signals that are treated as data by the digital computer when it is operated in accordance with this invention. The digital computer operates on each of the successive groups of data signals representing the sequences of instructions or instruction groups of the program to be documented, and determines therefrom what type of instruction is represented and the length of the column display required for presenting each instruction as a diagrammatic block along a column of a flow chart. Chains of such blocks between successive transfer types of instructions are established and the length of the chain for display on the flow chart page is determined. Destination tags in the data blocks are identified, and a tag table is developed of those tags and their relations to the associated blocks. The locations of successive chains in a main path of the program are allocated to successive columns of the display pages. Branch instructions of those chains are handled specially by identifying the destination (or tagged) chain to which the program branches and, where adequate space is available for the tagged chain in a column adjacent to the column of the main path, allocating the tagged chain to the adjacent column. Chains are allocated in sequential order to the main flow except where branch instructions are encountered; for the latter, the main path allocation is interrupted to allocate the branch chains. The locations of the connecting paths between blocks in the same and adjacent columns and of connectors to blocks in non-adjacent columns of the same or different pages are established. A display is provided of the interconnected chains with connecting paths where possible to represent paths between blocks in the same and adjacent columns and connector symbols are drawn where such paths cannot be drawn.

In other embodiments, modified forms of the invention are used.

BRIEF DESCRIPTION OF THE DRAWING

The foregoing and other objects of this invention, the various features thereof, as well as the invention itself may be more fully understood from the following description when read together with the accompanying drawing, in which:

FIG. 1 is a schematic block diagram of a data processing system incorporating a control program for flow chart documentation in accordance with this invention;

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FIG. 2 is a schematic block diagram of a fragment in outline form of the type of chart that is produced by means of this invention;

FIG. 3 is a general schematic block diagram of a computer program embodying this invention and used in the system of FIG. 1;

FIG. 4A and B are schematic block diagrams of outline fragments of a form of computer program flow chart produced in accordance with modifications of this invention;

FIG. 5 is a schematic block diagram of an outline fragment of another form of computer program flow chart produced in accordance with modifications of this invention;

FIGS. 6A, B, C, D and E are schematic graphical diagrams of flow chart patterns produced in accordance with another modification of this invention;

FIGS. 7A, B, C, D and E are a series of schematic flow chart diagrams that together illustrate graphically details of a modified form of computer program embodying this invention and used in the system of FIG. 1; and

FIGS. 8/1 to 8/30 are a series of schematic flow chart diagrams that together illustrate graphically details of a computer program shown in FIG. 3.

In the accompanying drawing, corresponding parts are identified by similar reference characters throughout.

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data stored in the memory **102** is supplied under the direction of the controls **106** to the arithmetic unit **104** for processing and then returned to the memory to be stored at appropriate locations. The controls **106** determine the memory locations from which information is taken to be processed in the arithmetic unit **104** as well as the memory locations to which it is returned after processing, and also determine the timing of the flow of electrical signals representing the information. The controls **106** also determine the particular operations performed by the arithmetic unit as well as the time interrelations thereof. The flow of signals to and from the memory is indicated by solid lines, while the flow of control signals from the controls **106** to operate the memory **102** and the arithmetic unit **104** is shown by broken lines. In practice, various kinds of such control lines are required, and the details are omitted here since appropriate arrangements are well known in the art and they are unnecessary to an understanding of the invention.

The data processing system of FIG. 1 may be any of various well-known types of systems such as that employing a stored program in the form of signals stored in the memory and representing sequences of control instructions or commands which select the operations to be produced by the controls **106** as they are needed to perform the desired operations. Such a stored program is effectively a part or extension of the controls **106** and is commonly stored in the memory in a section set apart for that purpose. A section **102a** of the memory **102** is labeled to indicate that the control program for the documentation of a flow chart (F/C) is shown as part of that memory. Alternatively, the controls **106** may have a fixed program wired and/or built-in which may take the form of logic combinations of gates and other circuits to perform the proper sequence of operations that make up the logic of the program, all in accordance with techniques that are well-known in the art. The operations required of the flow chart documentation system of this invention are generally large in number and have complex interrelationships. Therefore, a stored program is the preferred form of control system for presently available computers, and an embodiment of such a program is described below.

The input portion of the data processing system is represented as a magnetic tape unit **108** operated by appropriate signals from the controls **106** to supply groups of combinatorial signals to the input area **102b** of the memory **102**, which input signals represent the program which is to be documented. That is, the data inputs of the system are themselves successive sections of another computer program, which are processed as data to provide a flow chart representation of the logic of the input program. Another magnetic tape unit **109**, operated by the controls, carries the signals that form the flow chart control program, and this program is read into the memory section **102a** when the computer is to be operated in accordance therewith. The entire F/C program may be read into memory section **102a** at the beginning; or, since the program is formed as a sequence of subdivisions or passes, the subdivisions may be separately read into memory as required by the program. Other magnetic tape units **110** and **111**, operated by the controls **106**, receive records from the memory output area **102c** and supply these records back to the input area **102b** thereof during different stages of operation of the system. In use, the tape units operate effectively as portions of the computer memory system. An output display device **112** is also operated by the controls **106** to produce a graphical display of the final flow chart produced by the system. This display device **112** may be a high speed printer (e.g. a line-at-a-time printer), a digital plotter recorder, a cathode ray tube display or recording system, or any other appropriate form of display or recording system. The display device may be operated on-line directly from the memory **102**, or the output from the memory may be

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GENERAL SYSTEM

In the embodiment of the invention in FIG. 1, a data processing system constructed in accordance with this invention is shown, and it includes a memory **102**, an arithmetic unit **104**, and a set of controls **106**. This system may be a suitable form of digital computer in which

written on to an appropriate tape and used to operate the display device off-line in any suitable manner.

FLOWCHART SYMBOLS AND FORMATS

FIG. 2 illustrates an outline of a flow chart that is produced in accordance with the system of this invention. The flow chart documentation program of this invention examines specific fields of each instruction line and other data sections of the program to be documented and produces a standardized flow chart by means of the display device 112. In accordance with one form of the invention, the input program is in an assembly language and the flow chart that is produced is divided into four columns (the invention may be used to supply flow charts of any desired size having one or more columns). The first and third columns 114 and 116 may be used for depicting the main flow of the program and the other two columns 118 and 120 would then be reserved for branches from the main flow columns. However, as described below, it is preferred to have all four columns available for display of the main flow logic of the program and to use the next adjacent column for display of the logic that branches from a main flow column. Various types of graphical symbols are used to represent the different types of instruction or instruction groups of the program that are being documented. One such symbol is a rectangular box 122 representing a PROCESS; another symbol is a diamond-shaped box 124 that represents a DECISION or branch (that is, conditional transfer) instruction. Seven other different symbols are utilized and illustrated in FIG. 2 as is described hereinafter. Symbols are assigned numbers sequentially beginning with "01" on each page, and all cross-referencing to other symbols is in the form "XX.YY," where XX is the page number and YY the number of the symbol on that page, which provides for 99 symbols for each of 99 pages, and which can be readily modified for larger numbers if needed. Symbol sequence numbers are printed above and to the right of all symbols on the chart, as shown in FIG. 2, where box 122 carries the symbol sequence number "01." The symbols are numbered in sequence from top to bottom of a column, and, in the main flow columns, from the bottom (or exit) of one column to the top (or entry) of the next reading from left to right. Branches from a DECISION symbol are from the side corners of the diamond, and these branches may be connected to the entry point of another symbol in the same column (see line 123) or to the entry of a symbol in the adjacent column (see line 125). The symbols in a branch column are sequentially numbered (in one embodiment) starting from the number of the DECISION symbol from which it branches; and the symbols below that DECISION in the main flow column have numbers that continue after those of the branch column. Thus, the number of branch block 126 is "03" following that of DECISION 124, and the number of block 127 is "08" following that of the last symbol in branch column 118.

The documentation program in one embodiment, examines four fields of an input program presented in a fixed formal assembly language and produces the flow charts therefrom; these fields are the symbolic tag, the comments of the programmer interpreting each instruction or group of instructions, a special flow chart code located in a predetermined part of the Comments field, and operands (such as the operation code and certain addresses) that supply the destination tags or addresses of transfer and branch instructions. The following eleven flow chart codes are utilized in such an assembly language format for the system of FIG. 2:

"P"—PROCESS	"H"—HALT
"C"—CONTINUATION	"N"—NOTE
"S"—SUBROUTINE	"B"—CHART TITLE
"J"—JUMP	"D"—DECISION
"E"—EXIT	"W"—SWITCH
	"T"—TEXT

Eight different chart symbols are used to represent the eight different classes of data processing, and "B" and "T" identify TITLE and TEXT that are to be printed on the chart. "C" is used to identify a continuation of a comment that started in a preceding record of the program. It has been found convenient to insert the F/C code after the Comments field separated therefrom by spaces. Due to the records being fixed in size, the number of characters usually being that of a punch card, successive records are used to carry extensive comments.

In FIG. 2 the PROCESS (or P-code) symbol is shown as a rectangular block, such as the block 122, and the Comments portion of the corresponding instruction is inserted in the symbol in the manner illustrated in FIG. 2 and as shown in greater detail in FIG. 8 (the latter may be referred to as an illustration of a flow chart, in two-column form, produced by this invention). This PROCESS symbol is variable in column length depending on the length of comments. The TITLE represented by code-code-B is placed on the top of each page of the flow chart, as indicated at the top of column 114 of FIG. 2. The T-code for TEXT indicates that the textual material supplied in the assembly language is to be printed out without a special chart symbol (see column 116). In column 114, diamond-shaped symbol 124 represents a DECISION, and the accompanying comments are edited and inserted in that symbol as indicated. Either one or two branches of coding may be shown as coming from the side corners of the DECISION symbol. Labels are supplied to the branch paths from the DECISION symbol in accordance with a special code or in accordance with the Comments field of the input instruction. The documentation program of this invention determines how the lateral branches are to be depicted on the chart; they may be shown by a horizontal line 125 connecting to an adjacent column as indicated by the branch EQL from the box 124 to process box 126 in column 118; or by a connecting line 123 from the LOW branch of box 124 down to the input of block 128 in column 114; a third method (where such connecting lines cannot be drawn) is by means of a connector symbol in the form of a circle 130 connected to the DECISION branch and containing the cross-reference identification XX.YY of the flow chart symbol to the input of which it connects. Thus, in the case of connector 130, the HIGH branch from the DECISION block-10 goes to block-13 on page 1, which is shown in column 120.

Another symbol is that for SUBROUTINES (S) such as that shown by block-05 in column 118. This is a hexagon-shaped box in which the Comments field is written, and in a separate section in the left of the box in the cross-reference is given to the location of the details of the subroutine. As illustrated, SUBROUTINE-05 is cross-referenced to 01.20, which one can readily locate in column 116 on the same page. The cross-reference back to SUBROUTINE block 01.05 is shown at the input to box 01.20 so that the reader of the flow chart may readily determine the entry into that section of the program and interrelate the different positions thereof.

The JUMP code (J) is represented by a circle as indicated by block-07 at the bottom of column 118. The circle has at the lower right the tagged destination to which the program jumps, and contains within it the page and box number of that destination in the flow chart so that it can be readily located. The tag of the entry point of a block is at the upper-left of a block, as the tag, "JAN" for block 126.

The E-code representing EXIT is illustrated by a circle that terminates a chain of blocks and contains the word "exit," as shown by circle-23 in column 120. Similarly, the H-code for HALT terminates a chain of coding as shown by circle-17 in column 116. The N-code, used for NOTES, is represented by a rectangular block (see block-5) which can be varied in column length to contain the associated comments, and which is offset with its left

side indented to distinguish it from a PROCESS rectangle.

The W-code is a diamond-shaped symbol similar to a DECISION and it contains the word "switch." Its branches are handled in a similar fashion to the DECISION box as illustrated by the example of block-14 in column 120.

The destination tags where they are provided in the comments field of the input program (or otherwise in the operand codes of that program) are picked up and supplied to the transferring symbol, where it is not directly connected to the destination symbol. For example, block 130 is a connector symbol that has the destination tag "CAN" printed next to it, for that tag is the input of block-13 in column 120. "CAN" is also printed at the input of block-13 and if the other blocks have input tags they are similarly printed. Cross-references to show the originating points for entries to blocks of the flow chart are by way of the page and block number of each symbol. An example is the cross-reference back to block 01.05 that is set forth in parentheses at the input to block-20 in column 116; another example is the tagged entry "RAN" to block-05 that is shown in parentheses (namely, 01.10), which cross-references back to the branch from connector 132 in column 114. In addition where more than one cross-reference entry exists, rather than indicating all of the symbol locations, an asterisk is provided, as at the entry of block-13 in column 120, and a separate cross-reference table lists all such entries, as is discussed in further detail herein below.

In accordance with the documentation program of this invention, each chain of symbols is terminated by a JUMP or EXIT symbol. In addition, where space limitations do not permit the printing of a long chain of symbols until it terminates in that fashion, a special connector symbol is generated by the F/C program to set forth the continuity of the main flow of the program. Thus at the bottom of the column 114, such a pseudo-connector symbol is shown as a circle containing the location (01.11) of the next instruction in the main path, which in the case illustrated is to block-11 on the same page, namely, the first block in column 116. Also at the bottom of column 116 a pseudo-connector circle is generated containing a cross-reference to the first block of page 2, for the next symbol of the main flow of the program.

PHASES OF FLOWCHARTING SYSTEM

FIG. 2 is referred to hereinafter to illustrate the development of the flow chart as the operations called for by the documentation program are sequentially performed. FIG. 3 is a block diagram of the four main phases of the F/C program. In the specific embodiment of the invention described hereinafter, these four phases correspond to four separate sections of the program and four passes of data. During Pass I, block 140, the editing phase is performed in which the input data is accepted and edited, the column lengths of individual symbols are fixed, individual chains of logic are established and their lengths determined, and a Chain Table and the skeleton of a Tag Table are constructed. During Pass II, block 142 output records of Pass I are processed to complete the Tag Table; individual symbols are assigned locations on successive pages of the flow chart, and assigned to a specific part of a column within the page. Successive main flow chains are processed along with the branch chains. During Pass III, block 144, the Tag Table is used to print a table of contents, the successive pages of the flow chart are laid out, connecting lines and symbols set out, and each page is printed. During an additional pass, Pass IV, block 146, a table of cross-references is developed and printed out. The program of this invention is not limited in its form to any particular number of passes of the data; the particular number varies with the computer that is used and with the availability and division of memory space (say, in the random-access section of the memory as against drum or tape memory storage) in the computer

as well as with the complexity of the processing operations that can be performed by the computer and of the flow chart that is desired.

INPUT DATA

The input data for the assembly language program generally takes the form of successive records corresponding to the quantity of information that can be developed on an individual punch card. This record may have the information arranged in any prescribed order, and for one assembly language, it contains the following ordered fields of information:

Assembly line #; tag, if any, of current record; operation code; A-Address; index instruction code; B-Address; Comments.

The Comments field carries the programmer's interpretation of the data processing operation, so that the operation code is not needed and is not used for that purpose in the present embodiment; nor is the index instruction code and B-Address. However, for the purpose of picking up destination tags, if they are not carried by the Comments, the operation code and A and B-Addresses are examined. That is, the operation code is examined to determine if it is a conditional branch instruction, and if so, the A-Address is used for one branch destination and the B-Address for any second branch destination that may be involved. The Comments field may be used to carry any desired destination tags to which the current record transfers, which destinations are set off parenthetically at the beginning of that field by asterisks, e.g. as follows: *TAG*. In addition, the labels for the branches of DECISION symbols may be supplied with those destination tags in a special format described below. In addition, in accordance with the present embodiment of the documentation system, an additional field of flow chart code is provided; that is, one of the aforementioned eleven flow chart code characters; and, in the case of DECISION OF D-codes, an additional optional code character may be provided representing the different classes of labels for common DECISION branches, as follows:

"Y"—YES; NO
 "Z"—NO; YES
 "I"—HIGH; LOW
 "L"—LOW; HIGH
 "Q"—EQUAL; UNEQUAL
 "U"—UNEQUAL; EQUAL
 "#"—PLUS; MINUS
 "—"—MINUS; PLUS
 "3"—EQUAL; HIGH; LOW

Program Logic

FIG. 8 is a schematic flow chart diagram of the general logic of the documentation program; it consists of 30 sections or pages of flow chart, and FIG. 8 is numbered as FIG. 8/1 to FIG. 8/30 to identify those 30 flow chart pages. Each of these pages of FIG. 8 is presented in the form of flow chart that would be produced under the automatic control of the documentation program itself; and, in fact, that program was used to develop the flow chart of FIG. 8. Due to the size limitations of the patent drawings, only a quarter page of the high-speed-printer page was utilized, which afforded room only for two (of the available four) columns in width and a half-column in length. Thus, FIG. 8 illustrates an actual flow chart developed by the documentation program, but is much simplified in that less information is provided on a single page. Reference is made to FIG. 2 for a representation of the four-column format that is produced by a preferred embodiment of the documentation program; and FIG. 8 illustrates that a two-column format may also be produced with minor modifications of this program. Each of the pages of FIG. 8 has a program title and chart title together with a page number in the same fashion as is developed by the documentation program itself. As explained above, each of the blocks is referenced as XX.YY,

where XX is the page number and YY is the block number thereon.

PASS I

In FIG. 8/1, following the program and chart titles and page number at the top, a body of text material sets forth the functions of Pass I as that of editing the source information and constructing Tag and Chain Tables. Block **01.01** is a PROCESS symbol whose Comment field indicates that its function is that of setting the record sequence number counter to zero. For simplicity of illustration, the flow diagram of FIG. 8 omits certain other preliminary operations such as those known as "house-keeping" operations, and those of calling in Pass I of the F/C program, clearance of memory areas where required, and entry of constants, and the like, which are routine in nature, and which would be readily apparent to those skilled in the art and are not needed for an understanding of the invention.

The main flow continues with the second block **01.02**, and the process of reading the first source record is performed. This operation of the computer may call for a series of detailed instructions, depending on the computer construction, by which the input tape **108** is controlled to operate momentarily, and the first available record of the input program is read and stored in a predetermined order so that its fields in a fixed format are placed in prescribed sections of the input area **102b** of the memory **102**. The first record that is read in may be assumed to be a special control record, which may be marked with a special code identification if desired, and which should contain the name of the program. This program title is stored, under the operation of block **01.03**, in a prescribed primary storage area of the memory for use by Pass III in composing and printing out each page of the program; then this first card is dropped.

The next available source record is read, block **01.04**, from the input tape **108** into the input area **102b**; and the field for the F/C code is checked, block **01.05**, to determine if it is a valid F/C code other than "C." If the check indicates that the answer is NO, the program branches in a loop back to block **01.04** to read the next available input record into the input area **102b**, and DECISION **01.05** checks the F/C code of that record in the same way. This loop is repeated until the answer to the check is YES, and the program then continues on the block **01.06** (as indicated by the reference in the pseudo-connector symbol at the end of the first column of FIG. 8/1). As indicated by NOTE **01.06**, the input memory area **01.06** now contains a source record that is to be processed, since it contains a valid F/C code other than "C." A C-code is not processed, where it does not follow another code type.

The main-flow logic of Pass I then begins with block **01.07**, and all of the fields of the first record to be processed are moved from the input area **102b** to corresponding sections of the work area **102d**. That is, the following fields of that record from the memory input area **102b** are moved to individual memory work areas **102d** to which these fields are assigned: assembly-line #; tag; Comments field; F/C code; and, if a transfer instruction, the operand that contains the destination tag. Prior to moving the F/C code to its work area, the previous contents of that F/C code work area are transferred to a field of the primary storage area **102e** in memory, identified by the mnemonic LSTCD; this previous F/C code is used under a certain condition, as is explained hereinafter. The next block **01.08** controls the reading of the next available source record into the input area **102b**, which is now free to receive it since the previous record was moved to the work area **102d**. The two memory areas **102b** and **d** each contains a record, which records are processed in the order they were received.

The next block, DECISION **01.08**, checks to see if the record in input area **102b** contains an identifier character

EF indicating the end-of-a-file, which is commonly provided in all magnetic tape systems in one form or another. If the result of this test is that such an end-of-file identifier is in the record, the program branches to PROCESS block **04.08** (the last block of FIG. 8/4) which calls for the setting of an end-of-pass switch. After completion of the processing of the record that lies in the work area **102d**, that switch is tested (block **04.07**) and if it is set, operations of block **06.07** are initiated for rewinding the input and output tapes **198** and **110** and calling in Pass II of the program; thereafter, by means of the JUMP instruction **06.08**, the main flow of the program is transferred to block **07.01**, the first block on FIG. 8/7, at the Pass II starts. However, if DECISION test **01.09** indicates that there is no end-of-file identifier in the record in input area **102b**, the program steps to DECISION **01.10** to test if there is a valid F/C code in that field of the input record that sits in the input area **102b** of memory. If the answer is NO (that is, if it is an invalid character or a space), the program ignores that record and returns as a loop to PROCESS block **01.08**, and the next source record is read. This loop is repeated until the next record is found that contains a valid F/C code in the proper field of the input record. When it is found, the main flow of the program continues with DECISION **02.01** (FIG. 8/2) which determines if the F/C code of the record in the input area is a "C," which identifies a CONTINUATION record. If input area is a "C," the program branches to block **02.02** where the process is performed of appending the Comments field of that CONTINUATION record to the Comments field of the previous record lying in the memory work area. A CONTINUATION record serves only to carry continued lines of Comments of a previous record; it is not otherwise processed, and except for the Comments, it is not moved to the work area **102d**. Thereafter, as indicated by JUMP **02.02**, the program then loops back again to block **01.08**, and the same process is performed on the next record. If it is another C-code, the Comments field of that input record is again appended to the Comments field of the previous record lying in memory work area **102d**. This loop is repeated until another record is established in the memory input area which contains an F/C code other than "C" or space. At that time, as indicated by NOTE **02.04**, memory work area **102d** contains all of the necessary information for processing the code record therein. The record in the input area **102b** is left there and it is processed after the record developed in the working area has been processed. A locator RHECOM is used to maintain the memory location of the right-hand end (RHE) of the Comments work area, since this work area increases as additional CONTINUATION records are read and appended. This locator is a field of the primary storage area **102e** which contains the address in the work area of the RHE of the Comments. When the record in the work area is complete and ready to process, a control symbol is inserted into the memory location at the RHE of the Comments to mark the end of that record.

Blocks **02.05** and **02.06** begin the processing and they represent a number of processing operations that are performed to prepare for the output area **102c** to receive the record from the work area. The output record of Pass I has a certain format, which is set forth in Table I below. The same format is used for the output record of Pass II, and Table I indicates which Pass is used to fill each field. The beginning and end of each "line" of the Comments field is denoted by control symbols since the "lines" are of variable length, as explained below. A maximum size of the Comments field is arbitrarily set at about 15 lines of TEXT. The output memory area includes the following prescribed memory fields for receiving the characters that compose each record. The number of characters specified in the following table are those found suitable for particular embodiments; the number of characters may vary for other embodiments. Each output record

assembled in the output area is written successively on work tape 110, so that the latter contains all of the records in order at the end of Pass I; in turn, work tape 110 becomes the input to Pass II. Each output record is processed successively and completed during Pass II and then written out to the output file on tape 111.

TABLE I.—OUTPUT RECORD

Field	No. CHAR	Filled in by—	Remarks
No. Lines.....	1	Control Symbol.
Column No.	3	Pass I.....	No. lines required on a F/C page to contain this record.
Line No.	1	Pass II.....	Column No. on F/C page assigned to this symbol.
Page No.	3do.....	Line No. on F/C page at which symbol begins.
Box No.	2do.....	F/C page No. assigned to this symbol.
Code.....	2do.....	Box No. assigned to this symbol.
Code.....	1	Pass I.....	Type of symbol.
Tag.....	7do.....	Blank if record has no tag.
"LHTC".....	1	Pass II.....	Connection indicator for field "LHT".
"LHT".....	6	Pass I.....	Contains destination tag on D, J, S codes—blank for other codes.
"RHTC".....	1	Pass II.....	Connection indicator for field "RHT".
"RHT".....	6	Pass I.....	Second destination tag on DECISIONS (may be blank).
SEQ No.	4do.....	Ascending Seq. No.
Length.....	4do.....	Length—lines of this record.
Comments.....	Variabledo.....	(Max. of 456 char.).
			Control Symbol.

As indicated by block 02.05, the record-sequence-number counter is incremented by "1," and its new number is moved to the prescribed field of output area 102c. The processing continues with block 02.06, and the F/C code is moved from the work area to the corresponding field of the output area. Also several memory fields are set to initial conditions by block 02.06: Line Counter-A in primary storage 102e is set to 0; Comments Locator is set to the initial address (which is a constant) of the left-hand-end (LHE) of the Comments portion of the work area 102d, and Output Area Locator is set to the initial address of the LHE of the Comments portion of the output area 102c. All other output fields, other than the Comments field, are fixed and have predetermined addresses, and therefore, require no locators.

The main flow of the program steps to block 03.01 (as indicated by the pseudo-connector at the bottom of FIG. 8/2) where a test is made on the F/C code. If it is a B-code, the program branches to block 03.02 where construction of the Tag Table in primary storage is initiated for the first such B-code; and each succeeding B-code initiates a new section of the table for all tags following it. The entry in the Tag Table consists of insertion of the title of the chart, and upon printing out of the table of contents during Pass III this chart title is printed out as a heading for the tags associated with that chart. Thereafter, as indicated by JUMP 03.03, the program jumps to block 03.05 where the editing operations are initiated.

If the test 03.01 determines that the F/C code is not a "B," the program steps along the main flow to DECISION 03.04 where a test is made for a tag in the record, which identifies the entry point of that record. If it does not have a tag the program steps to block 03.05 for the editing operations; if it does have a tag the program branches to block 05.01, which controls the creation of a Tag Table entry. The Tag Table is begun by Pass I and completed by Pass II; it is used by Pass III and this remains as part of the primary storage area 102e of the

memory throughout the program. Each entry in the table consists of the following items:

TABLE II, TAG TABLE

- (a) Tag
- (b) Assembly line # of tag
- (c) Sequence # of output record (from Pass I) containing this tag
- (d) Page and box number assigned to tag
- (e) Column and flow chart line number of tag
- (f) Indicator to show cross-reference by other symbols, and page and box number of the first cross-reference symbol.

Item (b) of the Tag Table is used for information purposes in composing the table of contents during Pass III. Item (c) is used by Pass II in determining whether a DECISION branch can be processed and allocated to a secondary column (which is an adjacent column in the present embodiment). Item (d) is used for cross-reference connectors in Pass III. Item (e) is used by Pass III for drawing connecting lines on DECISION symbols. Item (f) is used by Pass III in placing "from-connectors" or cross-references on the charts. The entry of items (a), (b) and (c) is performed during Pass I; items (d), (e) and (f) are entered during Pass II.

Block 05.01 creates the Tag Table entry by inserting items (a), (b), and (c) of the record currently being processed. That is, it determines whether the record has a tag, and enters that tag as item (a), enters the assembly line # from the record in working storage as item (b); and enters the current reading of the sequence number counter as item (c). The sequence number counter was stepped forward by block 02.05 to establish the sequence number of the current record. Space is left for the additional items (d), (e) and (f) of the Tag Table to be added during Pass II, and a locator (TLLOC) of the RHE of the Tag Table is incremented appropriately so that the next Tag Table entry may be made at the proper location. In addition, a control field NUMTAG containing the number of items in the Tag Table is also incremented. The tags associated with C-codes do not have to be placed in the Tag Table and in fact it is not necessary to pick them up from the input area except where the preceding input record has a blank tag area. The operation of block 05.02 moves the tag field of the record in the work area also to the tag field of the output record. Thereafter, JUMP 05.03 transfers the program to block 03.05 for the editing process. It is seen from the logic flow after the test 03.01 for the B-code that this tag routine is bypassed for such B-code records, so that their tag fields are not examined either for the purpose of entry in the Tag Table or for transfer to the output record. The T-code records are similarly tested, and the logic flow therefrom also bypasses the Tag Table routine. Since B- and T-codes do not generate a flow chart symbol, any tags that they may have are not referenced on the flow chart.

After the tag processing has been completed or bypassed as required, three subroutines (EDLIN, PAREN, CHENT) are performed on the different types of records in the various ways described in detail below. These subroutines are represented in simplified form in FIG. 8. At block 03.05 the editing routine (EDLIN) of editing successive lines of the Comment field of the current record is entered. The operations that are performed are those of editing the Comments field into intelligible lines and moving them into the output field as summarized generally by the comment of block 03.05. This operation varies with each of the code symbols and is described in detail below. As the editing is performed, the number of lines of flow chart required by each symbol is measured by block 03.06; this number is fixed for the fixed format symbols such as D, S, W, E, H and J; the number varies for the P and N symbols as well as for the lines of text identified by the T-code.

The number of lines for the current symbol (from Line Counter A) is added to the running total (in Line Counter B) for the current chain of logic flow by block **04.01**. A chain is defined as all coding between successive J- or E-codes, and includes the first group of coding so terminated. Examples of chains are marked off by slant bars in the following sequences of codes; /BTCPCCNJ/PSCPHPCDJ/PPCDE/.

Thus, block **04.01** produces a running total of the lines within any chain by means of Line Counter -B, which is cleared at the end of the chain, as noted below. The details of the EDLIN subroutine for the individual codes are discussed below.

A test is made for D-, W-, S- and J-codes by block **04.02**; and, if the current record contains such a code, a branch from the main flow is taken to a subroutine PAREN starting at block **05.04** and continuing through **06.02**. This subroutine is used to identify and extract from the Comments field, the destination tags and labels, if any, that are parenthetically included therein by asterisks, or to extract the destination tags from the input operands, which operations are described below for individual codes. After the PAREN routine or if the test of block **04.02** is negative, the program jumps to text **04.03**, which determines whether the current record contains a J- or E-code. If the answer to the test is "YES," the program branches to a subroutine CHENT of blocks **06.04** and **06.05**, which controls the construction of the Chain Table and the entry of new items in that table. The chain table contains the following three items:

TABLE III, CHAIN TABLE

- (a) Sequence number of the Pass I output record at which the chain ends.
- (b) Number of lines on a flow chart page required for the chain.
- (c) Field for an indicator character to show that the chain has been processed during Pass II.

From the definition of a chain, every record in the output record file is necessarily part of some chain, which is identified in the Chain Table by the sequence number of the last record in that chain. The Chain Table entries have sequence numbers in ascending order, since they are created sequentially in Pass I as the records themselves are processed, and the records are numbered in ascending order. Block **06.05** transfers the current value of the record sequence counter to the Chain Table to establish the Chain Table identification of the current chain then ending, and also transfer the current value of Line Counter B, which is a direct measure of the number of lines in a column required on the flow chart page for the chain. Line Counter-B is cleared to "0," so that it is in condition to accumulate the lines for the next chain. This completes the operation of subroutine CHENT, and it exits back to block **04.04**. The record in the output memory area is not complete and a record terminating symbol is added to the RHE of the record, which is then written to the output file of work tape **110**.

At this point, the aforementioned end-pass switch is tested (NOTE **04.05** and DECISION **04.06**) to determine whether it is the end of a pass. As explained above, this switch would have been set by block **04.08** if the following record (now in the input area) contained an end-of-file (EF) identifier indicating that the input tape had been completely processed. By testing the switch at the end of each cycle, it is determined whether there are no further records in the input area to be processed, and whether the input tape and output work tape are to be rewound and Pass II called in. When an input record in the input area contains an EF identifier, the switch is set, the last record lying in the work area is processed, and then the switch is tested, which leads to the termination of Pass I. If the end-pass switch is not set, the program jumps, via block **04.07** to block **01.07** to move the fields of the record that is then in the input area **102b**

to the work area **192d**, which starts another cycle of processing. Block **01.08** reads the next source record into the input area, and it is tested for an end-of-file identifier (block **01.09**), a valid F/C code (block **01.10**) and a C-code (block **02.01**), all in the manner described above. If the new input record has a C-code, its Comment field is appended (block **02.02**) to that of the record in the work area, and the loop repeated until a record with a F/C code that is not a "C" is in the input area. Thereafter, the record in the work area is processed in the manner described, starting with block **02.05** and as indicated by NOTE **02.04**.

Processing of Individual Codes—Pass-I

The flow chart of FIG. 8 has been simplified and omits a number of detailed tests and paths of coding that are sufficiently outlined hereinafter for an understanding of the invention. Many of the codes do not enter the EDLIN subroutine (block **03.05**) since the Comments fields of such codes are not displayed (e.g. J- and E-codes). In addition, the PAREN subroutine is different for D- and J-codes, which differences are noted below in the detailed discussion thereof.

As noted above, each code is processed in detail by an individual set of instructions to perform functions peculiar to that code; two functions that all codes require are: (1) determine the setting of Line Counter A (the number of lines required by that particular symbol on a flow chart), which is inserted in the output record as "# lines" (see Table I), and (2) adding that number to the running total of Line Counter B to get the total number of lines required for the current chain, which is entered in the Chain Table. For the purpose of processing the individual codes, separate legs or paths of coding are provided, with each leg being a branch from an individual comparison test of the current code against one of a set of constants that respectively represent the codes that are employed.

J-code

The test **04.02** initiates the PAREN subroutine **05.04**, which is used to extract the destination tag from the asterisk field of the Comments. Block **05.04** determines whether the Comments field contains an asterisk field (located at the beginning of the Comments field) and if it does not, the program branches to block **06.01**. The later controls the picking up of the destination tag from the operand field of the input record which currently sits in the work area; the operation code of the input record is examined and the appropriate operand address is picked up. If the Comments do contain an asterisk field, block **05.05** develops a subroutine for extracting the destination tag from that field. If there is a label associated with that tag, it is ignored in the processing of J-codes, since labels are handled only in connection with D-codes. From whichever source the destination tag is obtained, block **05.06** moves it to field LHT of the output memory area, and the program jumps to a test **04.03** for the J-code, which leads to a branch to the CHENT subroutine **06.04** to create the Chain Table entry. The editing operation of block **03.05** is bypassed for J-codes (the Comments field is ignored completely for the processing of J-codes except for the asterisk-field search). Block **03.06** sets Line Counter-A to 10 lines (which is suitable for the fixed format symbol plus an extra space left between the Jump symbol and the next chain in the column); block **04.01** similarly increments Line Counter-B, which provides the information needed for the Chain Table entry to be made at blocks **06.04** and **06.05**. The record is then complete and can be written to the output tape **110** via block **04.04**.

H-code

Line Counter A is set to 10 lines covering the fixed format of this symbol, and Line Counter B is incremented by the value of Line Counter A; the editing subroutine

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is bypassed. Thereafter, block **04.04** writes the record to the output tape. The H-code may be considered to be the end of a chain, if desired; and the Chain Table subroutine CHENT would be entered accordingly. However, preferably it is not so considered, since the computer operator may thereafter push a start button and the program would pick up the next line in the flow of coding. Accordingly, by not treating the H-code as the end of a chain, the next intended line of coding is naturally followed in the flow chart. The record is written to the output tape without the Comments field.

E-code

This code is processed in the same fashion as the H-code, except that it is considered the end of a chain and the Chain Table subroutine CHENT is entered and followed in the manner described above for J-codes.

P-code

Initially, the editing subroutine EDLIN represented by block **03.01** is entered to carry out the P-code editing. The coding supplies a constant for the length of line that is to be moved from the Comments field; for the P-code, this length is arbitrarily set for 19 characters as a maximum. An "intelligible" line is one that does not exceed the maximum and does not break up a word; hence, it may and usually will be somewhat smaller than the maximum length. EDLIN operates by finding the start of the Comments field which is supplied by the Comments Locator and then counting successive characters of the Comments field until that character is located which would mark the last character for the maximum length of the prescribed line. If this marked character is a space, then the intelligible line is exactly the length desired. If this character is not a space the subroutine steps back to the left until it does find a space, which signifies a word break; thus, the character immediately to the left of the space is the RHE of the line to be moved to the output Comments area. After the line is moved to the output area a control symbol is placed to its right as a line delimiter. The Output Area Locator is adjusted to the new RHE of the Comments field in the output area, and the Comments Locator for the work area is advanced to the first character of the next line to be moved out. A test is made before each line operation is performed to determine if there is an indicator for the end of the Comments field, which was placed there when the Comments were moved into the work area. If it is not, the logic recycles back to pick up a new line until eventually the end of the Comments area is found and the editing operation terminates. Line Counter A is stepped for each edited line of Comments that is moved to the output area; in addition, a count of 5 lines is added to allow for the fixed format of top and bottom symbol lines and a 3-line vertical connection to the next box. Line Counter B is similarly advanced.

A special case arises where the line length that is desired is too small to pick up even a single word; for example, where the word in the work area is 20 characters, and for a P-code the line length desired is only 19 characters. If this occurs, an artificial "word" is made by forcing a space into the twentieth character which results in a 19 character word and line, and permits the EDLIN loop to operate.

The P-code does not enter the PAREN subroutine; after the editing operation is complete the record is written to the output tape.

N-code

This code is handled in the same fashion as the P-code, except that each intelligible line is 15 characters long. Pass III arranges for the offset position of this symbol.

B-code

This code is handled without editing by EDLIN, and the Comments field is picked up in its entirety and moved

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to the corresponding output record area. The Line Counter A is set to zero, since the space for the chart title is allocated for each page, and it does not vary from page to page.

S-code

The PAREN subroutine **05.04** is entered from test **04.02** to place the destination tag in the output field LHT. If there is no asterisk field, the A-Address in the operand work area is used instead via block **06.01**. If asterisks were present, after the contents were moved via blocks **05.05** and **05.06**, the Comments Locator ends up pointing to the first character after the right-hand asterisk. If this character is a space, it is deleted by advancing the Locator to the right; and the Locator is so advanced until the first non-space character is found. A subroutine to perform this operation is also used for the D-code editing. Thereafter, the subroutine EDLIN (as described above for the P-code) is entered 5 separate times to move the Comments to the output area; the fixed odd-shaped format of the SUBROUTINE symbol (see, for example, block **13.02**) allows for 5 lines that are respectively 14, 15, 15, 15 and 14 characters long. After each call for a line, a test is made of an indicator (set by the subroutine itself) to see if all the Comments have been processed, and when the indicator is set the remaining calls are by-passed. No calls are made after the fifth one, so that any remaining Comments are dropped.

In processing S-codes, supplementary editing operations take place when there are only one or two lines to go in the symbol. Since the block is of a fixed column length and line processing is from top to bottom, Comments of one or two lines are moved down in order to center them within the box. Thus, if an end-of-Comments-area indicator is set after the first call on EDLIN (which indicates a single line), a subroutine arranges to shift this single line two positions to the right in the output Comments area, and each of the vacated spaces has an end-of-line symbol place therein. This anticipates the printing-layout operations of Pass III, which automatically moves the single line to be printed to the middle line of the box, so that it is centered in the block without additional coding. If the second EDLIN call produces an end-of-area indicator, the second line is shifted two lines to the right, the first line is shifted one line to the right, and end-of-line symbols are inserted between the two lines and in front of the first line, which have the effect in Pass III of moving the two desired lines into the second and fourth lines of the symbol to provide neat centering.

In the S-code processing, the Line Counter A is advanced 10 lines because of the fixed symbol format, Line Counter-B is correspondingly advanced and the record is written to the output tape.

D-code

The PAREN subroutine is entered via asterisk field test **05.04**. If there is an asterisk field, the subroutine is performed twice for the D-code, since two asterisk fields may be provided for the two lateral branches from a DECISION symbol. If only one asterisk field is present, the second entry to the subroutine has no effect. The destination tag extracted by the first call on PAREN is put in field LHT in the output area; and the destination tag extracted by the second call is placed in field RHT. If the second call produces no tag, the field RHT is cleared. If the first call on this subroutine produces no tag, then there is an input error, both LHT and RHT are cleared, and the DECISION symbol is printed out on the flow chart during Pass III with no lateral branches being indicated. The labels for the main flow branch and a single lateral branch from a DECISION symbol may be supplied by a special code that supplements the flow chart code, as noted above. Alternatively, these labels may be provided by

means of the asterisk field in accordance with the following format:

LABEL*TAG,LABEL*TAG,LABEL*

The first label is the main-flow branch, the second label corresponds to the associated tag placed in LHT, and the third label corresponds to the associated tag placed in RHT. These three labels are extracted via block 05.05 and stored via block 05.06 in the first 15 characters of the output Comments area in the order given. Labels of more than 5 characters are truncated by the subroutine, and if a label is missing its corresponding output field is cleared to spaces. If there is no asterisk field, the destination tags are picked up via block 06.01 from the A- and B-Addresses of the input operands, and the labels are picked up via block 06.02 from a stored table of contents as determined by the special label code.

A subroutine is used which shifts the Comments Locator to the right, if necessary, to bypass any spaces between the end of the asterisk field and the first non-space character of the actual comment, in a manner similar to that described above for the S-code operation. The Output Area Locator is advanced 15 characters so that the first actual Comments line is laid down after the labels. Thus, on a DECISION symbol, the actual start of the Comments in the output area is the sixteenth character of the field.

A test is then made to see if the total number of characters in the Comments work area is 13 or less; 13 corresponds to the room available along the middle line of the DECISION box, whose size is arbitrarily set to permit six Comment lines of 7, 11, 13, 11, 7 and 3 characters, respectively. If the total number of Comments characters is 13 or less, two end-of-line symbols are inserted into the output area (representing blanks for the first two lines), an EDLIN call of 13 characters is made and all remaining calls on this subroutine are bypassed. If the Comments work area contains more than 13 characters, then successive EDLIN calls of 7, 11, 13, 11, 7 and 3 characters, respectively, are made. The logic operation is similar to that described above for the S-code. After each call, the end-of-area indicator is tested, and if set, all remaining calls are bypassed. If EDLIN indicates an intelligible line is impossible, a word is forced, as explained above, by inserting a space in the last character that fits. Due to the earlier test for a total of 13 characters, it is not possible for the first call of 7 characters to produce an end-of-area setting; however, if the second call produces such a setting, both lines are shifted on to the right in the output area, and an end-of-line symbol is inserted in front of the first line to produce a more attractive line spacing, as explained above in connection with the S-code. If the third call of 13 characters produces an end-of-area setting, no shifting is performed since otherwise the 13-character line that has already been moved might not fit into the smaller available space in a line below it. If the end-of-area setting is not reached after the final, sixth call, the Comments work area is simply truncated; if a large Comment is desired by the programmer, a NOTE can be used for that purpose.

After the EDLIN subroutine is complete, Line Counter A is advanced 13 lines corresponding to the fixed format used for the DECISION symbol. Line Counter B is similarly advanced, and the output area is written onto the output tape.

W-code

This code is initially routed down the path followed for the D-code to process the asterisks field in the associated PAREN subroutine. Immediately thereafter, a test is made to determine if it is a W or D-code, and if the former, it is processed down its own branch leg of logic. The remainder of the Comments field of the W-code is ignored, and instead, a constant is moved into the output Comments area beginning with the sixteenth character, and the

code of the output record is changed to a D-code, so that it may be processed in that fashion from then on. The constant which is moved into the Comments area consists of the following: an end-of-line symbol, 11 minus signs, another end-of-line symbol, the word "SWITCH," another end-of-line symbol, 11 more minus signs, and two more end-of-line symbols to form the fixed format symbol shown as block 04.06. Since the code is now changed to D, Passes II and III treat it as a D-code and the constant in the Comments field is so arranged that it is printed out as the desired SWITCH symbol. The Line Counters A and B are handled in the same fashion as in the D-code described above, and the output area is written out to the work tape.

T-code

In order to avoid ambiguity in the flow chart, the TEXT of T-codes is printed out without a special symbol only at the start of a chain of flow, for otherwise the T format would interfere with the appearance of the chart. Thus, if the T-code occurs other than at the start of a chain, it is converted to an N-code and processed as such in the manner described above to be printed out in a NOTE symbol. This is done by an initial test to see if the code of the previous record is a J, E, or B; this previous record code was saved by moving it into the memory filled area field LSTCD prior to transferring the current record from input to work area. If the T-code does come in at the start of the chain, it is processed by making successive calls on EDLIN until the end-of-area indicator is set; the line requested is 30 characters long. Line Counter A is incremented for each such line, and after the last line, it is incremented by "3" to provide a space between the final line and the first part of the next F/C symbol to be printed in the column.

Summary of Pass-I

This first Pass examines successive records of the input program to be documented and extracts all the information needed to produce an F/C symbol from each record that contains an F/C code. Where the Comments field for an F/C symbol is greater than that carried by one record, the Comments fields of succeeding records (denoted by C-codes) are tacked to the previous record. A sequence number counter is incremented for each record and used to identify that record in the subsequent processing. Any tag that identifies the entry point of each input record is extracted and used to construct the skeleton of the Tag Table. The destination tags carried by records for the JUMP, SUBROUTINE, DECISION and SWITCH records are extracted as well as branching labels for the latter two symbols. Editing operations of the Comments fields of P-, N-, S-, D-, and T-codes are performed, and the number of lines along a column is determined for each symbol. In some cases, the symbols are a fixed format, and in other cases they are variable in format and their column length is determined by the number of Comment lines that have to be printed out.

Each unconditional transfer, i.e. J- and E-code, defines the end of a "chain" of coding, and a Chain Table is constructed that has an entry for each chain; each chain is formed as a sequence of symbols terminating with a JUMP or EXIT symbol and represents a section of program logic that is referenced on the flow chart to one or more other sections and that can be treated as a separable entity of logic for display on the flow chart. The Chain Table entries are identifiable by the sequence number of the last symbol of each chain and include the column length required for recording each chain on a page of the flow chart. The flow chart codes for HALT, SUBROUTINE, DECISION and SWITCH are not treated as chain terminating symbols, but rather as parts of a chain. The HALT symbol in some respects is like an unconditional transfer symbol in terminating a section of logic flow; however, it may be followed, as the program is performed, by the operator of the computer restarting the program,

which would lead to the next record of the original program sequence being the one to be processed. Consequently, the logic flow from the HALT symbol, in effect, is an entry to the next symbol in the original sequence, and these symbols are preferably considered as part of the same section of logic and not separated on the flow chart. The SUBROUTINE symbol refers to a sub-section of logic which continues the main flow processing and does not terminate it, though the details of it are ordinarily separately reviewed and are therefore best left to a separate SWITCH symbols, each also have a branch that continues the main flow processing and are therefore incorporated as parts of the chain and not as branches from it.

PASS II

The flow chart for Pass II is shown in summary in FIGS. 8/7 to 8/18. As indicated in the text at the beginning of FIG. 8/7, the function of this pass is to complete the construction of the Tag Table and to assign all flow chart symbols to certain positions on the F/C pages. The input records for Pass II are supplied by the Pass I output work tape 110 and the output of Pass II is written on the second work tape 111.

The sequential operation of Pass II is determined overall by the sequence of chains in the Chain Table, and certain chains are processed out of that sequence. Within any chain, successive records are generally processed in sequence as received, except when DECISION records are encountered; at that point the processing of the chain containing that DECISION is interrupted and the branch chains are investigated and (in the present embodiment) processed. As indicated in block 07.01, a Chain Table Locator in the primary storage area is employed, which always points to (carries the address of) the Chain Table entry which is currently being processed along a main flow column. As the processing of each chain is completed, this Locator is advanced to the next chain in sequence.

Other initial operations performed in Pass II, block 07.02, include that of setting a page-number (Page #) counter and a column (Col #) counter to "1," a box-number (Box #) counter to "0," and a Line Counter 1 (LNC-1) to "1."

The main loop of Pass-II then begins at block 07.03 with the first input record from work tape 110 being read into the memory area 102b (an additional memory work area is not required in this pass). A test is made, block 07.04, for an end-of-file indicator, and if it is found, the program branches to 09.06 which controls the re-winding of the tapes 110 and 111 and the initiation of Pass-III. If it is not the end of the file, the sequence number of the current record is stored in field TPOS for later use; the program steps to a test 07.05 for a B-code, and if it is found, the program branches to block 12.09 where the end of a page is forced by setting Col # to "4" (assuming a 4-column page). Then the program transfers to an end-of-column subroutine ENDCOL, where a pseudo-connector record is developed. Since a B-code involves the development of a new chart title, a new page is normally started, which is the function of the ENDCOL subroutine. However, as explained below, in certain situations (as where it is the first B record of the chart) is not necessary to start a new page and thereby needlessly skip a page; and the ENDCOL subroutine is essentially bypassed to block 08.06, which writes the first B-code record to the output file. The program is then recycled via the test 08.07 back to block 07.03 to read the next input record from work tape 110.

The next record which is not a B-code is processed via block 07.06 where the current value of LNC-1 (which is "1," for the first record of a page) is established as the Line # of the record being processed by moving it to LNC-3. As indicated in NOTE 07.07, the Line # of a record is allocated as the Line # of the corresponding

symbol to be printed on the F/C page and thereby fixes the position of the symbol in a column. Thus, in the example of the first record after the chart title, which may be text or some symbol, the Line # is set at 1. Block 07.08 extracts "# Lines" (i.e., the length) of the current record therefrom (see Table 1) and adds it to LNC-1. Thus, the previous LNC-1 represents the ending line-number-plus-one of the previous symbol, and the new LNC-1 becomes the starting line number of the following symbol; the current LNC-3 is the beginning line number of the current symbol and is used as a temporary store of that number before it is moved to the output record.

The program steps to block 08.01 to compare the new value of LNC-1 to an end-of-page constant EPCON, where EPCON is the total number of lines allowed in a column, that is the address of the last line (which is set to 106 for the high-speed printer page, and would vary for different types of recorder and display devices). If the end of a column has been reached, the program branches to block 11.01 for the ENDCOL subroutine; but if the new LNC-1 is less than EPCON, there is room in the current column for the current symbol and the program steps to block 08.02. Box # is incremented by 1, and the Page #, Box # and Col # are moved to the output record in the output memory area.

Thereafter, test 08.03 determines if the input tag field of the current record contains a tag or is blank. If it contains a tag, the program branches to a block 09.08 to locate a Tag Table entry for that tag. The input parameter for this operation is a locator which points to the left-hand character of the tag field for which a search is desired. If this field contains a tag, the Tag Table is then searched by a straight series of compares beginning at the start of the table and running down until the tag entry is found. The starting address TBIN of the Tag Table was set during Pass I, and a counter NUMTAG was established during Pass I containing the number of entries in the Tag Table. Thus, NUMTAG tells the routine when it has exhausted the table as it makes its series of compares; upon exhaustion of the table if the Tag Table does not contain the desired tag an indicator is set to show this. The main output of this operation is the setting of a locator to the address of the left-hand character of the proper entry in the Tag Table; a subsidiary output is the setting of indicators to reflect "tag found," "no tag in field," or "Tag-Table entry missing." Block 09.09 places the page location data (Page #, Box #, Line # and Col #) of the current record in the Tag Table entry for the associated tag; this completes the basic structure of the Tag Table entry, and the supplementary data of cross-references to "from connectors" is subsequently added as explained below.

Modified logic is employed for T-code records; that is, Box # is not incremented and the Tag Table search is bypassed, because T-codes do not produce a symbol on the F/C page and, thus, do not carry box numbers or tags.

The Tag Table receives any "from connector" information in the current record; and this operation starts with a transfer back to a test 08.04 for a D- or J-code. If it is found, the program branches to block 10.01 where a cross-reference subroutine is performed in order to place Box # of the current record as a cross-reference in the Tag Table to any destination tags contained by the current record. Thus, block 10.01 obtains the destination tags (LHT and RHT) from the current record and looks those tags up in the Tag Table (in a manner similar to the above described subroutine of block 09.08). A test 10.02 determines the results of the search; if no tags are found in the table the program returns to the next operation 08.05 of the main flow; but if the tags are found a test 10.03 determines whether this is the first reference to the tag. If it is, block 10.06 puts Page # and Box # of the current record in the cross-reference field of that Tag Table entry, and the program returns to the next main flow block 08.05. If it is

not the first reference, the program branches to block 10.04 which sets a signal in the Tag Table entry to indicate that there is more than one such cross-reference (which signal determines that an asterisk is to be printed at that entry point; see for example, the branch input to block 07.03); the program returns to the next main path block 08.05, to which it would pass if test 08.04 had proved negative. The cross-reference subroutine of 10.01 is entered twice for D-codes, since two such tags may be carried by such a record.

Upon return to the main logic path of the program the test 08.05 determines if the current record is a D-code. If so, the logic branches to block 13.01 for processing the DECISION record; if it is not a D-code, block 08.06 writes the record in the output memory area to the output tape 111.

Thereafter, block 08.07 tests to determine if there is a J- or E-code in the current record, which codes indicate the end of a chain. If it is not the end of a chain, the program recycles back to block 07.03 to read in the next input record and repeat the processing loop described above. If it is the end of a chain, as indicated by NOTE 08.08, the program proceeds to locate the next chain to be processed via test 09.01, which determines whether the current record is being allocated to a main column or to a branch or secondary column. If the latter, the program branches to block 15.07, which is described below; if the former, block 09.02 advances the Chain Table Locator to the next table entry, and test 09.03 determines from the indicator of that entry whether this chain has yet been processed. If so, the program loop continues until the test 09.03 finds an unprocessed chain. This loop is an important part of the processing system since by the very nature of the processing of DECISION branches, as explained below, it is possible that the next chain in sequence may have already been processed and assigned to a secondary column by the DECISION branch logic.

When the next unprocessed chain is found, its sequence number is used to locate the corresponding data record on the input work tape 110 via block 09.04. It should be noted that the Chain Table entry for any chain contains the final sequence number of that chain. Thus, if the Chain Table Locator points to the chain that it is desired to process, the sequence number of the first record of that chain corresponds to the final sequence number of the previous chain plus 1. The tape position sequence field TPOS contains the sequence number of the last record that has been read; therefore by subtracting TPOS from the Chain Table entry sequence number, the difference corresponds to the number of input records that have to be skipped to get the desired record. Ordinarily, the next chain to be processed is the next physical chain on the tape; which is indicated if the sequence number of the Chain Table entry is the same as TPOS, and no further records have to be skipped. Thus, block 09.04 computes the difference between TPOS and the sequence number of the first record of the desired chain; it then proceeds to skip that number of records on the input tape 110 so as to position the tape at the first record of the desired chain. After the input tape 110 is so positioned a test is made to determine if there are 20 or more lines left in the current column being allocated. This determination is made by subtracting LNC-1 from EPCON (the total column length) and comparing the result with the constant 20. This test is of assistance in insuring good page format in that a new chain is not initiated near the bottom of the page instead the remainder of the column is left blank and the new chain is started at the top of the next column. If this test shows that there are less than 20 lines left in the column, then an end-of-a-column subroutine similar to ENDCOL is entered to determine which column is currently being processed and thereby begin a new column, before reading the first record of the new chain. Whichever direction is taken by the last mentioned test, the program recycles back to the start of the main pro-

cessing path at 07.03 to read in the first record of the new chain and perform its processing.

ENDCOL Subroutine

If the aforementioned test 08.01 indicates that the end of a column has been reached, the program branches to perform a further test 11.01 to determine whether the current record is a J- or E-code. If it is such a code it can nevertheless be allocated to the current column since EPCON has its value chosen so that there is enough room at the end of every column to contain an additional 10 lines required for a J- or E-connector (or for a pseudo-connector). Consequently, if the test shows a J- or E-code, the logic is routed directly back to the main processing path at block 08.02 as if the test 08.01 against EPCON had gone the other way, since the rest of the processing beginning with block 08.02 can be properly performed on the current J- or E-record. The next record will then be directed by the test 08.01 into the ENDCOL subroutine to start a new column.

When the code is not a J- or E-, then the ENDCOL subroutine is entered at block 11.03 to arrange for the proper termination of the current column by the generation of a pseudo-connector record in the output file, and by the proper initiation of the new column which involves resetting the various indicators and locators. A pseudo-connector record is a short record, 13 characters in length, which goes to the output file and is used by Pass III to create a connector symbol at the bottom of each column that is not terminated by a JUMP or EXIT symbol.

TABLE IV.—PSEUDO-CONNECTOR RECORD

Field	CHAR	#	Field	CHAR	#
Control symbol.....	1	1	Line #.....	3	3
Letter "X".....	1	1	Page #.....	2	2
Spaces.....	2	2	Box #.....	2	2
Column #.....	1	1	Control symbol.....	1	1

The letter "X" identifies the record as a pseudo-connector. Column # and Line # (LNC-3) fix the location of the symbol on the page. Box # and Page # are those of the currently-processed record and are printed inside the connector symbol to indicate the next symbol in the path of flow. If it is a connector for the bottom of a main column (other than the last) of a page, the connector symbol contains the current Page # and current Box # plus 1; if it is a connector for the bottom of the last main column on a page, it contains the current Page # plus 1 and a Box # of "01."

A pseudo-connector is not needed when the previous record was a J-, H- or E-code record terminating a chain, since the J-, H- or E-symbol satisfactorily terminates the column. A test 11.02 for this condition is made by examining the contents of LFTCD, where the previous record's F/C code was saved. If the test is negative, block 11.03 proceeds to generate the desired pseudo-connector record in the fashion explained above and write it to the output tape. The program continues with block 11.04; if the test 11.02 indicates that the previous record was a J-, H- or E-code, this pseudo-connector operation 11.03 is bypassed and the program proceeds directly with the test 11.04. The latter tests the Col # counter to determine whether it is set to a value of "4," if so, the current column is the last column on a page, and the program branches to block 12.03, where the Page # is increased by 1; then to block 12.04, where the Col # counter is set to "1" and the Box # to "0." Then block 12.05 writes an end-page record to the output tape 111, which record comprises a control symbol that Pass III uses to determine when it has read in all of the records needed to create a page. With the end-page symbol a complete page of records has been written to tape 111, and the program then jumps to block 11.07 where the line counters are set to "1." Thereafter, a test 12.01 determines whether the currently processed record is a B-code, and if it is the program branches to

18.06 where the record is written to the output file, and the program proceeds to process the next record via blocks 18.07 and 07.03. If it is not a B-code, the program jumps to 07.06. Jump 12.02 transfers the program to block 07.06 to repeat the initial processing of the line number of the current input record in view of the resetting of the line counters at the start of this new column. Thereafter, the program proceeds in the manner described above.

If the test 11.04 indicates that the current column being allocated is not the last one on the page, the next test 11.05 determines if any symbols have been allocated to the next adjacent column, which condition can occur upon processing of branch chains from DECISION symbols, as explained below. Thus, if the adjacent column has been already allocated, the program branches to block 12.07, where "2" is added to the Col # counter. This has the effect of skipping the adjacent column to obtain the next column thereafter for the current allocation. If this new column number is greater than "4," the program operation is via blocks 12.03 to 12.05 to start a new page as described above. In any case, the program transfers back to block 11.07 to reset the line counters and start the processing at the top of a new column.

If the test 11.05 determines that the adjacent column has not yet been allocated, block 11.06 adds "1" to the Col # counter. Block 11.07 resets the line counters, and the processing continues at the top of a new column, in the manner described above.

B-code Processing

As previously described a test 07.05 for the B-code is made shortly after each record is read. When such a code is found, a new page is started by setting the Col # field to "4" and then entering subroutine ENDCOL at block 12.09. This subroutine, via test 11.02, blocks 11.04, 12.03 to 12.05, 12.01 and 08.06, sets up a new page and returns control directly to the main path at the point of writing the record to the output.

There are two special cases where the ENDCOL subroutine is preferably not entered. The first is where this is the first record on the input file, which is normally a B record. To avoid skipping a blank page, a test is made of the sequence number of the B record, and ENDCOL is bypassed if it is "1." The second case is where B-symbol by chance is the start of a new page; that is, where the current symbol location is at the start of a new page, and the previous page was already properly terminated. This condition can be tested for by testing Box # for "00"; which number indicates that the current operation is at the top of a new page and that ENDCOL can be bypassed. Normally, a B-code should be preceded by a J- or E-code, which would properly terminate the previous page. If by error that should not occur the B-code will nevertheless start a new page, and a pseudo-connector for the previous page is generated in block 11.03.

DECISION Branch Processing

The aforementioned test 08.05 for a D-code, if affirmative, directs an immediate branch to the associated processing logic at block 12.01, and there-preciding is TEXT that sets forth the function of this processing. The DECISION record and associated data, are analyzed to determine how best to illustrate its branches on the flow chart. The processing includes the secondary-column subroutine SCOL, block 16.07, which, as the TEXT there-preciding indicates, analyzes a chain of code branching from a DECISION symbol in order to determine if it is possible to assign that chain to a secondary-column. When it is determined that such an assignment to the secondary-column should be made, the DECISION-branch logic functions (in this embodiment) as a control routine for processing directly to process that chain in the secondary-column.

There are two key indicators in the input records which must be set for all DECISION records: These are in-

dicators LHTC and RHTC, which are used by Pass-III to determine how the page is to be laid out. LHTC refers to the destination tag located in field LHT, and RHTC performs a similar function for field RHT. Each of these indicators can have three possible values that correspond to the following courses of action, respectively: (a) This branch requires a connector symbol; (b) The coding for this branch is in the adjacent secondary-column; (c) The adjacent column is not being used for this branch; it is not known whether a connector is to be used or a line can be drawn to show the path of flow.

Of the above three possibilities, courses-a and b are definite; course-c indicates that the final result is unknown and is to be finally resolved by Pass III. Pass I, when it sets up its output record, sets both LHTC and RHTC to indicate course-c. Pass II may change them to course-a or b, or leave them set at course-c. Under certain conditions, subroutine SCOL will set the two indicators; under other conditions the DECISION-branch logic itself does the setting.

It should also be noted that subroutine SCOL determines the suitability of a chain for display in the adjacent secondary-column, and the final determination of so displaying that chain is made by the DECISION-branch logic itself. For a single branch DECISION, field LHT is used. For a two-branch DECISION, both LHT and RHT are used. The subroutine SCOL generates a "no good" signal if the testing field has no tag (thus, on a one-branch DECISION, the field RHT has no tag, and SCOL generates "no good").

The DECISION-branch logic is divided essentially into two sections. The first section starts at block 13.01 and determines (by means of SCOL) if either branch chain is suitable. The second second section (beginning at block 14.03) is a control routine which sets up a branch chain to be processed and starts the processing at block 07.03; after the branch chain is processed, the program control picks up processing of the main column of flow again from the point it was temporarily halted to handle the branch chain. In the event that no chain is allocated to the secondary-column, the second section of the DECISION-branch logic is not entered.

The first test 13.01 determines whether the current processing is taking place in a secondary column by checking an indicator EVOD which assumes one value for processing main columns of flow and another for secondary or branch columns. If EVOD indicates "secondary column," no further branching to subordinate columns takes place from the secondary-column (in this embodiment); LHTC and RHTC remain unchanged (Pass III determines whether to draw a connecting line or a connector symbol at the branch point), and the program branches back to block 08.06 to continue with the processing of the branch chain in that column. If EVOD indicates the current processing is in the main column, the program steps to the SCOL subroutine 13.02, and the latter symbol indicates that SCOL starts at block 16.07.

SCOL subroutine

The SCOL logic begins, block 16.07, with setting the exit from the subroutine back to the main-flow reentry point. This subroutine is entered from block 13.02, 13.04 or 15.01, and the reentry point in each case is the main-flow block immediately thereafter; namely, block 13.03, 13.05 or 15.02, respectively. The input parameter for this subroutine consists of a locator pointing to the left-hand character of the tag field in question (LHT for the first-branch analysis and RHT for the second). This locator also fixes the location of LHTC or RHTC, as the case may be, since the latter indicators are located one character to the left of their associated tag fields. The output of the subroutine is a signal stating whether or not the chain involved can be put in the adjacent secondary-column, namely "ok" or "no good," respectively.

Block 16.08 looks up the destination tag of the first

branch LHT for an entry in the Tag Table (in the manner similar to the operation of block 09.08 described above). Test 16.09 checks if the record's tag field contained a tag, and if the tag could be found in the Tag Table; if either condition is negative, block 17.02 sets the indicator LHTC for a "connector" symbol, and the subroutine jumps to EXIT 18.08.

If a tag and its corresponding Tag-Table entry are located, SCOL determines whether the chain identified by the tag is suitable for assignment to the adjacent secondary-column. All six of the following criteria must be satisfied or a "no good" signal is established, and LHTC is set to "connector" by block 17.02.

(a) The record identified by the destination tag must not have been assigned. Test 17.01 checks this criterion by determining whether the Box # field in the current Tag Table entry is a blank; if so then the associated record has not yet been processed and allocated (for block 09.09 makes the Tag Table entry upon such allocation). If the Box # field is completed, then the destination has been allocated a place on the chart (this of course is the case for all DECISIONS which jump to an earlier section of the coding).

(b) The destination record must not be in the current chain. To check this criterion, block 17.04 extracts from the Tag Table (Table II) the Sequence # of the destination tag and places it in field TSQ. The Chain Table Locator points to the Chain Table entry for the current main-column chain and thereby the final Sequence # of the current chain can be extracted and compared with the field TSQ, the Sequence # of the destination tag. The TSQ must be greater than the current-chain Sequence # or else the destination tag is in the current chain (since it has already been determined that it is not in a prior chain by test 17.01). Thus, the result of this comparison supplies the answer to the test 17.05 of this criterion.

(c) The destination record must be at the start of a chain, rather than at some other point thereof (for this embodiment of the invention). This operation is performed by locating (block 17.06) the chain that contains TSQ; beginning at the start of the Chain Table (Table III), the successive Sequence # fields thereof are compared with TSQ until the Chain Table entry is higher than TSQ. The chain of the tag is thereby located. To obtain the start of that chain, block 18.01 extracts the Sequence # of the previous Chain Table entry (which is that of the last record of the previous chain), and "1" is added to it to obtain the Sequence # of the first record of the chain containing the tag; the latter result is compared with TSQ (in block 18.2) to determine if they are equal. If so, the destination record is the start of a chain. Block 18.01 also picks up the length of the chain containing the tag for subsequent use in this subroutine. An extra test has been found desirable to determine if the destination record is the second one of the chain; and, if so, whether the first record is TEXT or NOTE. If it is, the second record is then considered as satisfying this criterion; and the chain can be printed in the secondary-column with NOTE or TEXT at the beginning thereof.

(d) The adjacent secondary-column must be free at this branch point of the DECISION symbol (for this embodiment). This criterion determines that the adjacent column at the branch point has not been previously assigned to a chain coming down from another DECISION symbol in the same main column, but located above the current DECISION symbol. Block 18.03 determines this condition by comparing LNC-3 against LNC-2. LNC-3 is the column Line # at which the current DECISION symbol begins; and LNC-2 is a field completed by the DECISION-branch logic when the secondary-column is allocated and it represents the Line # at which the last chain in the secondary-column ends. If LNC-2 is greater than LNC-3, the adjacent column is occupied at the branch point and therefore not free.

(e) There must be enough room in the adjacent column to contain the entire chain (for this embodiment). The test 18.04 for this criterion is performed by taking LNC-1 (the ending line for the current DECISION symbol), subtracting "6" to get the Line # of the branch point of the DECISION symbol (which, by an arbitrary rule, is where branch chains should start on the flow chart), adding the number of lines in the destination chain (obtained by block 18.01) and comparing the results with the constant EPCON (the column line-length). If the latter is smaller, there is insufficient room in the secondary-column for the entire chain.

(f) The number of records to be skipped on the work tape 110 to reach the desired chain must not be excessive. This criterion is checked via block 18.05 by taking TSQ, subtracting "1," and then subtracting the field TPOS (the Sequence # of the record that was last read). The result gives the number of records to be skipped to reach the chain in question, and if test 18.06 determines that this number is excessive then it would be too time-consuming to pick up the chain. This number would vary depending upon the apparatus and individual choice as to efficiency. For example, a skip of 40 records or more has been considered excessive for some purposes.

If all of the above criteria (a) to (f) are satisfied, the "ok" indicator is set by block 18.07. If the chain is "ok," no further action is taken by SCOL. The subroutine exits via block 18.08 and returns to the main flow reentry point of the DECISION-branch logic (e.g. to block 13.03 after analysis of the first branch chain). Failure of any of the above six criteria results in the "no good" signal being set by block 17.02.

If a chain is found to be "ok" then the number of records to be bypassed (as computed by block 18.05) to reach the desired chain is preserved in an appropriate memory field, since it will be used by the succeeding DECISION-branch logic.

DECISION-branch logic continued

The "ok" and "no good" signals from SCOL are used in the test 13.03; and if "ok" the program steps to subroutine 13.04, which directs another entry into the SCOL subroutine for the second branch tag RHT. Upon completion of the second tag analysis by SCOL, the subroutine exits back to 13.05 to test if the second tag can be assigned to the secondary-column. The different possible combinations are handled as follows: If the first branch could not be assigned, as tested at block 13.03, the program branches to 15.01 for the second-branch operation of SCOL. If that second branch likewise could not be assigned (test 15.02), then the DECISION-branch logic exits by pumping back to the main flow at block 08.06, the current DECISION record is written to the output tape 111, and main-column processing continues. However, if the second branch is tested to be "ok" in block 15.02, the program branches to block 15.05 to set the status indicator for the second branch to the secondary-column. The program then jumps to block 14.02, which writes the current DECISION record to the output file and steps to block 14.03 for processing the second-branch chain. However, if the first branch tests "ok" in block 13.03, while the second branch tests "no good" in block 13.05, the program jumps to block 14.01 to set the status indicator for the first branch to the secondary-column. Thereafter, block 14.02 writes the current DECISION to the output file, and block 14.03 initiates the processing of the first branch.

Where both the first and second branches test "ok," a decision is made (in this embodiment) to process one of the branches and mark the other one for a "connector"; a criterion found to be suitable is that of determining which branch is the closer one, as shown by test 13.06. The relative closeness is readily determined from the number of records that have to be skipped to reach each chain, as computed in block 18.05 of the SCOL subroutine. If

the first branch is closer, the status indicator for the second branch is set to "connector" in block 13.07 and the status indicator for the first branch to the secondary-column by block 14.01. If the second branch is the closer one, the program branches to block 15.04 which sets the status indicator for the first branch to "connector," and block 15.05 sets that of the second branch to the secondary-column, and the program jumps to 14.02 for writing the current DECISION to output. Appropriate coding is provided to preserve the information derived during the SCOL analyses for LHT, so that it is not lost when SCOL is repeated for RHT.

When a chain is finally chosen for the adjacent secondary-column, its corresponding indicator LHTC or RHTC is set to indicate this; this indicator for the other chain is said to use a "connector." This is done even though there may be only one branch in the DECISION, since the setting of RHTC for a nonexistent tag in RHT is ignored by Pass III.

Secondary-column processing

When block 14.02 writes the current DECISION to the output file, the current value of TPOS (the Sequence # of that DECISION record) is stored for later reference and reentry to the main column. Block 14.03 starts the second section of the DECISION-branch logic and uses the calculation (block 18.05) of the number of records to be skipped on the input tape to position the tape in front of the first record of the branch chain which is to be assigned to the secondary column. An indicator is then set (14.04) in the Chain-Table entry for that branch chain that the latter has been processed; this prevents further processing of that chain later on in the operation of Pass-II when it would normally be picked up in turn.

The column number is advanced (14.05) by "1," so that it points to the adjacent column next in order. The indicator EVOD is set (14.06) to indicate that a secondary-column is being filled; this indicator is needed, because the processing of the secondary-column chain is via the main processing logic; and it is tested upon completion thereof for return to the DECISION-branch logic.

LNC-1 (which indicates the bottom line number of the DECISION symbol from which the branch occurs) is stored (14.07) so that it may be subsequently picked up upon return from the secondary-column processing. The new LNC-1 for the secondary chain is obtained (14.08) by subtracting "6" from the previous LNC-1, which has the effect of starting that secondary chain 6 lines above the bottom of the DECISION symbol from which it branches, which is at the branch point of the diamond-shaped symbol. Control is then transferred (14.10) to the main processing path at the point 07.03 where it starts processing a new record, the first in the branch chain. As indicated by NOTE 14.09 the main processing path allocates assignments for the secondary-column in its normal fashion, since LNC-1 and Col # have been appropriately set. DECISION symbols that are encountered in the secondary chain tend to move the logic into the DECISION-branch coding via the test at 08.05, but the program is immediately returned to the main processing path by the test at 13.01 which initiates the branch logic. Eventually, a J- or E-code is found by the main processing path via the test 08.07; the next test 09.01 finds the EVOD indicator set to "secondary," and the program branches to block 15.07, which is effective to rewind the input tape to the record following the DECISION symbol from which the secondary-column chain branched. Since the SCOL subroutine measured the secondary chain and found that it would fit in the adjacent column, when the first J- or E-code is reached, the secondary-column chain terminates and it is proper to return to the main-column processing as indicated in NOTE 16.01. Block 16.02 transfers the then current value of LNC-1 to LNC-2, so that the latter represents the last line assigned to the secondary-column. LNC-2 may be required by SCOL in the event

that another DECISION occurs further down in the main column. The value of LNC-1 of the main column DECISION symbol, which was stored (at block 14.07) upon entry into the secondary-column processing, is returned to LNC-1 so that the main column processing begins where it left off prior to processing the branch. Block 16.03 resets EVOD to "main," and block 16.04 reduces the column number to restore it to its original main-column value. The indicators are all restored and the input tape is then at its proper place to continue the processing (NOTE 16.05) of the main-column chain that was interrupted for the secondary-column branch. The main-column DECISION record that initiated the branch operation has been completely processed so that the input tape is positioned to the succeeding record by block 15.07. For this purpose the TPOS of the DECISION record, which was stored is now subtracted from the current value of TPOS, which is the sequence number of the last record of the secondary-column chain. The result is the number of records that the input tape 110 must be backspaced for proper repositioning. After repositioning of the input tape, control goes back to 07.03 to read the next input record of the main column and continue the processing of the chain that was interrupted.

The foregoing operations of Pass II are performed iteratively on all of the records of each chain in the manner described. Successive chains are processed in order until a DECISION record calls for a branch chain from the main flow. At that time, the branch chain is analyzed to determine if it is suitable for allocation adjacent to the main flow chain, and if so, it is processed. After the last record of the input file is processed, the end-file indicator is detected (07.04) the output tape 111 is rewound (09.06) and Pass III is called in for operation.

Summary of Pass II

The primary function of Pass II is the allocation of flow chart locations to the symbols. The input is the output from Pass I. As each record is read, it is assigned a Box # and a Line # in the current column (the associated portions of the Tag Table are thereby completed). When the column is filled, a new column is started, until the rightmost column is reached. After this column is filled, a new page is started. DECISION records involve considerably more complicated processing than other types of records. Whenever possible, the F/C program attempts to place the coding that branches off from a DECISION in the column immediately to the right of the DECISION symbol, which is termed an adjacent or secondary column. However, before this can be done, a number of conditions must be satisfied by the branch chain, including the following:

(1) The section of coding that branches from the DECISION must be further down the input tape; otherwise it would have been allocated at an earlier point.

(2) The entire chain must fit in the adjacent column without overflowing the bottom of the column, to avoid breaking up a chain. (In other embodiments, branch chains that fit in larger page sections than a column may be used.)

When all of the above conditions have been satisfied, a chain branching from a DECISION symbol is considered eligible for assignment to an adjacent column. All of the information necessary to test the above conditions is contained in the Tag and Chain Tables and the record itself. In the event of a three-way DECISION, where there are two branches to be considered, it is possible that both branches will be eligible for the adjacent column; in this case, the chain closer to the current record is chosen. Each of the two destination tags in the record has an associated indicator which is set by Pass II to one of the three possible values—"adjacent column," "connector," or "unknown." If a chain is selected for the adjacent column, its corresponding indicator is so marked in the DECISION record. Chains not eligible for the adjacent col-

umn may be designated either "unknown" or "connector," depending on which of several conditions was not met. With a destination indicator of "unknown," Pass III attempts to draw a connecting line to show the path of flow instead of using a connector.

After a branch chain has been selected for assignment to the adjacent column, the input tape is advanced to the first record of this chain. An indicator is then set in the Chain-Table entry for this chain to show it has been processed; this is necessary to avoid reprocessing this same chain at a later time in a main column. The column number is advanced by "1," appropriate indicators and counters are set, and the program logic is then routed back to the same coding that processes records for the main columns; thereby common coding is used for processing chains in both main and adjacent columns. The first J- or E-code encountered while in the adjacent-column mode indicates the end of the branch chain; at this point the input tape is rewound back to the original DECISION record, and main-column processing continues where it left off.

Pass II also makes the necessary "from" connector entries in the Tag-Table to allow handling of cross-references by Pass III. Only the first such reference to any tag is noted, along with a signal if there is more than one.

PASS III

The function of Pass III is to form an entire flow chart page in memory, draw the necessary connecting lines from DECISION symbols, and produce a finished flow chart. The input consists of the output from Pass II on tape 111 and the Tag-Table located in memory. The output either goes to an on-line printer or to a tape for off-line printing. Pass III also produces the Table of Contents at the beginning of the flow chart.

Page layout

Throughout Pass III, a section (e.g. 15,000 locations) of memory is reserved (in this embodiment) for holding an entire F/C page internally; for computers having limited memory capacity, storage tapes or drum may be used to supplement the memory. This page-layout memory area is structured as contiguous "lines" of 120 characters each. The first line of each chart is represented by the first 120 characters of this area, the second line by the next 120 characters, etc. The location of any symbol on a page is given by its Column # and Line #. A subroutine 19.06 is used to convert these two factors into a memory address that represents the centerpoint of the first line of that symbol; it operates by multiplying the Line # by 120 and adding to the result one of four factors depending on the Column #. An index register is reserved for use as a locator; this index register always contains this base location on the page that the program is currently concerned with, and hereinafter it is referred to as the "Page Locator."

To "move" this Locator around on the page, "120" is added to the Page Locator, which moves it to the same position on the next lower line; subtracting 120 moves it to the same position on the line above. Moving to the right or left on the same line is accomplished by adding or subtracting the appropriate number of positions from the Page Locator. Since the Page Locator is an index register, indexing techniques may be used in place of actually modifying the Locator.

Every symbol has fixed dimensions for which constants are stored; that is, the horizontal dimensions are fixed for all symbols, and the vertical dimensions are fixed for some and variable for others (e.g. P- and N-codes). For all symbols, the location of the Box #, tag, and other related information is always located in a fixed position relative to the center-line of the column. Detailed specifications of these dimensions will be apparent from the flow chart of FIG. 8, which illustrates suitable values and conventions that are followed in printing the flow chart.

GENERAL DESCRIPTION OF PROCESSING

The first function performed by Pass III is to print out the Table of Contents (block 19.01). All information present in this table is obtained from the Tag Table, which is in memory throughout the entire program. Printing of the Table of Contents involves moving the necessary information from the table to the page-layout memory area; that is, for each chart a list of the tags and the page and box numbers therefor. A tally NUMTAG set by Pass I determines when processing of the Tag Table to Table of Contents is completed. An entire page is constructed prior to writing anything out and as many pages as needed to contain all the tags are produced. It has been found suitable to list the information in two columns or sections on a page, with the entire left column of the output memory area being filled before any entries are made in the right column thereof.

After the Table of Contents is completed, production of the flow chart pages begins. A full page is processed at a time; nothing is printed until the entire page has been formed in the page layout area, at which time the entire page is printed. A control symbol on the input tape 111 designates the end of a page; this control symbol is developed in Pass II. Each input record is processed separately, and a new record is not read until all processing for the previous record is completed.

Immediately upon reading a record by block 19.02, tests 19.03, 19.04 and 19.05 are made respectively for end-file and end-page indicators and B-code, as discussed below. Thereafter in the program, a subroutine 19.06 computes the chart location of the F/C symbol for the current record. This chart location is stored in the Page Locator and is obtained from the Col # and Line # of the symbol in the input record. After determining the location of the symbol, a check 20.01 is made for a cross-reference and the proper "from connector" is generated (block 20.02) if required. The record is then routed down a particular path, depending upon its F/C code. There is a separate page for each code. Each path performs the necessary layout for the particular symbol involved, puts in any connecting lines, box numbers, and tags necessary, and upon completion returns back to read and process the next record.

Any Comments text associated with a symbol is moved from the input record to the page layout area by means of a Move-Line subroutine 20.08. In an input record, each Comments line is delimited by a control symbol, and a second control symbol is used to indicated the last Comments line within the record. Each individual code path calls upon this subroutine when necessary in order to move the text from the record to the page area. The Move-Line subroutine picks up the next line of text from the input record and places it on the page centered about the memory position given by the Page Locator. Upon entry to the subroutine, therefore, the Page Locator must point to the center of the field where the line is to be placed. It should be noted that Comment lines, as they exist in the input record, are usually less than the horizontal dimensions of the symbol, and this subroutine centers them so that there are equal margins on the right and the left.

The Move line subroutine 20.08 also performs additional functions:

(1) After moving each line, it increases the Page Locator by 120 characters, thereby automatically setting the Locator in proper position for the next line within the symbol.

(2) Adjusts an input record pointer so that the next call on the subroutine will move the next sequential Comment line of the record.

(3) Tests to see when it has moved the last line of an input record to the page area, and sets a signal when this condition is encountered.

When an End-of-Page symbol is located in an input record, block 19.04, the entire page is written to a print

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ape, block **21.05**, or printed directly to the on-line printer **112**. The page layout area is then cleared to spaces and formation of the next page begins with the reading of the next input record.

An End-of-file symbol on the input file signifies that all input data for this program has been processed (blocks **19.03**, **21.03**; at this point Pass **IV** is called in.

Detailed Processing

As outlined above, every record (except **B**-codes) has the location of its symbol computed by means of a subroutine **19.06**, which location is stored in the Page Locator. Separate paths are then taken for each **F/C** code:

B-code

A **B**-code is the only **F/C** code which is not allocated to a column; its sole function is to supply a chart title. When a **B**-code is encountered (blocks **19.05**, **21.07**), its Comments field is stored in the sub-header area of the page memory area, where it remains until overlaid by the next **B**-code. No further processing is needed for **B**-codes.

J-code

A subroutine **20.06** is used to create an octagon of dots on the page, with the vertical connecting line pointing to the midpoint of the top line thereof. The **Box #** and tag (if any) of the record are placed outside the symbol in the proper memory locations (block **20.07**). The Move-Line subroutine **20.08** is bypassed for there are no Comments in **J**-records, and a test **20.09** leads to a branch **22.03** that locates in the Tag Table the destination of the **JUMP** record. The Page # and **Box #** of this destination are taken from the table and placed inside the symbol. If the destination tag is not found in the Tag Table, indicating an undefined tag, the center of the symbol is left blank. This completes the processing, and the next record is brought in (**19.02**).

E-code

The subroutine **20.06** creates the symbol in the page area. The word "EXIT" is then placed within the symbol and the **Box #** of the symbol and tag (if any) placed alongside the symbol (block **20.07**).

H-code

A **HALT** symbol is generated (**20.06**) with the word "HALT" inside of it. Tag, if any, and **Box #** are then placed on the page (**20.07**).

P-code

The Page Locator is first backed up one line (120 characters) and the tag (if any) and **Box #** placed on the page (**20.07**). The Locator is then advanced back to the top line of the symbol, and the top symbol line, consisting of a field of minus signs, is placed (**20.06**) on the page. The Locator is then advanced to the next line, the letter "I" is generated, one on the extreme left and one on the extreme right, to form a part of the vertical boundary lines of the symbol, and a call is made on the Move-Line subroutine **20.08** to move the first line of Comments text to the page. On return from the subroutine, a test is made to see if this was the last line. If not the last line, the logic is recycled back to where the vertical boundary line segments "I" are inserted, and the next parts thereof are inserted, and another call is made on the Move-Line subroutine. This cycle continues until the signal indicating "last line" is set, when the vertical boundaries are complete, as is the Comments area. Then, the bottom line of the symbol, consisting of a field of minus signs, is placed on the page. It should be noted that it is not necessary to adjust the Locator to the next line, since this is a function performed by the Move-Line subroutine.

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N-code

This code is processed identically to the **P**-code, with the following exceptions:

(a) Asterisks are used for both horizontal and vertical boundaries.

(b) Since a **NOTE** symbol has its left side offset by two positions to the right, the Locator is incremented by two, prior to doing any processing. (Incrementing the Locator by two positions sets it to the horizontal center of the offset **NOTE** symbol.)

S-code

The Page Locator is backed up to the previous line, and tag and **Box #** placed on the page. The entire SUBROUTINE symbol is then moved to the page layout area; this symbol is stored in memory as a constant and is moved to the page area from the constant area in a series of moves controlled by a tally (**20.06**). After each line of the constant is moved, the Locator is incremented by 120 positions to bring it to the next line. After the entire symbol is on the page, a subroutine is used to put in the three "I" symbols forming the connecting line leading down to the next box. The Locator is then adjusted back to the second line of the symbol, which is the first line to receive any Comments text (the line directly underneath the upper horizontal boundary). Successive calls are made on the Move-Line subroutine (**20.08**), until the last-line indicator is found to have been set. Test **20.09** leads to extracting (**22.03**) the destination tag of the SUBROUTINE (identified in field **LHT**) from the Tag Table, and its corresponding Page # and **Box #** (from the Tag Table) are placed within the symbol. If the tag is undefined, its Page # and **Box #** are omitted from the symbol. It should be noted that printing of the destination tag in parentheses within the subroutine symbol is not handled by Pass **III**; this field is inserted as a regular Comments line by Pass **I**, and Pass **III** handles it merely as another line of Comments. Also, any vertical editing of lines, for better spacing, is controlled by Pass **I** through the insertion of dummy control symbols representing blank lines, thereby effectively spacing the lines properly within the symbol.

T-code

A **T**-code generates no symbol, but merely results in the placing of text on a page. Successive calls are made on the Move-Line subroutine (**20.08**) until the last-line indicator is set.

Pseudo-Connectors

Pseudo-connectors are short records generated by Pass **II** to indicate a connection from the bottom of one column to the top of the next column. These records (Table **IV**) are of a different format from the other input records, and are identified by the letter "X" in a fixed position of the record. The only information contained within this record, in addition to its Column and Line #, are the Page # and **Box #** to which the connector is jumping. Pseudo-connectors cannot have tags, since they are generated internally, nor are they assigned box numbers. After the address is computed (**19.06**) and the symbol is generated on the page area, test **19.07** leads to block **22.01**, where the destination Page # and **Box #** from the pseudo-connector record are placed within the symbol. It is always necessary to add "1" to the **Box #** before placing it inside the symbol. This is because Pass **II**, when setting up the pseudo-connector record, uses a **Box #** which is one too low. There is no logical basis for this; it is purely a matter of convenience in the implementation of Pass **II**.

DECISION Records

The handling of **DECISION** records presents a far more difficult problem than other codes, primarily because of the many courses of action available on the

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branches from a DECISION. Each DECISION record has two fields, LHT and RHT, in which destination tags are stored, and two indicators LHTC and RHTC that indicate one of the following courses of action:

- (a) This branch requires a connector.
- (b) The coding from this branch is located in the adjacent column.
- (c) Use a line to indicate connection, if possible; if this is impossible then use a connector.

It should be noted here that the mnemonics LHT and RHT do not refer to "left" or "right" side; one of the decisions to be made by Pass III is which side of the DECISION symbol will indicate a particular branch.

Rules Regarding Connection of Branches

In laying out lines from DECISION symbols, the program follows certain pre-defined rules. These are:

(a) Wherever possible, connection lines are used in place of connectors.

(b) Where a connector must be used, the connector is always placed on the right side of the DECISION symbol unless this side is already in use (either by another connector or by a line to an adjacent column) in which case it is placed on the left.

(c) Connection lines may only appear (in this embodiment) in a lane on the left side of the DECISION symbol (23.04 et seq.).

(d) Connection lines in the same lane are allowed to go to different destination symbols so long as they do not overlap (23.08).

(e) Connection lines are drawn (in this embodiment) only if the destination is in the same column as the DECISION symbol (23.07).

Detailed Description of DECISION Processing

The outline of the DECISION diamond is first moved to the page. This outline is carried as a constant within the program and is moved (20.06) to the page area by a series of moves controlled by a tally. The main-flow branch label is then moved from the Comments field to its position on the page and a subroutine is used to drop a vertical connection line to where the next symbol will be. It should be noted that a DECISION symbol has a vertical connect line consisting of four elements, while other symbols use three elements; this is necessary in order to guarantee proper clearance between any connectors and following F/C symbols. The Page Locator is then moved back up to the top of the symbol and the tag and Box # placed on the page alongside the symbol. The Locator is then moved down two lines in place to receive the first line of text. Successive calls are made on the Move-Line subroutine 20.06 until an indicator shows that all lines have been moved. Overflow of the DECISION diamond is not possible at this point; if overflow did occur, the excess Comment was truncated by the Pass-I editing logic.

After the DECISION symbol is completely laid out on the page, with its related tag and Box #, the logic to examine the branches begins (21.01, 22.05). In processing the branches, the same physical coding is used for processing both fields LHT and RHT. The branch tag currently being processed is always located in field RHT. When the tag originally located in RHT is finished, LHT is moved into RHT for its processing. An indicator is used so that the logic knows when it has completed processing the second tag and can go fetch a new record.

It is possible for either or both of the tag fields, LHT and RHT, to be blank; for a two-way decision, field RHT is blank. Both tags may be blank due to an error condition in the source program; in this case a DECISION diamond is printed with no branches. If a field is found to be blank, it is bypassed: thus, the program does not have to formally distinguish between a two- and three-way decision since bypassing blank tag fields automatically

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handles the problem. The following indicators are used throughout the DECISION-branch logic:

(a) Two signals LSS and RSS tell the logic whether the left side and right side respectively of the DECISION symbol have been utilized.

(b) An undefined symbol indicator UNDS is set by the subroutine which searches the Tag Table. This indicates that a tag has not been found in the table.

(c) An indicator TAGTA tells the logic whether it is processing the first or second branch tag.

At the start 22.05 of branch processing, indicators LSS and RSS are set to OFF, indicating that both sides of the DECISION symbol may be available. Indicator TAGTA is set to indicate that the first branch is being processed.

A check is then made to see if the indicator RHTC says to go to the adjacent column, and if so, the fields LHT and RHT are reversed, along with their related indicators and labels. This is necessary in order to insure that the RHT branch is always processed first (an arbitrary convention), and consequently the right side of the DECISION symbol (which leads to the adjacent column) is initially made available for it. At this point, branch processing begins. A test is first made for the presence of a tag in this field. If the tag field is blank, the logic is routed to a test 25.01 of indicator TAGTA, which is described below. A subroutine 22.06 locates the branch tag in the Tag Table; and if found (22.07) the location data for this tag is extracted and stored (23.01) in an index register.

If the tag in question is not found, indicator UNDS is set to ON for later use and test 22.07 routes the program to the branch 25.01 for drawing a connector symbol.

When the tag data is obtained, one of three courses of action is taken depending on the status of the tag indicator (NOTE 23.02), which is then tested (23.03). If RHTC indicates a connection to the secondary (adjacent) column, a horizontal line is extended (24.06) to the right an appropriate number of positions. A vertical connecting line at the right-hand end of the horizontal line is then dropped (two elements in length) to connect with the top symbol of the adjacent column. The label is then placed just above the horizontal line, and RSS is set (24.07) to ON, indicating that the right side of the DECISION has been utilized.

If RHTC indicates (23.03) that a connector is to be used, a test 25.01 is made to see if RSS is ON, if it is ON, then the connector must be drawn (25.05) on the left side of the DECISION symbol. If RSS is OFF, then the right side of the symbol is used (25.02). Depending on the status of RSS, the Locator is either advanced or retarded to the left or right side of the DECISION symbol. The label is put on the page and the connecting lines between the DECISION symbol and the connector are drawn in, as is the symbol itself. After the symbol is drawn, the undefined tag indicator UNDS is tested. If the tag data is undefined, then the tag itself is placed within the generated symbol. Otherwise, the Page # and Box # of the destination are picked up from the Tag Table and placed within the symbol. Either LSS or RSS is then set (25.06 or 25.03) to ON, depending on which side of the symbol the connector was drawn.

The third course of action to be taken is when field RHTC indicates (23.03) that a line should be used if possible (NOTE 23.04). The program determines whether it is feasible to draw a line; if not feasible, then a Connector is used. First LSS is tested (23.05); if it is ON, then the left side of the DECISION symbol is all ready in use. Since connection lines may only be drawn to the left (NOTE 23.06), a Connector must be used on the right side, and the program is routed down that logic path 25.02. If LSS is OFF (test 23.05), then a test is made of UNDS; if this indicator is ON, then a Connector is used and the logic is routed down the path 25.05 for left-side Connectors. However, if the tag data is available, a test 23.07 determines if the tag symbol is in

the same column. That is, the Page # and Col # of the branch tag (obtained from the Tag Table) are compared with the Page # and Col # of the DECISION symbol. If they are not the same, then the program is routed through the Connector logic 25.01. If the Page # and Col # of the branch tag match that of the DECISION symbol, then the branch is in the same column as the DECISION and a line may be feasible. A check is then made to see if the destination of the branch is above or below the DECISION symbol. This is done by comparing the Line # of the DECISION symbol with the Line # of the branch tag. If the destination is below the DECISION symbol, then a "down" line must be used. A check 23.08 is made of the left lane reserved for connecting lines; if it contains an "I" then there is already a line in that column. This line is then traced back to its destination, to see if its destination is the same as the destination of the current DECISION. If so, then a simple horizontal connection is made to the line that already exists. If the destination of the existing line is not the same, then a Connector must be used for current branch tag and the program is routed to that logic 25.01. If the test 23.08 determines that the column reserved for vertical lines is unoccupied, a down-line is drawn in (24.01). The ending point of the down-line is known from the Line # of the destination tag (obtained from the Tag Table). Appropriate horizontal connecting lines are drawn on the page, the label is placed in its proper position, and LSS is set (24.02) to ON.

If the destination tag is above the DECISION symbol, a similar type of logic is followed, with the following exception: in testing for the presence of an existing vertical line, the down-line logic had only to test one location—that element of the vertical-line lane immediately below the DECISION branch point. For an up-line, however, every element of the lane between the DECISION symbol and the destination must be tested. Otherwise, an up-line might interfere with an earlier up-line placed further up towards the top of the page. An additional complication may also arise whereby an up-line may interfere with a "from" connector, which is discussed below.

As explained above, one of several courses of action will be taken for each branch, depending on the setting of RHTC and the feasibility of drawing a line. At the end of each path, return is made to a common point 24.03, where the indicator TAGTA is tested to see if this is the second or first branch just completed. If the indicator shows that the second branch has been processed, then processing for the entire record is not complete and control is returned back to block 19.02 to fetch a new record. If only the first branch has been processed, then fields LHT and LHTC are moved (24.04) into fields RHT and RHTC, respectively. The label for LHT (first five positions of Comments field) is moved into the label area for RHT (second five positions of Comment field), UNDS is set to OFF, and the logic recycled (24.05) to begin the process 22.06 for field LHT, now located in the area previously reserved for RHT.

Cross-References

The information for inserting cross-references ("from" Connectors) is contained in the Tag Table, where it was placed by Pass II. For each tag entry in the Tag Table, five characters are reserved for dealing with cross-references. The first four characters contain the Page # and Box # of the first reference to that tag. The fifth character is a counter of the total number of references to that tag. After a record is read from the input file, a test 20.01 is performed for the presence of cross-reference (this test is made prior to splitting each code down its own branch). This test examines the tag field for presence of a tag, since the absence of a tag indicates that there is no cross-reference. If a tag is present, the tag is looked up in the Tag Table. If the tag is not found in the Table,

then it is undefined and cross-references are not possible. If the tag is found in the Tag Table, a check is made for the presence of a cross-reference by testing the Tag Table field for the first cross-reference; if that field contains spaces, no cross-reference exists.

If a cross-reference does exist, then it is necessary to determine (20.02) the relative position of the current symbol on the page. There are three possibilities:

- (a) Middle of a chain.
- (b) Top of a chain in a main column.
- (c) Top of a chain in a secondary column.

How a cross-reference is shown on the page depends on where the current symbol is located on the page. At this time, the Page Locator is pointing to the center of the top line of the current symbol. This Locator is now backed off to three lines above the top line of the current symbol and a test is made for the character present in this location. From the page design used, it follows that if this character is the letter "I," then the current symbol is in the middle of a chain; if a minus sign, it is at the top of a chain in a secondary column. If the character is a space, the current symbol is at the top of a chain in a main column.

If the current symbol is at the top of the chain in a main column, the cross-reference is placed on the page centered about the column's centerline. The reference placed on the page is extracted from the Tag-Table entry, and an asterisk is inserted if the Tag-Table indicator is set to show more than one cross-reference (20.03, 20.04). If the current symbol is at the top of a chain in a secondary column, then it is not desirable to show the Page # and Box # of the first reference, since this is the DECISION symbol connected by a line to this point, and showing the Page # and Box # here would be redundant and possibly confusing. However, a check 20.03 is made for more than one reference and if there is more than one reference, an asterisk in parenthesis is placed to the left of the centerline. If there is only one reference, then nothing is placed on the second column of the page.

If the current symbol is located in the middle of a chain (either main or secondary column), a check is made to see if there is interference with an existing line. Cross-references for this case are always inserted to the left of the column centerline and are two lines above the first line of the symbol. If there is a vertical down-line coming into this point from a DECISION symbol above, then this horizontal line is already occupied. The check is made by positioning the Page Locator to two lines above and one space to the left and checking the resultant location for a minus sign. If there is no minus sign, then the line is free; the first cross-reference is placed on the page, along with an asterisk if the indicator in the Tag-Table entry is set for more than one. If there is a minus sign in that location, a connecting line is being drawn and there is no need to place the Page # and Box # on the line, since this usually is the same as the symbol from which the line is drawn. In this case, a test is made for more than one reference and, if found, an asterisk in parenthesis is placed on the line, if there is only one reference, then no action is taken. The asterisk which may be placed on this line becomes part of the horizontal line coming into the centerline and thus provides notice to the user that there is at least one more reference besides the one shown via the connecting line.

If a cross-reference is placed to the left of the centerline, in the situation where the symbol is the middle of a chain, it is still possible that the subsequent drawing of an "up" connecting line will erase it. This can happen because up-lines come from symbols which are further down in the column and have not yet been processed. It cannot happen with down-lines, since they must come from symbols above the current one and hence will have already been drawn. That part of the DECISION branch logic which places up-lines on the page tests for the pres-

ence of an asterisk prior to drawing the horizontal connection back to the centerline of the column. If an asterisk is present, its position is moved up over the connecting line. If an asterisk is not present, then the line is either free of interference or there is but a single reference, which must be the one for which the line is presently being drawn. In either case, the horizontal connection line back to the centerline can be put in without any complication.

When the end-of-page indicator is detected (19.04) in the last input data block, the F/C page is complete in the memory layout. Thereafter, the entire page is put out to the printer (21.05), and the next page is started in the same fashion as described above. When the last page has been printed out, test 19.03 detects an end-of-file indicator to initiate (21.03) the rewinding of the tapes and calling in of Pass IV.

PASS IV

Production of the Cross-Reference List is accomplished by a separate pass, following the completion of the last page of flow chart. Input to the Cross-Reference pass is the same tape 111 that served as input to Pass III; output consists of the Cross-Reference List, either to an on-line printer 112 or to a magnetic tape for off-line purposes.

Two tables, both kept entirely in memory, serve as the basis for producing the listing. These tables are designated:

- (1) Abridged Tag Table
- (2) Reference Table

Tag Table

The Tag Table used by this pass is an abridged version of the main Tag Table (Table II) used by the first three passes. Each entry of the Abridged Table consists of the following items:

- (1) Name of Tag
- (2) Page # and Box # assigned to this tag
- (3) First reference (Page # and Box #) to this tag
- (4) An indicator which tells whether there are more references to this tag
- (5) Memory address of the Reference Table entry containing the next reference (if there are any more references).

At the start of the Cross-Reference pass, the main Tag Table is still in memory from the previous pass. The first job 27.01 is to set up the Abridged Tag Table from the main table. Since each entry of the Abridged Table is shorter than its corresponding entry in the main table, the same physical memory area may be used for the Abridged Table. Every tag in the main table has a notation as to whether that tag is referenced by another symbol (created by Pass II for the flow chart layout). Only those tags which have references to them are moved to the Abridged Table; all others are dropped.

Each entry of the Abridged Table has four characters reserved for the first reference to this entry. This information is already available from the main Tag Table. However, for ease of implementation, this information is not transferred between tables, but is dropped. The only information transferred between the two tables, therefore, is the name of the tag and the Page # and Box # of that tag. Room is reserved in each entry for the remaining three items, which are filled in later on in the pass.

After all appropriate entries from the main Tag Table have been transferred to the Abridged Table, the latter is internally sorted 27.02 into a Page # and Box sequence. Any one of several known sorting techniques may be used for this purpose.

Reference Table

Each entry in the Reference Table consists of three items:

- (1) Identification of this reference (Page # and Box #).
- (2) An indicator telling whether this is the final reference or whether there are additional reference.
- (3) Memory address of the entry in the reference table containing the next reference for this tag (if another reference exists).

The three fields of each reference entry are identical in format to the final three items of each Abridged Tag Table entry. One entry is created in the Reference Table for each reference (after the first one) for any given tag. A chaining technique is used to connect these references back to the Abridged Table entry to which they refer. Thus for any given tag, the first reference is in the Abridged Table entry, and succeeding references are spread out throughout the Reference Table, with each reference giving the location of the next reference in the chain. The chain is ended when the indicator in a particular entry says that is the last reference.

It should be noted here that each Abridged Table entry contains the first reference to that tag. Consequently a tag has entries in the Reference Table only if there is more than one reference. This choice of format was made for reasons of efficiency; that is, for most flow charts, the great majority of tags have only one reference, and for these, the Cross-Referencing can be handled entirely within the Abridged Table itself eliminating the need for access to the Reference Table. Other techniques may be used to collect the cross-reference data which has been developed by the F/C program and to present it in a simple table.

Setting up Reference in the Tables

The input file 111 to Pass-III also serves as the input file to this pass. Only the records which represent J-, D- or S-codes need be processed, since they are the only codes which involve "jumps" to other locations; consequently all other records may be bypassed without any processing.

Each D-, J- or S-record contains three fields that are of interest to this pass. The first field is the Page # and Box # assigned to the symbol on the flow chart. The second and third fields are the destination tags to which a transfer is called for by these records. In the case of J- and S-codes, only one of the latter fields contains a tag; for a D-code, either one or both of the fields contains tags, depending on whether the decision has one or two branches. It should be noted that the input records do not contain the Page # and Box # of the destinations, but only the tags of the destinations.

Prior to reading any input records, a locator must be set up for the Reference Table. This locator always contains the current RHE-plus-1 of the Reference Table. Since successive Reference Table entries are constructed extending to the right in memory, the locator always contains the memory locations at which the next entry is to be created. The Reference Table immediately follows the Abridged Table in memory, and is placed in its initial condition prior to starting the input file. Therefore, as each new entry in the Reference Table is created, the locator is incremented by a fixed constant.

When a D-, J- or S-record is read via blocks 27.03, 27.04, 27.05, the first destination tag is extracted 27.06 from the record. A search 28.01 is then made of the Abridged Tag Table to locate the tag entry; if an entry is not found for that tag, the destination is undefined and no cross-reference is made. Upon locating the Table entry for that tag, the first-reference field is examined; if it only contains spaces, the current record is the first reference to this tag. The Page # and Box # of the input record (i.e. of the D-, J- or S-symbol) is then placed in the "first reference" field and the indicator in

he entry is set to "last reference" status, which only indicates that the reference just inserted is thus far the final link of the chain for this entry.

However, if the first-reference field of the entry is already filled, then an entry in the Reference Table is created **28.02** for the reference from the current record. To create a linkage to previous reference (**28.02**), the existing chain of references for this tag is traced down to its end via a simple loop. The indicator in the Abridged-Table entry is tested; if set to "last reference" a new entry in the Reference Table is created at the next available address set up by its locator by extending the Table to the right by the length of the new entry. This address of the new Reference Table entry is placed **26.03** in the Abridged Table entry of the tag and its indicator is set to "not last reference" status. The Page # and Box # of the D-, J-, or S-symbol are extracted from the input record and placed in the newly created entry, and the indicator of the new entry is set to "last reference" status. The Reference-Table locator is incremented so that it again contains the RHE of the table. On the other hand, if the Abridged Table entry is set to "not last reference," the address of the Reference Table entry containing the next reference is picked up from the Abridged Table entry, and the indicator of that Reference Table entry is tested. The latter entry is either the "last reference" or it in turn leads to the next reference. Eventually the "last reference" entry in the Reference Table is located, and a new Reference Table entry is created, and filled in at the next available address set up by the locator. This address is placed in the previous Reference-Table entry for this tag, and the indicator therefore is reset to "not last reference." The Reference-Table locator is incremented to supply the next available address for any new entry to be created.

The above process is repeated if the input record is for a D-code and contains a second destination tag (via test **28.04**, NOTE **28.05**, and blocks **28.06**, **27.06**). After processing the second tag, or after the first tag processing if the second tag is not present, a new input record is read **27.03** and the entire process recycled. This process continues until an end-file indicator is found **27.04** in the input; construction of the Tables is then complete, and all the information needed for the listing is now contained within the Tables. Accordingly, the program transfers to the output section **28.07** of this pass.

Output

Production of the Cross Reference List consists of combining and printing out the contents of the two internal tables. Each entry in the Abridged Tag Table produces at least one line on the listing. Additional lines are used if the number of references to a particular tag overflows the amount of room available on the first line. The Abridged Table entries are handled successively with a locator being set **28.07** to the initial entry. The tag name and its Pages # and Box # is moved **28.08** from the Abridged Table to the output area; the first-reference field of this Table entry is also moved to the output area. The indicator of the entry is then tested for "last reference"; if not the last reference, the address of the next reference in the Reference Table is picked up **28.09**, **29.01** from that entry. The Page # and Box # of the next-reference entry is moved **29.02** to the output, and the indicator of the new reference entry is tested **29.06**, and the logic recycled to block **28.09** if "last reference" is not found. The process is then repeated to locate the next entry and extract the desired data therefrom. When the indicator of any link specifies "last reference" (test **29.06**), the cycle ends and the current output line is printed **30.01** or written to the listing tape. Appropriate locators and counters are maintained for controlling placement of the references in the output line. When a counter indicates **29.03** that the output line is filled, the

line is printed **29.04** or written to tape and a new line begun by checking **29.06** for further references.

The above processing is repeated by checking **30.02** for more tags in the Abridged Table, advancing a locator (**30.03**) to the next entry thereof, and recycling via block **30.04** and **28.08** to repeat the process until all entries in the Abridged Tag Table have been processed. After processing the final Abridged Table entry, test **30.02** determines that the listing is complete, and the program transfers to block **30.05** to "wrapup" any housekeeping details, such as rewinding the tapes, and the operation terminates (**30.06**).

The Cross-Reference List affords a valuable body of information that assists in reading and studying the flow chart. That is, each entry point marked with a cross-reference is known to have but a single transfer into that point, except where it is marked with an asterisk. In the latter case, the Cross-Reference List provides, under the tag of the entry point, a complete list of all other such transfers, which makes it possible to determine various interrelationships of the documented program.

MODIFICATIONS OF THE INVENTION

By modifications of the flow chart documentation system of this invention other forms of flow charts may be produced, such as those having characteristics illustrated in the fragmentary charts of FIGS. 4A and B, 5 and 6. FIGS. 4A and B and 5 present diagrammatically the interrelationships of D-symbols, each represented by a diamond **150**, and the other types of symbols all represented, for simplicity, by a rectangle **152**, except for J-, E-, H-symbols and connectors which are represented by circles. A four-column chart is assumed by way of example.

These flow charts may have one or more of the following features:

(1) A branch chain from a main flow column may be presented in an adjacent secondary-column as described above, and in addition the branch chains from the secondary-column may also be illustrated in the next adjacent column. See FIG. 4A where column **154** contains the main flow, column **156** contains branch chains from DECISIONS **150** and **162**, column **158** contains a branch chain from DECISION **164** in column **156**, and column **160** contains a branch from DECISION **166** in column **158**.

(2) A branch chain may be entered at an intermediate point of that chain as well as from the first block of the branch chain and branch chains may be shifted up or down so that they fit in the available space. See FIG. 4A where branch chains **158** and **160** are entered at intermediate points, chain **160** is shifted up and chain **170** is shifted down.

(3) A branch chain need not be entered directly opposite the branch output of the DECISION in the main column; the branch connecting line may be formed as a combination of horizontal and vertical line segments so that the branch chain may be positioned in any suitable place within the adjacent column. See FIG. 4A, branch chain **160** and connecting line **168**, and branch column chain **170** and line **172**.

(4) If a branch chain is not provided in a column adjacent to the main column, that adjacent column may be used for the continuation of the main flow, and all four columns of a page may be used for the main flow where appropriate and where branch chains are not or cannot be illustrated. See FIG. 4B, columns **174** and **176**. Each column has two possible vertical-line lanes, one on each side of the symbol, to permit connection in the same or adjacent columns (e.g., the lanes for lines **178** and **180** of column **158**, and the lanes for lines **182** and **184** of column **160**). The vertical lines can be connected up or down in each path. Thereby, in a four-column chart, eight vertical-line lanes are available for appropriate interconnections (and the use of all eight is illustrated in FIG. 4A).

(5) Vertical and horizontal lines may cross, (e.g., lines 168 and 180 in FIG. 4A) but provision is made to try alternative non-crossing paths.

(6) Connecting lines may be drawn between any two of the four columns, and these connecting lines may be directed either from left to right or from right to left, and may be a combination of vertical and horizontal line segments (e.g., line 168 of FIG. 4A, and lines 186 and 188 of FIG. 4B).

(7) Unconditional transfers (jumps or exit instructions) are represented by a line being drawn wherever possible, either to the same or to another column on the page (e.g., lines 186 and 188 of FIG. 4B). Similarly, pseudo-connectors are avoided where connecting lines can be drawn to the same page.

(8) Branch chains are connected either to the left or to the right, or both, of the main column containing the DECISION symbol from which the branch or branches occur (e.g., in FIG. 5 the DECISIONS in main-flow column 190 have respective branch chains 192 and 194 that are presented in columns on opposite sides of main column). Thereby, any column may be used for the main flow or for branch chains.

(9) A branch chain is picked up and printed if it fits in the space remaining on a page, be it one or more columns.

(10) Each flow chart page may be developed as a cluster of chains, with any one column or columns containing the main-flow logic and the remaining column or columns containing the chains branching from the main-flow.

A form of flow chart incorporating the last two features is shown diagrammatically in FIG. 6, in which the main-flow column of logic is illustrated in a simplified fashion by a relatively wide strip, and branch chains by a narrower strip (so that they can be readily distinguished) and JUMPS and pseudo-connectors at the ends of columns by circles. The simplified diagrams of FIG. 6 indicate the branching of chains from DECISIONS of the main-flow logic and from DECISIONS of the branch chains; the various F/C symbols are omitted to illustrate the general nature of the flow chart configurations that are handled. The aforementioned features of FIGS 4 and 5 are applied in illustrating the "cluster" feature.

FIG. 6A illustrates the four columns 200, 201, 202, 203 used for the main-flow logic (where no branches from DECISIONS occur that would fit in the remaining space on the page). The successive columns are connected by lines; alternatively pseudo-connectors may be used to terminate each column. FIG. 6B illustrates a column 204 of main-flow logic, from a DECISION of which a branch chain 205 is connected; and from a DECISION of the latter a sub-branch chain 206 is connected; and another sub-branch chain 207 connects from a DECISION in branch chain 207. Where only a single branch chain (e.g., chain 205) develops from the main-flow logic 204 and does not itself develop additional branch chains, only the main-flow logic 204 and branch chain 205 are printed on the page. Thereafter, the next page continues initially with the development of the main-flow logic and with the processing of branch chains as DECISIONS arise (and in the manner described with respect to FIG. 8). FIG. 6B illustrates the facility of displaying sub-branch chains to the right of the main chain.

FIG. 6C illustrates in the first column the main-flow logic 208, from a DECISION of which there is a branch chain that is a long one and has sections 209, 210, 211 in three remaining columns of the page. The single column of main-flow logic and the single branch chain make up the page. If the branch chain terminates in the second or third column, the page likewise terminates.

FIG. 6D shows two columns 212 and 213 of main-flow logic and a branch chain 214 from a DECISION in the second column 213, as well as a second branch chain 215 from a DECISION of the first branch 214. Where the second branch chain 215 is not suitable for presentation in the fourth column, the page terminates with branch

chain 214 and presents a cluster of the three columns 212, 213, and 214. This cluster feature of the F/C program does not attempt to use all of the available page space, but rather it is constructed to display as much of the branch interrelationships of the program being documented as the page size limitations permit. For practical reasons, the page size limits the amount of information that is presented as a unit.

FIG. 6E shows a column 216 of main-flow logic with a branch chain 217 connected from the right side of a DECISION thereof, another branch chain 218 connected from the left side of a DECISION thereof, and a sub-branch chain 219 connected from the left side of a DECISION of the left branch chain 218. FIG. 6E illustrates the cluster feature of presenting branch chains on either side, or both sides of the column containing the main-flow logic, and the feature of sub-branch chains being allocated to the left for versatility in the display of branch information in each cluster.

In implementing the feature of forming successive pages as "chain clusters," the F/C Control Program is constructed to start each page with a column of main-flow logic. Upon reaching a DECISION record, the F/C program branches in the manner described above for the first embodiment, and processes a branch chain from that DECISION. In the course of processing that chain (or a second chain from the same DECISION), further branches may be developed from DECISIONS within the branch chain. If these sub-branches can be presented on the same page, they are developed in the same fashion. The information regarding the space required for the first branch is already known before the second or succeeding sub-branches are processed; therefore, the available space for the second branch or for sub-branches is then known, and the F/C can determine whether the sub-branches are suitable to be placed on the same page or not.

The implementation of an automatic system for producing a flow-chart formed from clusters of chains may assume different forms including that of a three-phase program of the same general type as shown in FIG. 3 (the fourth, cross-reference phase is optional). The detailed program logic may be similar to that described above in connection with FIG. 8, with certain modifications of Pass I and II, as described hereinafter.

Pass I

This logic is the same as described above for FIGS. 8/1 to 8/6, with the following additions. Each chain is given a separate identifier (e.g., a sequence number) in addition to the one already provided in the chain table, and this chain identifier is used in the Tag Table to associate all Tags that are part of a particular chain. Thus, the first chain is so identified and set forth in the Tag Table, and all Tags occurring in the first chain are listed in the Tag Table under (or in association with) the first chain. Thereafter, each time a new entry is created in the Chain Table for a chain ten ending (block 06.04, FIG. 8) a corresponding new entry is likewise created for the Tag Table for the following chain, under which all associated tags are listed. An additional field is also provided in the Tag Table, which is used to furnish the relative line position of a particular tag within its associated chain. This information is available from the cumulative count in Line Counter B (block 04.01) and inserted in the field upon the creation of each Tag Table entry (block 05.01). Thereafter, a search for a particular tag in the Tag Table supplies the identifier of the chain in which the tag lies together with its line position within that chain.

Pass II.—Main Chain Processing

The general logic for this portion of the processing is shown in FIGS. 7A to E; details will be apparent from the following description and from the foregoing of FIG. 8. Upon the start 300 of the program, block 302 operates

to reset counters and work storage areas to their proper conditions. The input data is the output tape from Pass I, in the manner described above. Thereafter, block 304 locates the next main chain via a pointer in the Chain Table; the first chain is assumed to be a main chain, and succeeding chains are also assumed to be main chains until they are assigned as auxiliary chains. Test 306 determines if there are any more main chains; and if not, the program branches to block 308, which operates to rewind the tapes and bring in Pass III. If there is another main chain, the program proceeds with block 312 to read the next symbol record from the located chain. Test 314 determines whether the symbol is a J or an E, and if so, the program branches to a connector 316 leading to the auxiliary chain processing, ACP, described below with respect to FIG. 7B. If the current record is not the end of a chain, test 318 determines if the current record is at the end of a main column (i.e., if EPCON is exceeded, see block 08.01 of FIG. 8); if so, the program branches via block 320 to ACP as indicated by connector 324. Block 320 creates the appropriate pseudo-connector symbol used to identify the end of the column and the page block number to which the program connects from that point, and it provides a supplementary record in the Chain Table. Each chain when processed, be it a main chain or an auxiliary chain, is marked in the Chain Table as processed. Under the circumstances of block 320, a main chain cannot be so marked; however, the main chain can be marked as partially processed, with a store of the return record number to which the program will go to continue the processing after the remainder of the current cluster is completed.

If the current record is not at the end of the main column, the program continues with block 326, which stores the record in a storage in memory that receives in order all the records making up a page and which starts with address LM, and also steps a counter LC to provide the next address in the record storage area for receiving the succeeding record; it also steps Line Counter-I an amount corresponding to the number of lines in the current record to obtain the line number in the main-chain column of the succeeding record. Thereafter, test 328 determines if the current record is a DECISION, and if it is, the program steps to subroutine CHSUB 330, which determines if the branch chains from that DECISION can be pulled and inserted in auxiliary (adjacent) column. An indicator is set to identify for CHSUB that branch chain is a transfer from a main chain to distinguish from transfers from another branch chain. After subroutine CHSUB, the program returns to the main-chain processing at block 312. If the current record is not a DECISION, the branch from test 328 is also back to block 312 to read the next symbol record from the main chain and repeat the above-described process.

Successive records of the main chain are processed in this manner until test 314 finds an end-of-chain record or test 318 finds an end-of-column record. In either case the program branches to ACP. The main chain processing is not interrupted for auxiliary chain processing, though it is interrupted for CHSUB to determine whether an auxiliary chain is appropriate to be pulled; all branch chains are examined to determine their suitability for display as an auxiliary chain. The actual processing of such auxiliary chains that are found follows the completion of the main chain processing. The actual assignment of page and block numbers to the main and auxiliary chain records is performed after the branch chains have all been identified and pulled and their records stored in memory.

FIG. 7B illustrates the logic flow for the subroutine CHSUB, which determines whether a branch chain can be pulled and utilized in an auxiliary column. Initial block 334 operates to locate the chain name for the tag to which the program branched from a DECISION record. The tag may be at the start of a chain or anywhere within a chain, and a search is made of the Tag Table

to locate the associated chain name which was stored during Pass I. The line position of the tag within the chain is also extracted from the Tag Table. Thereafter, block 336 performs a calculation to determine CPPOC, the preferred position of the chain in the auxiliary column. CPPOC is calculated by subtracting the line position of the tag in the chain from the line position of the DECISION record from which the branch takes place. If CPPOC is a negative value, it is reset to zero; this represents a condition of the tag having a line position lower down on the page than the branch-point of the DECISION and CPPOC cannot be assigned a useable line number. Thereafter, block 338 stores CPPOC in a temporary storage field until it is determined whether that chain can be pulled for an auxiliary column. In addition, three sets of stores in memory are provided to hold the names of branch chains that may be pulled for auxiliary columns together with other chain-locating information such as CPPOC; these temporary stores are identified as CR, CL, and CB, representing respectively branch chains to the right, those to the left, and sub-branch chains that branch from the left or right auxiliary chains currently being processed. Thus, the CB stores contain a buffer storage of sub-branch chains, which branch to the right of right auxiliary chains or to the left of left auxiliary chains, whichever is currently being processed, and which sub-branches are to be processed thereafter. Each set of these chain-locating records may contain an arbitrary maximum number, say 10, which indicates the maximum number of branch chains that may be actually utilized in a particular auxiliary column.

Thereafter, test 340 determines whether the branch chain being investigated is a branch from a main-column chain or from another auxiliary-column chain. If it is from a main column, test 342 then determines whether the length of the current branch chain plus the length of the other branch chains already assigned to CR-1 to 10 would be greater than the column length EPCON. If not greater, there is still room in the right auxiliary column for the current branch chain, and block 344 sets up the next chain-locating record CR and stores the name of the current chain, its length, sequence number, and CPPOC in the appropriate fields thereof. Thereafter, test 346 determines whether there is a second, unprocessed branch in the current DECISION record being processed; if not, processing exits from CHSUB. If there is a second branch, block 347 locates the chain and calculates CPPOC for the second branch and test 348 determines whether the length of this branch chain together with the combined lengths of the other left auxiliary chains already assigned are greater than EPCON. If EPCON is exceeded, the program exits from the subroutine; if it is not, then block 350 stores the chain information in the appropriate fields of the next one of the chain-locating records CL-1 to 10 for the left auxiliary column, and the program exits.

If test 342 indicates that the branch chain is too large for the right hand column, test 352 determines whether the branch chain is suitable to fit in the first auxiliary left hand column. If it does not fit, test 354 determines if there is a second, unprocessed branch chain from the current DECISION record, and if not, the program exits from the subroutine; however, if there is a second branch, an indicator is set for processing the second branch and the program recycles back to the start of CHSUB to process it in the same fashion as the first branch chain was processed (and the latter is identified as processed).

If test 352 indicates that the branch chain will fit in the left auxiliary column the program branches to block 358, and the chain information is stored in the next left chain-locating record CL-1 to 10. Thereafter, test 360 checks to see if there is a second branch chain from the DECISION record, and if not, the program exits from the subroutine. If there is a second branch, block 360 locates the chain and CPPOC is computed and stored in

the manner described above, and test 362 determines whether it will fit in the right auxiliary column (the left having already been preliminarily assigned); if not, the program exits from the subroutine, if it will fit, block 364 stores the chain information in the next right chain-locating record CR-1 to 10, and the program exits.

If the test 340 indicates that the chain is being pulled by an auxiliary column, test 366 determines whether it is from the right or left auxiliary column from the setting of an indicator. If from the right, block 368 inhibits the pulling of any chains to the left and inhibits pulling more than one branch chain from the current DECISION record; this inhibition of pulling to the left, once initiated, continues for the remainder of the sub-branches pulled from the current auxiliary column being processed. Block 368 also sets an indicator to store the information for sub-branch chains from the current right auxiliary column in buffer chain-locating records CB-1 to 10, and thereafter the program continues with the processing from test 342 in the manner described above. Similarly, if test 366 indicates that the chain is being pulled from the left auxiliary column, block 370 inhibits the pulling of any chains to the right column, inhibits the pulling of more than one branch chain from the current DECISION record, and sets up CB-1 to 10 to receive the information of sub-branch chains from the current left auxiliary column. Thereafter, the program continues with the processing of the chain via test 342.

ACP

The processing of auxiliary chains, ACP (FIG. 7C), begins with a subroutine SCHPL 380 (FIG. 7D) for calculating the exact starting line position SCHPL of each branch chain from CPPOC (the preferred starting line position of the chain with the tag opposite the DECISION branch point as calculated by block 336, FIG. 7B) and storing the new value in the field FCPPOC of the associated one of the records CR, CL, CB. Each time the program passes through SCHPL and performs the calculations, a counter is incremented so that its count represents the number of auxiliary columns processed thus far. Thereafter, test 382 determines whether any right or left branch chains remain to be processed (by examining the contents of CR-1 and CL-1, as explained below) and whether any of the columns available on a page for a cluster remain unassigned; and if so, test 384 determines whether they are right auxiliary chains. If so, block 385 locates the chain specified in CR-1 on the input tape so that it can be processed. In addition, an indicator is set for the CHSUB subroutine to identify any sub-branch chain as coming from a right auxiliary chain.

Thereafter, block 386 begins the processing of this auxiliary chain by reading the first symbol record thereof from the input tape. Test 388 determines if the current record is an end-of-chain record (J or E). If not, block 390 stores the record in the next available location in memory as indicated by the address LC (appropriate marker signals are provided at the beginning of the records for each auxiliary column), and thereafter LC is adjusted to indicate the next available memory location for the next record to be stored.

Test 392 determines whether the current record is a DECISION record, and if not, the program returns to block 386 to read the next record and process it in the manner described. If the current record is a DECISION record, subroutine CHSUB is entered (with an indicator set to identify that the branching is from an auxiliary column) to determine whether the sub-branch chain from the current auxiliary chain can be pulled. In the subroutine CHSUB (FIG. 7B) test 340 steers the program down the auxiliary processing section and test 366 determines whether it is a right or left auxiliary currently being processed to provide appropriate steering, in the manner described above. When CHSUB is processing an auxiliary column, the sub-branch chain information is stored in the records

CB. In addition, five counters CTR-1 to 5 are used by CHSUB to maintain cumulative counts of the chain lengths for the five possible auxiliary columns, the three right and two left columns, in order. The right or left column indicator identified which type of column and the number of successive auxiliary column passes through CHSUB determines which column in order is being processed. Upon existing from CHSUB, the program returns to block 386 to process the next record in the branch chain for the current auxiliary column.

This loop continues until test 388 indicates that an end-of-chain record has been reached and the program branches to block 396, which proceeds to store the record in memory and adjusts the setting of address LC. Thereafter, test 398 determines whether there are any more chains in the current column; this test may be performed by determining whether CR-2 contains any data. If there are more chains, block 400 shifts the contents of CR-2 to 10 into CR-1 to 9, respectively, so that the previous contents of CR-2 are stored in CR-1, CR-3 in CR-2, and so on. Thereafter, this routine begins again at ACT to process the chain now specified in CR-1 in the manner described. If test 398 indicates that there are no more CR chains to be processed, the program branches to block 401, which moves the contents of CB-2 to CB-10 into CR-1 to 9, respectively; thereby, the chain-locating records for the sub-auxiliary column are moved into position to be processed. Test 402 then checks an indicator to determine whether the auxiliary column just processed was a left or a right auxiliary; if a right auxiliary, then the new sub-auxiliary column is also a right auxiliary column, and the program returns to ACP to start the auxiliary chain processing for that sub-auxiliary column. If the auxiliary column just processed was a left column, the program branches to block 404, which sets a field LL-2 to the current value of LC; LL-2 contains the memory address of the first record of the second left auxiliary column, and the program returns to ACP to begin the processing of that second left auxiliary column. An indicator is set to inhibit entering the subroutine CHSUB for any branches from this second left column, assuming a maximum of two left auxiliary columns. Similarly, that inhibit indicator is also set when test 402 steers the program down the "right" branch the second time, since the program is then starting to process the third right auxiliary, which is assumed to be the maximum; when a third auxiliary is processed, there is room (in the assumed 4-column example of a page) only for the first left auxiliary, and accordingly the processing of the second left is inhibited.

The indicator for identifying whether the auxiliary column being processed is a right or a left is set initially during the first pass through test 384 of ACP, which determines first whether there are data in the contents in CR-1. If there are, then it is known that it is a right auxiliary; and it is processed first on a priority basis with the right column indicator being set. If no data is in CR-1 it must be a left auxiliary (since test 382 had indicated that there is an auxiliary), and the program branches to block 406, which sets the left column indicator and sets LL-1 (the address of the first record of the first left auxiliary column) to the current value of LC. Thereafter block 408 moves CL-2 to 10 into CR-1 to 9, respectively, so that the left auxiliary chain-locating records are in condition to be processed in the same manner as the right auxiliary chain-locating records, and the program returns to the start of ACP for processing. After the first left auxiliary is processed, test 402 steers the program via block 404 to process the record left auxiliary. After the left auxiliary columns are processed, test 382 finds both CR-1 and CL-1 empty of data, and steers the program to the output subroutine 409, OUTSUB, from which it returns to the beginning of the pass at STRT, FIG. 7A, to process the next main chain and from the cluster therefrom.

In operation, the first right auxiliary column chains are

processed initially by following the contents of the CR records, and then the second right auxiliary column, if any, is processed with the sub-branch chains from the first right auxiliary column using the records in CB, which are transferred to CR for the processing operations. Thereafter, the third right auxiliary column, if any, is processed using the sub-branch chains that were pulled from the second right auxiliary column; the records for the third right auxiliary column are set up initially in CB as the second right auxiliary column is being processed. After all of the right auxiliary columns are processed, if space permits, test 384 steers the program to processing of the left auxiliary chains via blocks 406 and 408, with the chains for the first left auxiliary columns being located by means of CL records, which are relocated into CR. After the first left auxiliary column is completely processed, test 398 indicates that there are no more chains for that column, and the program branches to block 401 to move the chain records CB for the second left auxiliary column into CR. Test 402 steers the program to block 404 to set up the address of the first record of that second left auxiliary column, and the program then proceeds to process those records.

When the records are written out to memory, the address of the first record of the main chain is LM, and the address LC is then used for storage thereafter of the successive records of the main chain, followed by those of the first-right, the second-right, and the third-right auxiliary columns in that order. Thereafter, the first-left and second-left column records are stored in that order. The records of the different columns are separated by appropriate marker signals. LL-1 and LL-2 provide the starting addresses of the records for the two left columns.

The subroutine SCHPL, shown in detail in FIG. 7D, calculates for each chain in an auxiliary column, its exact starting line position and stores it in the field FCPPOC of CR-1 to 10. It starts with CPPOC, the preferred line position for the start of the chain, already stored in FCPPOC (block 338), and the subroutine terminates with the exact position determined. In addition, the subroutine starts with the fields CTR-1, 2, 3, 4, 5 (developed by SCHUB) which contain respectively the total number of lines required by all of the chains in the first, second, and third right auxiliary columns and the first and second left auxiliary columns, in that order. The cumulative counts in the CTR fields are based on the packed lengths of the branch chains; i.e., it is assumed that the first branch chain in each auxiliary column starts at the first line and succeeding chains are positioned thereafter without extra spaces therebetween. This subroutine terminates with the CTR value adjusted to include any spaces inserted by readjustment of chains downward within the associated auxiliary column. Thus, SCHPL starts with the branch chains fitting in a column at least if they are moved up all the way, and proceeds to determine if they also fit when moved down to prepared positions.

Initially at the start of each column, block 410 sets CHPL equal to "1," and thereafter, each chain in the current auxiliary column is processed in order. Test 412 looks for the next branch chain in CR-1; if there is none, the subroutine exits. If a next chain is set up in CR-1, test 414 determines if the stored value of CHPL in working storage is greater than the value of the field FCPPOC of that branch chain; if it is, block 416 stores the value CHPL in FCPPOC as the exact line position for the chain. That is, if CHPL is the greater value, the starting line position is already far down in the column, and no further downward adjustment of the chain is desired; and CHPL is therefore used as the starting line position. Thereafter, block 418 adds the line length of the current chain to CHPL to obtain a new value of the latter, so that the initial CHPL value of "1" may be applied only to the first chain in a column. Thereafter, the subroutine returns to test 412 and exits. If test 414 indicates that CHPL is not greater than CPPOC, the pro-

gram branches to determine if space is available in the column to move the chain down so as to use CPPOC for its starting line position. Only if the full space is available, will the chain be moved down. Block 420 calculates TEM, equal to the difference between CPPOC and CHPL, and representing the desired downward displacement of the branch chain in number of lines. Thereafter test 422 compares the constant EPCON (the column length in lines) with the quantity of TEM plus CTR for the current auxiliary column; if EPCON is the lesser, there is no room in the column for downward adjustment of the chain, CPPOC cannot be used and the program branches to block 416 to store CHPL in FCPPOC of the current chain and proceeded in the manner described above. If test 422 indicates that there is room in the column for adjustment, CPPOC remains unchanged in the CR record. The program continues with block 424, and CTR for the current column is augmented by TEM and the new value of CTR is stored in its own field. Thereafter, block 426 moves the line number in FCPPOC to working storage for a new value of CHPL; since the field FCPPOC remains unchanged, the preferred starting line position is actually used for the chain. The program then continues with block 418, where the value of CHPL in working storage is augmented by the current chain length to obtain a new value of CHPL for the next chain, and the program returns to test 412.

The output subroutine OUTSUB is shown in FIG. 7E; it is entered after completion of the auxiliary chain processing and processes each cluster of columns by determining whether a new page is required and assigning the page, box, column and line numbers in each flow chart record in memory and writes the records to the output tape.

Initially, test 430 determines whether a new page is required; this test involves checking the number of unused columns, if any, on the current page (i.e., the columns required by the previous cluster or clusters) and comparing it with the number of columns required by the current cluster. If a new page is not required, block 432 adds "1" to the column number; if a new page is required, block 434 operates to add "1" to the page number and to set the column # to "0" and box number to "1" (the box numbers on each page are assigned sequentially in each column and in order from column to column starting with the column on the left). Thereafter, in either case, test 436 determines if there are any left auxiliary columns in the current cluster by checking the contents of LL-1 and 2; if not, the main column is the first column of the cluster and is displayed to the left on the page, with the other columns of the cluster to the right in order. Block 438 sets a pair of output pointers to LM (which is the start of the record area in the memory containing the first record of the main column chain) and to LC (the end of the last right auxiliary record) thereby bracketing the memory area of records to be processed. Thereafter, these bracketed records are computed and written to the output tape by the RITEOUT subroutine 440, and upon exiting therefrom, the OUTSUB subroutine also exits as shown by connector 441.

If there are left auxiliary columns, test 442 determines whether or not there is a second left auxiliary column; if not, block 444 sets the output pointers to start with LL-1 and to bracket the memory area of the records for the first left auxiliary; thereafter, the RITEOUT subroutine 446 processes those records, and upon exiting goes to block 438 to set up RITEOUT for the records of the remaining columns. If there is a second left auxiliary column, block 448 sets the output pointers to bracket the memory area for the corresponding records, and the RITEOUT subroutine 450 processes those records. Thereafter, the program passes to block 444 to initiate processing the first left auxiliary records, and so on. In this fashion the records are completed and written

out starting with the column which should appear on the left on the flow chart.

In the RITEOUT subroutine (details are shown in block 440), an initial test 452 determines if there are any more records in the bracketed memory area to be processed, and if not, the subroutine exits. If the bracketed memory area contains additional records, a test 454 determines whether the current record is at the start of a column. If it is, block 456 adds "1" to the column number. Thereafter (or if test 454 proves negative), block 458 adds "1" to the box number and calculates the line number of each record by adding the line length of each record to its line number to obtain the line number of the succeeding record; main chains start at the first line and each branch chain starts at the line set in the field FCPOC for that chain. Block 458 also stores the page, box, column and line numbers in the record. Block 460 thereafter writes the record to the output tape, and the subroutine returns to test 452 to repeat the loop for each succeeding record until all available records are processed.

After the OUTSUB subroutine is completed for all of the records of the cluster, the program returns to the beginning of the pass to process the next main chain and start the development of the next cluster. When all of the chains have been processed into clusters, Pass II is completed and Pass-III is started, in the manner described above. Various modifications may be made in Pass III as indicated above and also as discussed below. For example, a main chain may terminate before the end of a column leaving room for a small main chain thereafter; the documentation program is readily adapted to recognize this condition and to utilize the space efficiently by inserting the small main chain in the available main column space.

The flow chart documentation of a computer program in the form of clusters of chains (main and branch chains in main and auxiliary columns) enhances the two-dimensional character of the chart. It tends to supply the user with a greater amount of information about branches from the main-flow, since only one of four columns on a page is devoted to the main chain of a cluster, and the remaining columns display any of the branch chains that occur. Moreover, the sub-branching is also displayed, and as much sub-branching can be displayed as space permits; thereby, various loops and processing interrelationships and complexities in the documented program tend to be presented to the reader as he reviews each page of the chart. The main chain continues from column to column where branch chains do not occur or do not fit. The main chain also continues from page to page where it is of any substantial length, so that the main flow can be followed by flipping successive pages of the chart. But the intricacies of the documented program at any stage thereof generally occur at conditional transfer operations and they tend to be presented on a page displaying the cluster of branches from the main flow.

Various modifications may be made in the control system of this invention depending on the size and type of memory facilities provided by the computer. For example, where the computer has a large random-access memory (such as a core memory), the processing of the data may be made more efficient in the operation that requires searching for branch chains in the input tape during Phase II. The control system, as described above, only searches for those branch chains that do not exceed a specified length, say, not greater than one column. This search can be reduced by placing all potential branch chains (those which are less than the maximum length) on a separate magnetic tape during the Phase I operation; since these short chains are the only ones that may not be processed in their natural sequence on the input tape. The chain length, of course, is a variable that cannot be preset, and a special memory storage of the records making up a chain would be necessary in order to determine whether it was small enough to be a poten-

tial branch chain or not. Such a memory storage of records that would form one page column is provided and serves as a standard measuring unit for all chains that are developed. Those that fit within this memory area of a column length are placed in sequence on a separate output tape and serve as a search tape for potential branch chains. All other chains go on to another output tape in the usual fashion described above, and they are treated as main chains in the operation of the program. Such an arrangement would not result in any loss of processing time during Pass I, but would result in a substantial saving of time during Pass II, since the potential branch chains would be segregated and more readily searched, and the searching would not run through those chains that could not be branch chains. The order of chain processing is maintained and set forth in the chain table, which would indicate the tape that a particular chain was on. Thus, in picking up successive chains to be treated as main chains, either of the two tapes would function as a source, with the location of a particular chain on one of the two tapes being set forth in a chain table.

The searching of chains can also be reduced by filling the memory area that is available with as many potential branch chains as possible to eliminate the serial searching of these chains on the input tape. Another technique to reduce search time is to search for several chains at one time on the input tape with the search performed in the sequence of appearance of the chains. For example, the branch chains for the right auxiliary column can all be pulled out in one search when that auxiliary column is to be processed; thereby rewinding of the main tape separately for each branch chain would be avoided.

For computers having random-access disc or drum storage devices, the searching may be eliminated since the branch chain can be identified by its address on the disc or drum and retrieved directly in accordance with that address, which would be stored in the chain table instead of the sequence number.

Where all of the branch chains are separated initially, these chains may be preliminarily analyzed to determine whether a DECISION in a subsequent main chain (or auxiliary chain) refers to a tag in the branch chain. Thereby, it would be possible to pull such branch chains and allocate them to the subsequent DECISIONS as well as to preceding ones.

This invention is not restricted to the processing of small branch chains that fit in the remaining space in a column or on a flow chart page. The branch chains can be broken by pseudo-connectors (similarly as the main chains), and the Chain Table entry is used to identify the portion of the chain that has been processed and the portion that remains to be processed. Thereby, the remainder of the branch chain is picked up subsequently in the processing operation and displayed on an appropriate page of the chart identified in the pseudo-connector. Such a procedure enhances the two-dimensional character of the flow chart that is produced, since a large number of branch chains tends to be displayed on the same page as the associated DECISIONS. Even though only portions of the branch chains are shown on a page, the reader of the flow chart is generally given sufficient information to identify the type of operation performed in the branch chain and thereby given in a single page a greater amount of information about the interrelationships of the program both in the main flow and the various associated branches.

Where it is found desirable, the control system may be arranged to duplicate any or all branch chains in subsequent portions of the flow chart where the same branch chain was developed as a transfer from a preceding DECISION. Thereby, the user's reading of the flow chart is made more convenient, in that he does not have to move to different pages of the chart to identify the branch chains. In such a system, the second and subsequent dis-

plays of a branch chain are identified as duplicates, and the chain table maintains a record of first and subsequent displays of such branch chains. Such an arrangement may be a selective one under the control of the programmer in that a special code symbol may be provided to be inserted at the discretion of the programmer where he feels it desirable for branch chains to be duplicated. That is, each DECISION which would carry branch chains of special significance could be marked to present the branch chain as a duplicate (and in as much detail as desired). The program then operates with such a control signal to display the branch chain even though it may otherwise not be displayed at that particular location on the chart.

For small computers, those with memory sizes of appreciably less than 32,000 characters (e.g., 8,000) additional tape units (or random-access disc or drum storage) can be substituted to compensate for the smaller core memory. In Phase I, an additional tape is used to record the Chain Table entries. The Tag Table need not be developed during Phase I, but rather during Phase II; the tag information is extracted from the symbol record tape by means of an extra pass through all of the records (after the page, box, column and line number assignments are completed). In addition, the record for each symbol includes additional information, such as whether the symbol terminates a chain or not, and the format of the symbol to be printed stores with all of the printed elements at their various locations set forth in the record.

In Phase II, an additional tape is used to record the Tag Table entries, as indicated above; a tape is used to record the second-left auxiliary chains, if any, and a tape to record the main and any right auxiliary chains. Additional information in the form of a control record is added to the second-left auxiliary tape, which record sets forth the information pertaining to the number of columns required for each cluster and the distribution of the cluster over the different columns. This control record precedes each group of second-left auxiliary column records and defines the individual cluster. During Phase II, clusters are stored on the tapes prior to being assigned page, box and line numbers. The input to Phase II is the symbol record tape and the Chain Table tape from Phase I. During Phase II it has been found convenient to develop all cross-references in the form of a table which can be then utilized in the final tape that is constructed representing successive pages.

In Phase III, four tapes are provided to record each column of page on a separate tape, and a tape to describe the actual vertical and horizontal lines that are to be drawn on a page; the latter is placed on one of the column tapes preceding the associated column data, or it may be placed on a fifth tape. The small memory may only be large enough to hold a single record for each of the four columns and the matrix of required lines for the whole page. The page image is segmented into columns, and the columns and the matrix of line connections are merged by page and line during the printing process. Several passes of the data are required for drawing line connections on a page and for printing the page. Each conditional and unconditional branch is assigned page, box, column and line numbers during the Phase II allocation, so that all destination points are established for Phase III.

The first pass of Phase III reads the symbol record tape for a page and writes out a "line coordinate matrix" based on the destinations of symbol records for that page. The size of the matrix in characters, is the column length in lines multiplied by the number of vertical line lanes. For instance, if a vertical lane can be drawn on the left or right of a column and if each column is 150 lines long, and if there are 4 columns to a page, then a 1200 character matrix is required. A matrix character may represent any one of several "horizontal" or "point" conditions. For instances: it may specify a blank, an up arrow, a down arrow, a horizontal line pointing to the left, a

horizontal line pointing to the right, a vertical line and others. This "matrix" describes line drawings on a page and is formed by examining each decision or branch record on the symbol record tape. Based on (1) the position of the decision or branch record and (2) its destination point, then (3) the matrix can be examined and the line drawings plotted within the matrix. For instance, if a decision in column 1, line 25 of page 1 is to branch to column 2, line 29 of page 1, the broken line to the right, downward in the lane, and to the right may be drawn and can be described by five characters stored for lane 2, as follows: one character (say, coded "A") for the horizontal line between the DECISION branch-point and the lane (formed of 3 dashes), three characters (each coded "I") for the three vertical lines in the lane running downward, and one character (coded "B") for the horizontal line from the lane to column 2 (formed of 5 dashes with an arrowhead). The number of combinations possible for any point in the matrix can be represented by one character (in different embodiments about 6 to 12 possibilities exist). In the above example, the characters A, I, I, I, B would be stored in the matrix that represents lines 25 to 29 of lane 2, of column 1, of page 1. As the records are read in page by page, they are written out on 4 tapes; each tape containing one column. The "line matrices" for all pages are written on a separate tape or they can be placed on one of the column tapes, with each page's matrix written before the column record for that page. The next pass of Phase-III reads in the matrix tape, reads in each column of the page, creates a line for output based on the matrix and the column information (as described above) and prints the output.

Another change resulting from the design of a system for a small computer is that Phase IV (the development of the cross-reference listing), is produced after Phase II and prior to the Phase III printing of the flow chart page, for convenience in developing cross-reference data for display on the flow chart.

The particular embodiment of the invention described above is directed to input data in the form of an assembly language program, a form that is quite commonly used. This invention is also applicable to programs written in a "machine-independent" language, i.e., a language which is not limited to a particular construction of a machine nor to its particular body of instructions. The "instructions" of a machine independent language are macro-instructions or statements that can be translated as a group of machine instructions. Where a program consists of such macro-statements presented in accordance with a consistent convention, they can be interpreted by a flow-chart control program constructed in accordance with this invention to produce a two-dimensional flow chart that is properly representative.

The documentation system of this invention can process higher-level languages such as COBOL, FORTRAN, JOVIAL and other languages by analyzing the source statement input. The analysis may be performed prior to Phase-I, say Phase-O, or performed concurrently with Phase-I. The source statements are analyzed without the programmer writing the special codes that are currently required where the input is in an assembly language input (e.g., the embodiment of FIG. 8) or the special codes may be used where desired as an alternative in each case and also to delete or combine various statements of the program. The source statements are analyzed to (1) determine the special code (i.e., the F/C symbol; e.g., a subroutine, a process, an unconditional transfer, a conditional transfer), and (2) determine the Comments to appear in the flow chart symbol. Depending on the amount of analysis desired, the Comment produced (1) may be dependent on the procedure statement itself or, (2) may be dependent on the procedure statement plus an analysis of the nouns (e.g., data fields) and their associated descriptors (e.g., file descriptions). In any event, the flow charting of this system is independent

of the language used as input. For example, preprocessors to Phase-I analyze the higher-level language and then produce input to Phase-I of the system.

The source statements are analyzed to determine the corresponding special code or the flow chart symbol or symbols which are to be used to represent those statements (i.e., symbols such as subroutine, process, unconditional transfer or conditional transfer, or suitable modifications of those basic forms. This statement analysis may incorporate a well-developed technique used to translate higher level languages; that is, the translators for these languages (e.g., compilers or interpreters) provide well-known techniques for translating each language statement (operation) or data field into a symbol which would properly represent the operation and its relationship to other operations (e.g., tags) that may be involved in transfer operations. Thus, the state of the art is such as to permit the development of program processing for analyzing the statements of such languages into a form suitable for this documentation system to operate. The Comment field may vary depending on the amount of analysis of the language that may be desired and the amount of detail desired in the Comment field of the F/C symbol; thus, the Comment produced may be simply dependent on the procedure statement itself (e.g., a simple repetition of the verbal or algebraic statement making up the procedure statement) or it may be dependent upon the procedure statement together with an analysis of the nouns (e.g., the data fields) and their associated descriptors (e.g., the file descriptions for those fields). Suitable techniques for this purpose are likewise well-developed, and appropriate forms exist for different compilers.

Where the language statements involve complex logical conditions, standard techniques may be used to present the involved statement as the Comment itself. Alternatively, known compiler techniques may be used to break down the involved logical statement into its simple components so as to present each one as a conditional transfer (DECISION), whereby a logical sequence of conditional transfers is developed in the flow chart rather than the single involved statement.

By means of this documentation system, a higher level language program may be documented in a flow chart, and various ones of the features described above may be employed. That is, F/C symbols for the process blocks may vary in size with the requirements of the Comment field; chains of logical sequences may be developed in the manner described above with each chain terminating with an unconditional transfer; and conditional transfers and associated branch logic may be analyzed to develop a two-dimensional chart showing the branch chains from the DECISION symbols. Employing the technique described above for presenting a sequence of symbols on a flow chart as called for by the program to be documented, and upon reaching a conditional transfer, the branch chains are analyzed to determine the suitability of their being displayed in adjacent columns. Likewise, the cluster techniques would also be applicable so that sub-branch chains may be shown in two-dimensional configuration.

Thus, this invention provides mechanisms for automatically producing flow charts by means of a data processor operating on coded signals representing the instructions of a computer program to be documented. The mechanism for developing a two-dimensional flow chart pulls branch sequences of symbols for display on the same page as the main flow sequence from which the branch takes place. It processes the branch logic out of the order in which the logic appears in the original program in order to show the branch on the same page as the associated main flow logic. This mechanism involves the development of symbols for individual instructions or groups of instructions, and then the development of sequences of symbols and chains from the individual symbols. The mechanism further develops columns of symbols from

the main flow symbols and from the branch chains and assembles the columns into flow chart pages of related symbols, with the relationships being shown by connecting lines or references to the page and box numbers of the symbols. A mechanism also pulls sub-branch chains that stem from branch chains, and displays them on the same page. The cluster mechanism uses the above mechanisms and, in addition, forms each page of flow chart from one or more chain clusters, where a cluster is developed as a column of main flow logic and as many branch and sub-branch chains associated therewith as may fit on the page. In one form of the invention, the main flow symbols are successively allocated along the column until a DECISION is processed, and then the branch chain therefrom is processed by determining whether it fits within the adjacent column. The allocation of the branch chain may take place before the allocation of the main column is completed, or it may take place after the main column is allocated. The allocation of symbols of a branch chain may similarly lead to a sub-branch chain upon reaching a DECISION in the branch chain, and the suitability of fit of the sub-branched chain can be determined at that point.

This invention also furnishes mechanisms for editing unnecessary details from the original program and to combine a group of instructions into a single symbol. Thereby, it permits a programmer to edit, simplify and explain the program and the various portions and sequences thereof by means of an informative flow chart without the labor of drawing the chart or of laying out the sections thereof. A mechanism scans the various parts of the COMMENTS field of each input instruction and extracts the pertinent parts as required: the explanatory COMMENTS; the code, if any, for the type of symbol to be displayed; the destination tags, if any, carried for branch and transfer instructions; and the code for DECISION labels, or the labels themselves. By means of the symbol codes, detailed program instructions that are unnecessary for an understanding of the program may be deleted (either by the absence of a code, or by applying a delete code thereto), and several instructions may be combined and displayed as a single symbol, which more clearly sets forth the overall function of the detailed instructions. For example, a group of 20 individual instructions of the program to be documented may be displayed by a single DECISION symbol fully representing the overall function of those 20 instructions, though none of them may be a DECISION instruction, or several of them may be subsidiary DECISIONS. The essential functions of the program, by this mechanism, are set forth in the flow chart with as much detail as the programmer may desire to show. The significance of the processing details may be incorporated in statements of COMMENTS, which can be set forth in a NOTE or TEXT. The COMMENTS field also permits the carrying of labels for DECISION symbols, since the branch conditions may not be readily apparent from the details of instructions that are edited. Likewise, the COMMENTS field serves as a vehicle for destination tags for branch and transfer instructions, since these tags may not be always available from the operands due to the editing process. The COMMENTS of explanation of the program may be as long as desired, and a mechanism appends continuation COMMENTS to preceding records and edits these COMMENTS so that variable size symbols for PROCESS, NOTE and TEXT can be drawn and set forth in the flow chart.

This invention may also be used to interpret assembly language programs without the use of special codes for the flow chart symbols. The instructions themselves may be interpreted to develop process, unconditional transfer, and branch instruction symbols. In addition, the commentary customarily provided in an assembly language program may be extracted to present NOTES and the contents of PROCESS symbols. Where the symbol codes

are not used, the resulting flow chart may contain a good bit of detail that is not ordinarily required; however, some editing of the program can be obtained. For example, where a DECISION is formed by two or more individual instructions such as "compare" and "conditional transfer," appropriate techniques may be used for combining those instructions as a single flow chart symbol. Similarly, where a long string of PROCESS instructions occur, the individual instructions may be set forth in one PROCESS box with suitable separations between the individual instructions; thereby, a great deal of flow chart space is not wasted on the separations and connecting lines between PROCESS symbols. In addition, a programmer versed in a particular assembly language program may perform a small amount of editing by marking certain instructions, or groups thereof, with a "delete" code symbol to avoid unnecessary detail in the final flow chart.

A mechanism is also provided to illustrate on the flow chart the allocation data (page and box number) of all branch and transfer instructions which are the origin symbols for a particular tagged symbol on the chart. Thereby, the chart furnishes cross references to all of the originating points from which a particular branch or other chain stems. The cross reference list furnishes a full listing of such originating points. Connecting lines can be drawn between symbols on the same page if conflicts do not exist; the program mechanisms attempt to draw lines in the available lanes starting from one coordinate point of a page and attempting to go to the other. If such connecting lines cannot be drawn, then cross references are set forth on the chart.

Any of various types of output devices may be used to develop a record of the flow chart. The record may be a printed page made by a line-at-a-time printer or a digital plotter printer, or a momentary display record or a permanent photographic record made by a cathode ray tube or similar display device. The record may also be a magnetic tape recording of the corresponding page format.

The page format may take various forms: for example, the arrays of main flow and branch chain symbols may be set forth in parallel columns (or rows) interrelated in a two-dimensional chart by means of connecting lines or cross references; alternatively, main flow columns and branch rows (and sub-branch columns) may be used to develop the two-dimensional chart, as may any other suitable arrangement of symbol arrays. This invention is not limited to a fixed format of a columnar page; for example, with a cathode ray tube display, the magnification may be varied to permit different sizes of columns and individual symbols. The magnification for different sections of a column or of a page can be varied so that a low magnification can be used to fit a long chain in a column, and the magnification can be varied to determine the spacing between the symbols or columns to achieve the most suitable presentation for an individual page.

As described above, this invention may be used to develop flow charts from input programs written in any desired program language, including assembly and higher order languages. The invention is also adaptable to various types of computers including those having large and small memory capacities.

As previously noted, the stored-program embodiment of this flowcharting invention described above is preferred in that it is comparatively less expensive to construct than a fixed-program or hardware embodiment in the present state of the art. A stored-program processor such as that of FIGS. 8/1-8/30, or the more specifically detailed RCA 501 program noted below, enables one to modify, enlarge or simplify the system or any part thereof without rewiring the circuitry and changing any other hardware portion of the system.

The flowcharting processor of FIGS. 8/1-8/30 is not dependent on any particular form of computer. It may be used as the basis for providing a stored program for

any of a number of specific general-purpose computers that are available, and the RCA 501 program set forth below is one such example; other computers for which such programs have been provided include the IBM 360, 1401 and 7090, Honeywell 200 and the RCA Spectra.

A general-purpose computer, as an elementary information-transformation machine or apparatus, has a built-in capability limited to the execution of basic instructions such as add, subtract, compare, branch, etc. The stored-program embodiment of this invention converts the general-purpose computer into a special-purpose or extended machine having a unique operation sequence or process. Thereby, the programmed computer takes on the character of the flowcharting processor (program) that controls and directs the operation of the computer's hardware processor. For example, the flowcharting program for the RCA 501 set forth below converts the RCA 501 computer into a flowcharting computer machine during the time it is controlling and directing the computer. In that general-purpose computer, and generally in others, the logic and control circuits for performing the various instructions of the machine are time shared. The control signals of a stored program embodiment of this invention specify the particular instruction or instruction combination to be carried out at each instant in a specific and interrelated sequence. Thus the program's signal combinations physically initiate the operation of various ones of the computer's logic and control circuits, and direct the activation and deactivation of the computer's circuits and devices in certain sequences and relationships physically determined by the signal combinations of the program.

The interrelations of the circuits, their operations, and the control signals of the stored flowcharting program retermine the unique character of the computer as a flowcharting machine when it is so programmed. The flowcharting program acts as a control mechanism for the general-purpose computer to establish the configuration of machine operations that form the process of this invention. The flowcharting program as a control mechanism also determines a particular machine configuration, which is uniquely established while that program is in control; that is, the aforementioned mechanisms for automatically producing flow charts are established by various sequences of particular instructions and by various combinations of those sequences. The machine system for making two-dimensional flow charts incorporates, for example, the mechanism for pulling branch sequences of symbols for display on the same page as the main flow sequence; the mechanisms for developing symbols, and the sequences of the symbols; the allocation mechanisms; the mechanism for pulling the sub-branch chains; the cluster mechanisms; the editing mechanisms; the layout mechanisms. These mechanisms are combined to form a special-purpose mechane, which in an illustrated embodiment is a stored-program controlled general-purpose computer.

This invention may also be embodied in various forms of fixed or wired program embodiments. For example, a program form of this invention, similar to the RCA 501 program noted below, may be established in a read-only type of memory for use in those computers employing such memories for the control program. This "firmware" embodiment of the invention is constructed and operates in the same way, in all material respects, as the stored-program embodiment described above; the term "firmware" indicates the relatively permanent character of the sequence control formed by the program in a read-only memory as contrasted to the "software" embodiment of the program temporarily written in a read-write memory.

Another form of construction of the invention may employ a read-only memory to establish a macro-instruction embodiment of the invention; the physical removal of one read-only memory with its program and the replacement by another changes the computer's control

NAME	AUTOFLOW	A	M	B	INDEX	4000/501	SEG #1	DATE	111665	PAGE 004
LSM	OP				OP	A ADDRESS	N	P ADDRESS	CC3000	COMMENTS
003000	71	004170	00	000000	SEG 01					HOUSEKEEPING
003010					TC	START				BEGIN PROGRAM
					DAR	5 0/0				OUTPUT AREA FOR
					000050	SLC	004571			PASS 1, 1/P AND O/P
					R/K					FOR PASSES 2 AND 3
003011					000060	LINES	DAR 3			# LINES ON CHART
003014					000080	ADDR	DAR 4			LOCATION OF BOX
003020					000090	BOXNO	DAR 4			BOX ADDRESS
003024					000100	CODE	DAR 1			PROCESS CODE
003025					000110	TAG	DAR 7			TAG
003034					000120	LUTC	DAR 1			LEFT CONNECTION
003035					000130	LMT	DAR 7			LEFT DEST TAG
003044					000140	RHTC	DAR 1			RIGHT CONNECTION
003054					000150	RMT	DAR 7			RIGHT TEST TAG
003060					000160	OPSD	DAR 4			SEG # OF REC
003064					000170	ORLTH	DAR 4			LENGTH OF RCD
003075					000180	COMU	DAR 456			COMMENTS FIELD
004001					000190	TEMP1	DAR 3,1			TEMPORARY
004004					000200	TEMP2	DAR 3,1			STORAGE AREAS
004014	01	000034	01	010101	000210	TEMP3	DAR 7			LENGTH OF ENTRY
					000220	TENTL	DAC 000034			IN TAG TABLE
					000230	TENTLX	R/K			LNTM#1 OF TABLE ENTRY
					000240	TENTLX	EOL 000033			# TAGS IN TABLE
					000250	NUMTAG	DAC 000000			CURRENT MSC TABLE
					000260	TLOC	DAC 000000			MSC TAG TABLE
					000270	TBIN	DAC 11AR			# OF CHAINS
					000280	NUMCH	DAC 000000			CURR MSC CH TABLE
					000290	CHLOC	DAC 000000			MSC CHAIN TABLE
					000300	CHIN	DAC 000000			SEG # COUNTER
					000310	SEGNO	DAC 000000			INPUT TAPE
					000320	TT1	DAR 1			WORK TAPE 1
					000330	TT2	DAR 1			WORK TAPE 2
					000340	TT3	DAR 1			WORK TAPE 3
					000350	TT4	DAR 1			
					000360	LSTCO	CON 1			
					000370	OPCOM	DAC 000000			MSC PLUS ONE OF
					000380	ADPAR	R/K			O/P COMM AREA
					000400	ALMT	DAC 000000			ADDRESSES OF
					000410	ARMY	DAC RMT			TAG AREAS IN
					000420	ATA0	DAC TAG			OUTPUT AREAS
					000430	AADDR	DAC ADDR			# CHAR IN COLUMN
					000440	EPCON	DAC 000152			PASS ONE
					000450	TITLE	DAR 30			START OF PROGRAM
					000460	INDEX	DAR 6			REWIND INPUT
					000470	R/K				IF OFF-LINE OUTPUT
					000480	START	OCT 001210			REWIND PRINT TAPE
					000490	START	OCT 001210			RWD PASS 1 O/P
					000500		R/K			SET UP INPUT TRUNK
					000510		TC # 3			N/T 1
					000520		OCT 001240			N/T 2
					000530		RWD			
					000540		OCT 001220			
					000550		RWD			
					000560	HEX	OCT 001210			
					000570		OCT 001220			
					000580		OCT 001230			

NAME	AUTCFLOW	MSW	OP	A	N	B	SEG #	TAG	INDEX	4000301	N	SEG OF	DATE	111665	PAGE	COMMENTS
004310	22	001240	00	004057					OP	A ADDRESS	TT4					C/P IF OFF LINE
004320	75	000000	00	010000					CSG	010000						CLOSE GATE
004330	71	004350	00	400000					TC	RFX2						B.P. 5 ON
004340	71	004750	00	000000					TC	RFT						READ PAPER TAPE
004350	22	004054	00	004365					RMK							END FILE NOEOD, YES
004360	14	023200	00	000000					OCT	TT1						IF EFF GO TO EOD
004370	72	005426	00	000000					LRF	IPAR1						RECORD OF PROGRAM
004400	43	023200	00	023200					SC	IPAR1						HAS AN ED READ
004410	61	004430	00	004430					CTC	NEX1						
004420	71	004640	00	000000					TC	NEX2						
004430	71	004470	00	400000					TC	NEX4						
004440	72	023521	00	000000					SET	IPAR2						
004450	43	023225	00	023234					SC	IPAR1,2						
004460	61	004610	00	004560					CTC	NEX6						
004470	22	005427	00	023235					OCT	NEX5						
004500	12	023225	00	770000					LW	IPAR1,2						
004510	72	001167	00	600000					RIS	STAP1						
004520	24	004540	00	004557					SET	0:1167						
004530	71	000340	00	000000					STC	2						
004540	70	004170	00	000000					TC	000340						
004550	00	000002	00	005770					00	02						
004560	22	004054	00	004575					R/K							
004570	06	720001	00	000000					RMK							
004600	71	004350	00	000000					OCT	TT1						
004610	12	023510	00	770000					U:5	770001						
004620	12	004710	00	770000					TC	NEX2						
004630	71	004330	00	000000					R/K							
004640	71	005240	00	400000					LW	IPAR2						
004650	12	023510	00	770000					LW	NEX5						
004660	12	004710	00	770000					LW	NEX6						
004670	71	004750	00	000000					TC	RFT						
004700	71	004650	00	000000					TC	NEX1						
004710	576156466140540155566301								R/K							
004724	4556644543015655015035576463								TC	EJ						
004742									R/K							
004750	25	000240	00	005100					R/K							
004760	14	023510	00	770000					R/K							
004770	72	000223	00	600000					LRF	ICAF2						
005000	45	004761	00	004761					SC	010223						
005010	61	005030	00	005030					SC	RPT1						
005020	71	005240	00	000000					ETC	NEX2						
005030	72	005431	00	600000					TC	EJ						
									SET	R7574						

PAGE 005
 COMMENTS
 C/P IF OFF LINE P
 CLOSE GATE D
 B.P. 5 ON S
 READ PAPER TAPE D
 END FILE NOEOD, YES D
 IF EFF GO TO EOD P
 RECORD OF PROGRAM C
 HAS AN ED READ D

 B.P. 5 ON Z
 COMPARE PROGRAM NO. D
 TO PAPER TAPE C
 NEX5 P
 TYPE OUT NUMBER C
 OF PROGRAM C
 BRING IN SEG 2 P

BEGIN FLOW CHART N
 FOR THIS PROGRAM C
 STAP1 J
 UNWIND AN EF ON P
 INPUT FILE C
 ERROR ON P/T J
 TYPE ERROR MESSAGE P
 AND BAD RECCD C
 TRY AGAIN J
 ED ON INPUT SOURCE T
 B.P. 5 ON D
 ILLGICAL IF RPS N
 ON, SHOULD NOT C
 REACHED ON I/P C
 REAC REST OF P
 P/T AND TYPE OUT P
 AS ERRORS C
 EODJ C
 J

REAL PAPER TAPE T
 MESSAGE TEST FOR C
 VALIDITY AND EF C
 SET EXIT P
 REAL PAPER TAPE P
 HAS AN EF READ D

15 "MESSAGE FAIL" C

NAME	AUTOFLOW	HSM	CP	A	N	B	SEQ	R	TAG	INDEX	40000501	N	B	ADDRESS	DATE	111665	PAGE	006	COMMENTS
005740	43	023511	00	023511	00	023511	00	023511	SC	IPAR2,1				IPAR2,1					
005750	61	005110	00	005110	00	005110	00	005110	CTC	RPT2				RPT2					
005760	43	023522	00	023522	00	023522	00	023522	CTC	IPAR2,10				IPAR2,10					
005770	61	005110	00	005110	00	005110	00	005110	CTC	RPT2				RPT2					
005100	71	777777	00	000000	00	000000	00	000000	TC	777777				770000					TYPE OUT MESSAGE
005110	12	023510	00	770000	00	770000	00	770000	L	IFAR2				770000					AND EXPLANATION
005120	12	005150	00	770000	00	770000	00	770000	L	RPT3									HALT
005130	01	000000	00	000000	00	000000	00	000000	PE	S									BACK TO READ
005140	71	004760	00	000000	00	000000	00	000000	TC	RPT1									
005150									CON	16				INVALID MESSAGE,					
005172		5655540335013015444626204644160101							CON	18				PUT CORRECTION IN					
005214		57646301425661614442635056501505501							CON	17				REAPER, HIT START					
005235		614440434613501475063016283406163							CON	1				75					END OF JOB
005240	22	004054	00	005255	00	005255	00	005255	R	K				* 1,25					REMINI INPUT
005250	17	000000	00	000000	00	000000	00	000000	OC	T	1			8P0					IS OUTPUT OFF-LINE
005260	71	005330	00	010000	00	010000	00	010000	TC	EOJ1				* 1,25					WRIT ED TO O/P
005270	22	004057	00	005305	00	005305	00	005305	OC	T	4								
005300	12	005426	00	000000	00	000000	00	000000	L	W	#73#								REMINI O/P
005310	22	004057	00	005325	00	005325	00	005325	OC	T	4								
005320	17	000000	00	000000	00	000000	00	000000	R	MD									
005330	12	005432	00	770000	00	770000	00	770000	L	W	#END OF RUN/#			770000					TYPE END OF RUN
005340	76	000000	00	000000	00	000000	00	000000	ST	6FOR				END					HALT
023000									DAR	120#0				END					FRONT OF PAGE AREA
023200									R	#K				END					
023400									DAR	15000,100				END					
023510									S	L	C	L	A	Y	O	U	T		
024020									DAR	200#0				END					
024044									S	L	C	I	P	A	R	2			
005400									DAR	3200,0				END					
005424									S	L	C	L	A	Y	O	U	T		
005427									CON	22,0				END					
005430									S	0001				00000000000000000000					
005432									S	0002				# 73					
005450									S	0003				# 75					
									S	0004				# 7574					
									S	0005				END OF RUN/					
									S	CON 0,10				00000000000000000001					

NAME	AUTOFLOW	A	N	B	SEG #	TAG	INDEX	40000501	N	SEG 02	DATE	111665	PAGE	007
MSH	OP						OP	A	ADDRESS	START			COMMENTS	
004170					001430		CON 5							
004200					001440		CON 5							
005104	01	004201	01	010101	001450	MACOM	DAR 45040							CONVENTS WORK AREA
005110					001460	COMAD	DAC MACOM,1							MSC COMMENTS
005117					001470	MDTAG	DAR 7							RESULT OF PAREN
005124	01	000000	01	010101	001480	MDLAB	DAR 5							LINES IN CURR CHN
005130	01	000000	01	010101	001490	LNCNT	DAC 000000							STOPAGE FOR RME
005134	01	000000	01	010101	001500	RHECOM	DAC 000000							OF COMMENTS
005140	01	023200	01	010101	001510	RK								USED BY PASS 1
005144	01	023510	01	010101	001520	STORI	DAC IFAR1							READ ROUTINE
005260					001530	STOR2	DAC IPAR2							
005260					001540	IPCOM	DAR 76,0							
005267					001550	IP	RK INPUT AREAS FOR FIELD 05-MOVED FROM READ AREA BY INPUT SUB							
005270					001560	IP	DAR 2470							
005276					001570	IP	SLC IP							
005310					001580	IPTAG	DAR 7							TAG
005317					001590	IPCDEF	DAR 1							F/C CODE
005320					001600	IPLA	DAR 6							LINE NUMBER
005326					001610	IPAD	DAR 7							A ADDRESS
005340					001620		SLC IP, P 1							
005341					001630	WA	RK STORAGE AREAS CORRESPONDING TO ABOVE IN PUT AREAS							
005350					001640	WA	DAR 2470							
005360					001650	WATAG	SLC WA							
005374					001660	WATAG	DAR 7							
005375					001670	MACODF	DAR 1							TAG
005404					001680	ASL4	DAR 6							F/C CODE
005404					001690	ADWA	DAR 7							LINE NUMBER
005424					001700		SLC M-AR 1							A ADDRESS
005444					001710	WALAB	DAR 1							
005464					001720	MADCA	DAR 7							STORAGE LABEL CODE
005504					001730	MADCB	DAR 7							STORAGE FOR ADDRESSES
005524					001740	LBCD	CON 11,0							ON DECISIONS
005544					001750	LBCD1	CON 1,0							
005564					001760		CON 1,0							
005604					001770		CON 1,0							
					001780		CON 1,0							
					001790	LbTh	CON 16							
					001800		CON 16							
					001810		CON 16							
					001820		CON 16							
					001830		CON 16							
					001840		CON 16							
					001850		CON 16							
					001860		CON 16							
					001870		CON 16							
					001880		CON 16							
					001890		CON 16							

YES NO
ZNO YES
HIGH LOW
LLOW HIGH
OEQUALUNEQL
UNEQUAL
FOFF ON
COM OFF
#PLUS MINUS

DAR 4
R/K TABLE CONTAINING LABEL CODES AND ASSOCIATED LABELS
R/K FIRST CHARACTER OF EACH LINE IS CODE, REST IS LABELS
DAR 224,0
SLC LBTB
CON 16

NAME	AUTOFLOW	HSM	CP	A	N	B	SEG #	TAG	INDEX	4000/501	A	SEG OZ	DATE	111665	PAGE	008	COMMENTS
005624		365750A462C1545C5564620101010101					001900	001900	CON 16								
005644		1454555616257536462010101010101					001910	001910	CON 16								
005744		26446CA442537504647015336660101					001920	001920	SLC LBTAP 1								
005750		412434447515557626366					001960	001960	DAC LBT8								
005770		22 004054 00 614535					001970	001970	CON 11								
006003		22 004045 00 610011					001980	001980	OCT T11								
006010		75 000000 00 010000					001990	001990	OCT T12								
006020		22 004054 00 000151					002000	002000	R4K								
006030		14 023200 05 000000					002010	002010	CSG T11								
006040		72 004155 00 600000					002020	002020	OCT T11								
006050		24 023302 00 023337					002030	002030	SET TITLER								
006060		72 005465 00 600000					002040	002040	STC IPAR166								
006070		24 023225 00 023234					002050	002050	SET INDEX,R								
006100		14 023200 05 000000					002060	002060	STC IPAR121								
006110		75 000000 00 000000					002070	002070	LPF IPARI								
006120		71 014300 00 600000					002080	002080	TC I-P1								
006130		72 006137 00 600042					002090	002090	SET P*7								
006140		43 005267 00 005267					002100	002100	SC ILCODE								
006150		61 00617A 00 006170					002110	002110	CTC FPI								
006160		71 006120 00 000000					002120	002120	TC R01								
006170		25 005104 00 000153					002130	002130	R4K								
006200		25 004064 00 000173					002140	002140	R4K								
006210		32 003024 00 004060					002150	002150	R4K								
006220		36 003010 00 003773					002160	002160	R4K								
006230		22 017756 00 003010					002170	002170	TCY COMAD								
006240		22 017757 00 004200					002180	002180	TCT OFCOM								
006250		72 005337 00 600000					002190	002190	OCT CODE								
006260		36 005260 00 005307					002200	002200	SCT OPAR								
006270		36 005260 00 005307					002210	002210	OCT #7#R								
006300		71 015730 00 000000					002220	002220	OCT #7#R								
006310		71 014300 00 000000					002230	002230	SET W2#P								
006320		72 006327 00 600042					002240	002240	SCT I#								
006330		43 005267 00 005267					002250	002250	SCT I#								
006340		61 006500 00 006500					002260	002260	TC MCOM								
006350		71 015730 00 000000					002270	002270	TC INP1								
006360		72 005310 00 600000					002280	002280	SET #7								
006370		43 017760 00 017760					002290	002290	SC IPCODE								
006400		61 006310 00 606310					002300	002300	CTC FPA								
006410		72 005316 00 600000					002310	002310	TC MCOM								
006420		24 005260 00 005266					002320	002320	R4K								
006430		72 005325 00 600000					002330	002330	SET WATAG								
006440		34 005270 00 005275					002340	002340	SC #01#								
006450		34 005260 00 005266					002350	002350	CTC FP2								
006460		34 005270 00 005275					002360	002360	SET WATAG/R								
006470		71 006310 00 000000					002370	002370	STC I-P1#G								
006500		22 005317 00 003024					002380	002380	SET ASL,R								
							002390	002390	STC IPL,R								
							002400	002400	SCT IPL#								
							002410	002410	SCC IPL#								
							002420	002420	TC FP2								
							002430	002430	OCT WACODE								

PLT TAPE NUMBERS P
IN I/O INSTRUCTION C
STORE TITLE FROM P
PIC LINE C
CLOSE GATE C
SET TAPE SLCT
REAC NEXT RECORD P
STORE TITLE FROM P C
SECOND LINE C
STORE INDEX # F
REC AREA 1
DEF GATE
PICK UP FIRST RECORD S
IS CODE A C
Z A
NOM HAVE FIRST
RCD TO BE PFCO
ESSED IN I/P AREA C
START MAIN LOOP K
INITIALIZE FOR P
NEW RCD C
SAVE OLD CODE C
CLEAR O/P AREAS P
PUT IN ST MESS
MOVE TAG, OPERANDS, P
LINE #, COMMENTS C
FROM INPUT AREA C
TO WORK AREAS C
READ NEXT I/P RECORD S
CODE EQUAL TO C D
APPEND THIS COMMENT P
TO PACK OF PREVIOUS C
ONE IN COMMENT #/A C
IF TAG FIELD IS
BLANK, MOVE TAG OF P C
C LINE TO TAG C
WORK AREA C
MOVE LINE # OF P
C LINE TO WORK AREA C
READ NEW RECORD J
MOVE CODE TO O/P P
CODE

PAGE 008

SEG OZ DATE 111665

INDEX 4000/501

NAME AUTOFLOW

NAME	AUTOFLOW	MSM	OP	A	B	INDEX	4000/301	OP	A ADDRESS	SEG 02	DATE	111645	PAGE 009	COMMENTS
006510	44	004053	00	017763		OP	A ADDRESS	000001R#1	5	77777			ADD 1 TO SEQUENCE #	
006520	25	004050	00	003054		TCT	S#0A0#3	OP50					AND MOVE TO OUTPUT	
006530	22	017764	05	777777		TCT	S#0A0	OP50					77 TO RHE COMM FIELD	
006540	45	000153	00	017767		OP	R#7H						STORE RHE OF	
006550	25	000153	00	005130		TCT	S5	R#00002R#1					COMMENTS FIELD	
006560	72	005107	00	000000		TCT	S5	RHECOM					IS COMMENTS FIELD	
006570	42	000151	00	000153		SET	CMAD,R	ST	000153				COMPLETELY EMPTY	
006580	61	006640	00	006610		CTC	# 4						CREATE DUMMY	
006590	72	000004	00	000000		SET	CMAD,R	ST	000151				COMMENT AND	
006600	61	006640	00	006610		TCT	CMAD						LINE COMM TO AMS	
006610	72	000004	00	000000		SET	CMAD						IF THIS IS A R	
006620	24	017770	00	017773		SET	WACODE		1,1				CODES PUT TITLE	
006630	44	005130	00	017777		SET	WACODE		ST				IN TAG TABLE FOR	
006640	25	005107	00	000153		OP	WACODE						TABLE OF CONTENTS	
006650	22	005317	00	006661		LVS	#T4						IF T, BYPASS TAG	
006660	72	770001	00	600000		CTC	FP5						ROUTINE COMPLETELY	
006670	31	020000	00	020001		RMK							T CODE NOSPP5,YES# D	
006670	61	007240	00	007030		RMK							MOVE 77 TO "SC OF	
006670	25	004024	00	000133		RMK							TAG AREA IN TABLE	
006670	22	017764	03	000001		TCT	TLOC						MOVE 26 CHAR	
006670	72	000033	30	600000		RMK							OF TITLE TO TABLE	
006670	24	004201	00	004232		SET	TENTLX						IF TITLE LFSS	
006670	72	770001	00	600000		SET	MACM#1		MACM#26				THAT 26 CHAR	
006670	31	000023	33	000033		SET	770001						CHARACTERS MOVED	
006670	61	007000	00	007010		LVS	0-0#02						INCR RHE TABLE	
006670	34	000001	23	000033		CTC	# 1						IS THERE A TAG	
006670	44	004027	00	004017		SCT	000001						PRP = NO TAG	
006670	72	004023	00	017763		TCA	TLOC,R						MAKE TAG TABLE ENTRY	
006670	43	020002	00	020002		TCA	N#TAG,R						TO OUTPUT	
006670	61	007240	00	007060		CTC	FP5						ADD TO TAG COUNTER	
006670	72	003033	00	600000		RMK							MOVE TAG TO	
006670	24	005310	00	005316		SET	TAG,R						TAG TABLE	
006670	44	004023	00	017763		SET	WATIG						SEQ # TO TABLE	
006670	25	004024	00	000133		TCA	N#TAG,R						CLEAR ADDR FIELD	
006670	72	000007	30	600000		TCT	TLOC						MOVE ASSEMBLY LINE	
006670	24	005310	00	005316		SET	0-0#07						NUMBER TO TABLE	
006670	36	000014	33	000023		TCT	S#0						INCREMENT RHE TABLE	
006670	72	000027	30	600000		SET	0-0#14						CHAINSET BY COMMENT	
006670	24	005322	00	005325		SET	0-0#27						SWITCH #RCATFR#	
006670	22	005321	00	000010		CTC	WATIG						CODE	
006670	22	017761	00	000000		SET	ASL#2						# CODE	
006670	36	000030	33	000033		OP	ASL#5						Z	
006670	44	004027	00	004017		TCA	TLOC,R						Z	
006670	72	014030	00	004000		RMK							Z	
006670	72	000003	00	600000		SET	F#C						Z	
006670	43	005317	00	005317		CTC	#1						Z	
006670	61	007310	00	007310		CTC	#2						Z	
006670	71	010050	00	000000		TCT	J#0#E						Z	
006670	72	000004	00	600000		SET	R#1						Z	
006670	43	005317	00	005317		CTC	#2						Z	
006670	61	007350	00	007350		CTC	#2						Z	
006670	71	014040	00	000000		TCT	BR0#E						Z	

NAME	AUTOFLOW	HSP	OP	A	N	B	SEG #	TAG	INDEX	40000501	N	SEG 02	DATE	111665	PAGE 011	COMMENTS
010210	25	02027	00	003011			003600	HCODE	OP A ADDRESS			B ADDRESS				SET # LINES TO 10
010220	44	005127	00	020027			003610		TCT RC00010#1			R000010#1				ADD TO LINE CTR
010230	72	020030	00	600000			003620		TCA L.C.T,R			01				CODE OF NEXT RCD
010240	43	005267	00	005267			003630		SET "J"			IPCODE				EQUAL TO J
010250	61	010400	00	010300			003640		SC IPCODE			H1				REDUCE # LINES BY 3#
010260	45	003013	00	017777			003650		CYC HI			R000003#1				PUT J IN FIELD RMT
010270	22	005267	00	003045			003660		TCS LINES,R			RMT				FOR USE BY PASS 3
010300	45	000173	00	017767			003670		OPC IPCODE							ANY IS RME OF RCD
010310	22	005267	00	010321			003680	H1	TCS #7			R000002#1				IF NEXT CODE IS A OR P
010320	72	770001	00	600000			003690		OPC IPCODE			* 1,1				T, SET # LINES IN
010330	31	020000	00	020001			003700		SET 770001			01				CURRENT CHAIN TO A
010340	61	010350	00	010360			003710		LNS "BT"			-BT				HIGH # TO AVOID
010350	44	005127	00	020033			003720		CYC #1			#2				USE OF SEC, COLUMN C
010360	71	007750	00	000000			003730		TCA L.C.T,R			R001000#1				GO TO WRITE
010370	25	020023	00	003011			003740		TC FP6							PROCESS E CODE
010380	44	005127	00	020027			003750	ECODE	TCT RC00012#1			LINES				SET # LINES TO 10
010390	71	017060	00	000023			003760		TCA L.C.T,R			R000010#1				ADD TO LINE CTR
010400	45	000173	00	017767			003770		TC CHEAT							CHAIN TABLE ENTRY
010410	71	016730	00	000000			003780		TCS #7			R000002#1				ANY IS RME PCD
010420	45	000173	00	017767			003790		TC FP6							GO TO WRITE
010430	71	007750	00	000000			003800		TC FP6							PROCESS P CODE
010440	25	020037	00	003011			003810	PCODE	TCT R000005#1			LINES				SET # LINES TO 5
010450	71	017060	00	000023			003820		TC ELLIN			D00023				EDIT LINE OF COMMENT
010460	61	010510	00	010450			003830		TC P2			P1				END OF COMMENT FIELD
010470	44	005127	00	003013			003840		RMK			LINES,R				ADD TO LINE CNTR
010480	71	007750	00	000000			003850		TCA L.C.T,R							GO TO WRITE
010490	71	007750	00	000000			003860		TC FP6							*ORC TOO BIG TO
010500	71	007750	00	000000			003870		TC FP6							FIT ON LINE
010510	25	010457	00	010527			003880		RMK							CREATE ARTIFICIAL
010520	22	017760	05	777777			003890		TCT P1,7			* 1,7				WORD WITH SPACE
010530	71	010450	00	000000			003900	P2	OCY #1#			05 777777				
010540	25	020037	00	003011			003910		TC P1							PROCESS N CODE
010550	71	017060	00	000017			003920		RMK			LINES				SET # LINES TO 5
010560	61	010610	00	010550			003930	NCODE	TCT R00005#1			000017				EDIT LINE OF COMMENT
010570	44	005127	00	003013			003940		TC ELLIN			R1				END OF COMMENT FIELD
010580	71	007750	00	000000			003950		CYC #2							ADD TO LINE CNTR
010590	71	007750	00	000000			003960		RMK			LINES,R				GO TO WRITE
010600	44	005127	00	003013			003970		TCA L.C.T,R							*ORC TOO BIG TO
010610	25	010557	00	010627			003980		TC FP6							FIT ON LINE
010620	22	017760	05	777777			003990	N2	TCT N1,7			* 1,7				CREATE ARTIFICIAL
010630	71	010550	00	000000			004000		OCY #1#			05 777777				WORD WITH SPACE
010640	25	020043	00	003011			004010	BCODE	TC N1							PROCESS B CODE
010650	25	005130	00	000133			004020		RMK			LINES				ZERO TO # LINES
010660	45	000133	00	000153			004030		TCT RHFCOM			\$3				SET UP ADDRESS
010670	44	000173	00	000133			004040		TCS #3			\$3				PCR MOVE
010680	44	000173	00	000133			004050		TCA #7			\$3				
010690	44	000133	00	000153			004060		TCA #3			\$5				MOVE COMMENTS TO
010700	44	000133	00	000153			004070		SET 010000			70 \$1				OUTPUT AREA
010710	72	000000	70	600000			004080		STC 777777			59 010000				GO TO WRITE
010720	24	777777	53	600000			004090		TC FFA							PROCESS S CODE
010730	71	007750	00	000000			004100		RMK							PICK UP DESTINATION
010740	71	016210	00	000000			004110	SCORE	TC P-RFM			* 4				FROM # FIELD
010750	61	011610	00	011610			004120		CF # 4							USE A ADDRESS
010760	72	003043	50	600000			004130		SET L1,R			\$1				

NAME	AUTOFLOW	A	M	B	SEC #	TAG	1NDX	40000501	N	SEG 02	DATE	111665	PAGE 013
MSV	OP						OP	A ADDRESS	N	B ADDRESS			COMMENTS
011550	25	000223	00	000153	004760		TCT	85TA	85				AM5 TO *
011560	25	000153	00	000113	004770		TCT	85	81				AM1 TO *01
011570	44	000153	00	017763	004780		TCA	85	8000001#01				CLEAR ISS FROM FRONT
011600	22	017760	00	004200	004790		OCT	801#	MACOM				OF COMMENTS W/A
011610	45	000113	00	017763	004800		R#K						FIND FIRST SPACE
011620	72	000001	10	600000	004810		TCS	81	8000001#01				TO LEFT OF STAR
011630	43	017760	00	017760	004820		SC	801#					MOVE 5 CHARACTERS
011640	61	011610	00	011610	004830		SC	801#					TO LAB AREA IN O/P C
011650	72	000015	70	600000	004840		CTC	803					CLEAR LABEL FIELD IN P
011660	24	000002	11	000006	004850		SET	000015	70 ST				COMMENTS AREA
011670	34	000001	15	777777	004860		STC	000002	11 000006				CLEAR OUT ANY CRAP
011700	72	150001	00	600000	004870		SCC	000001	15 777777				IN CASE LABEL IS
011710	31	000011	77	000015	004880		LMS	000011	77 000015				LESS THAN 5 CHAR
011720	61	011730	00	011740	004890		CTC	81	8000015				SAVE AMI
011730	34	000001	27	000015	004900		SCC	000001	27 000015				REPLACE ISS AT FRONT
011740	25	000113	00	004001	004910		TCT	81	TEMP2				OF COMMENTS W/A
011750	22	017757	00	004200	004920		OCT	874#	MACOM				EXTRACT 1ST TAG AND
011760	71	016210	00	000000	004930		R#K						LABEL IN STAR FIELD
011770	72	003043	00	600000	004940		TC	P.REN					TAG GOES TO
012000	24	005110	00	005116	004950		SET	L#T#R	8T				OUTPUT AREA
012010	72	000003	70	600000	004960		STC	000003	80TAG,R				LABEL GOES TO
012020	24	005117	00	005123	005000		STC	H0LAB	70 ST				OUTPUT AREA
012030	71	016210	00	000000	005010		R#K		H0LAB,R				IF 10 TAG OR LABEL,
012040	72	003053	00	600000	005020		TC	P.REN					SPACES ARE MOVED
012050	24	005110	00	005116	005030		SET	L#T#R	8T				TO OUTPUT AREA
012060	72	000010	70	600000	005040		STC	00010	80TAG,R				EXTRACT 2ND TAG AND
012070	24	005117	00	005123	005050		SET	00010	70 ST				LABEL IN STAR FIELD
012100	43	005133	00	600000	005060		STC	H0LAB	8T				MOVE 2ND TAG AND
012110	43	000151	00	000153	005100		SET	R#COM#R	8T				LABEL TO C/P AREA
012120	61	012130	00	012170	005110		SC	000151	000153				IF ASTERISK FIELD
012130	25	004001	00	000113	005120		CTC	81	000153				WAS AT END OF
012140	25	000113	00	005130	005130		TCT	TEMP2	DI				COM'FIELD# ADJUST
012150	22	017764	01	000001	005140		OCT	87#	81				TEMP2 TO RHE COM FIELD
012160	25	005104	00	000153	005150		TCT	COM#D	81 000001				RHE OF FIELD TO
012170	44	000173	00	020053	005160		R#K	IF ASTERISK FIELD WA S	85				LEFT OF LABEL
012200	72	020013	00	600000	005170		TCA	87	800017#1				SENTS NEEDED
012210	43	003024	00	003024	005180		SC	CODE	8T				SET AM7 TO POINT P
012220	61	012240	00	012240	005190		CTC	82	800017#1				TO COMMENTS FIELD
012230	71	013420	00	000000	005200		TC	WCODE	82				N CODE
012240	71	017440	00	000000	005210		TC	S:50					C
012250	25	005130	00	000133	005220		R#K						Z
012260	45	000133	00	000153	005230		TCT	R#COM	83				SKIP ANY SPACES TO
012270	45	000133	00	020057	005240		TCS	83	83				RIGHT OF * FIELD
012300	61	012410	00	012310	005250		CTC	83	85				LESS THAN 13
012310	72	000000	70	600000	005260		R#K		800014#1				COMMENT CHARACTERS
012320	24	020060	00	020061	005270		CTC	D6	81				C
					005280		R#K						# CHAR IS LESS
					005290		R#K						N
					005300		R#K						TMA 13
					005310		R#K						PUT IN 2 ISS FOR
					005320		SET	000001	70 ST				P
					005330		STC	87474#	87474#				2 BLANK LINES

NAME	AUTOCFLD	HS	CP	A	S	R	SEQ #	TAG	INDEX	ADDRESS	4000:501	A	SEG #2	DATE	111685	COMMENTS
012330	44	060173	00	017767			005340	TCA \$7	OP. A ADDRESS							PAGE 014
012340	71	017060	00	000015			005350	TC ELLIN		000002#1						COMMENTS
012350	61	000240	00	000240			005360	TC ELLIN		000015						ADJUST AMT
012360	25	020067	00	003011			005370	RMK		000240						PICK-UP SINGLE LINE
012370	44	065127	00	003013			005380	TC L,CTR,R								P
012400	71	007750	00	000000			005400	TC EP6		LINE5,R						S
012410	71	017060	00	000007			005410	TC ELLIN								P
012420	61	012730	00	012520			005420	CTC D5		000707						C
012430	72	060001	70	600000			005430	RMK								P
012440	24	003103	07	000000			005440	SET 000001		70 \$T						C
012450	44	000173	00	017763			005450	STC 000015		07 000000						C
012460	72	000001	70	600000			005460	TCA \$7 00001#1		000001#1						C
012470	24	060103	07	000000			005480	TCA \$7		07 000000						P
012500	44	000173	00	017763			005490	TC D2								P
012510	71	012360	00	000000			005500	TC D2								C
012520	71	017060	00	000013			005510	TC ELLIN		000013						P
012530	61	012730	00	012550			005520	CTC D5		000013						C
012540	71	012460	00	000000			005530	TC D4A		000013						P
012550	71	017060	00	000013			005540	TC ELLIN		000013						S
012560	61	012730	00	012600			005550	CTC D6		000013						Z
012570	71	012460	00	000000			005560	TC DAA		000013						S
012600	71	012460	00	000013			005570	TC ELLIN		000013						Z
012610	61	012730	00	012630			005580	CTC D-		000013						D
012620	71	012360	00	000000			005590	TC D-		000013						Z
012630	71	017060	00	000007			005600	TC D7		000007						S
012640	61	012730	00	012660			005610	TC D4		000007						D
012650	71	012360	00	000000			005620	TC D7		000003						Z
012660	71	017060	00	000003			005630	TC ELLIN		000003						S
012670	61	012710	00	012710			005640	CTC D5		00						D
012700	71	012360	00	000000			005650	TC D2		00						J
012710	44	000173	00	017767			005670	RMK								T
012720	71	012360	00	000000			005680	TC A \$7		000002#1						C
012730	25	000243	00	000133			005700	TC D2								P
012740	25	777767	30	012757			005710	RMK								J
012750	22	017760	05	777777			005720	TC D2								C
012760	71	777760	30	000000			005730	RMK								P
012770	25	005744	00	000113			005740	TCT 95TP		00						P
013000	72	000000	10	600000			005750	TCT 777767		00						P
013010	43	005340	00	005340			005760	CTC D1P		05						P
013020	00	013400	00	013400			005770	TC 777760		00						C
013030	72	000015	70	600000			005780	RMK								F
013040	24	000001	11	000005			005790	CTC D1Z		01						T
013050	72	000010	70	600000			005800	SET 010000		10						P
013060	24	000006	11	000017			005810	SC MALAB		MALAB						C
013070	22	005340	00	013101			005820	CTC D1Z		D1Z						D
013100	72	770001	00	600000			005830	SET 000015		70 \$T						P
013110	31	005374	00	005374			005840	SET 000010		70 \$T						C
013120	61	013130	00	013270			005850	SET 000006		11 000017						C
013130	72	003043	00	600000			005860	CTC 000006		00						D

SHIFT LINES DOWN
BY ONE FOR BETTER
VERT ALIGNMENT
SET LINES TO 13
ADD TO LINE COUNTER
GO TO WRITE
EDIT LINE OF COMMENT
END OF COMMENT FIELD

EDIT LINE OF COMMENT
END OF COMMENT FIELD
EDIT LINE OF COMMENT
END OF COMMENT FIELD
EDIT LINE OF COMMENT
END OF COMMENT FIELD
EDIT LINE OF COMMENT
END OF COMMENT FIELD
EDIT LINE OF COMMENT
END OF COMMENT FIELD
EDIT LINE OF COMMENT
END OF COMMENT FIELD

NOT ENOUGH ROOM
IN THE DIAMOND
FORCE END OF FIELD
DROP REST OF COMMENT
LINE TOO SMALL FOR
EVEN ONE WORD
CREATE ARTIFICIAL
WORD AND TRY
AGAIN

USE LABEL CODE TABLE
API TO HBC OF TABLE
COMPARE LABEL CODE
AGAINST TABLE ENTRY

MOVE LABELS
FROM TABLE TO
OUTPUT AREA

THREE WAY DECISION

MOVE A AND P ADDRESS

NAME	AUTOFLOP	MSV	OP	A	B	SEC #	TAG	INDEX	40000501	N	SEC #2	DATE	111665	PAGE	OIS
013140	24	005341	00	005347		005920		OP A ADDRESS			B ADDRESS			COMMENTS	
013150	72	005353	00	000000		005930		STC WADCA			WADCA,R			STORAGE AREAS TO O/P	C
013160	24	005350	00	005356		005940		SET RMT/R			WADCB,R				
013170	72	002071	00	000000		005950	09	SET #1515#			WADCB,R				
013200	43	005345	00	003045		005960		SC L-T			LMT			IF ANY OF THE	P
013210	61	013230	00	013230		005970		CTC # 2			LMT # 2			DESTINATION TAGS	C
013220	34	005335	00	003043		005980		SCC LMT			LMT/R			HAVE E-Z CODE #	C
013230	43	005345	00	003045		005990		SC R-T			RMT # 2			CONVENTIONS CLEAR	C
013240	61	013260	00	013260		006000		CTC # 2			RMT/R			THEY TO SPACES	C
013250	34	005345	00	003053		006010		SCC R-T							
013260	71	012170	00	000000		006020		TC D1							
013270	72	020070	00	000000		006030		RMK							J
013300	43	005341	00	005341		006040	D10	SET #15#						A ADDRESS HAVE	D
013310	61	013350	00	013350		006050		SC WADCA			WADCA			ASTERISK NOTATION	C
013320	72	005343	00	000000		006060		CTC D11			D11				
013330	24	005350	00	005356		006070		SET LMT/R						USE B ADDRESS	P
013340	71	013170	00	000000		006080		STC WADCB			WADCB,R			AS DESTINATION	J
013350	72	005343	00	000000		006090		TC D9						USE A ADDRESS	C
013360	24	005341	00	005347		006100	D11	SET LMT/R			WADCA,R			AS DESTINATION	P
013370	71	013170	00	000000		006110		STC WADCA							
013400	44	000113	00	020077		006120	D12	TC D9						J A#1 TO NEXT ENTRY	P
013410	71	013300	00	000000		006140		TC D#						RECYCLE	J
013420	44	000173	00	020103		006150	NCODE	RMK						PROCESS W CODE	T
013430	72	000000	70	000000		006170		TCA #7						ADJUST A#7 TO RHE	P
013440	24	013470	00	013527		006180		RMK						OF O/P COMM FIELD	C
013450	22	020005	00	003024		006190		SET O#0#00						MOVE IN CONSTANT	P
013460	71	012360	00	000000		006200		STC S#CON1			S#CON4,R			FOR COMMENT	C
013470						006210		OCT #43#						CHANGE CODE TO D	P
013472						006220		RMK						PROCESS AS D CODE	N
013505						006230		CON D2			7474				J
013515						006240		CON 2							
013530	25	017777	00	003011		006250		S#CON1							
013540	22	004060	00	013551		006260		S#CON2							
013550	72	770001	00	000000		006270		S#CON3							
013560	61	020104	00	020110		006280		S#CON4							
013570	61	013600	00	013750		006290		CON 8							
013600	25	005130	00	000113		006300		CON 11							
013610	72	150001	00	000000		006310		CON 11							
013620	31	004200	01	000000		006320		CON 8							
013630	61	013660	00	013660		006330		CON 11							
013640	22	017764	02	000001		006340		CON 8							
013650	71	013610	00	000000		006350		CON 11							
013660	71	017060	00	000034		006360		CON 8							
013670	61	013660	00	013660		006370		CON 11							
013700	72	005133	00	000000		006380		CON 11							
013710	43	005151	00	000153		006390		CON 8							
013720	61	013730	00	013773		006400		CON 11							
013730	44	005127	00	003013		006410		CON 11							
013740	71	007750	00	000000		006420		CON 8							

COMMENTS
 STORAGE AREAS TO O/P
 IF ANY OF THE
 DESTINATION TAGS
 HAVE E-Z CODE #
 CONVENTIONS CLEAR
 THEY TO SPACES
 A ADDRESS HAVE
 ASTERISK NOTATION
 USE B ADDRESS
 AS DESTINATION
 USE A ADDRESS
 AS DESTINATION
 J A#1 TO NEXT ENTRY
 RECYCLE
 PROCESS W CODE
 ADJUST A#7 TO RHE
 OF O/P COMM FIELD
 MOVE IN CONSTANT
 FOR COMMENT
 CHANGE CODE TO D
 PROCESS AS D CODE

INDEX
 OP A ADDRESS
 STC WADCA
 SET RMT/R
 STC WADCB
 SET #1515#
 SC L-T
 CTC # 2
 SCC LMT
 SC R-T
 CTC # 2
 SCC R-T
 RMK
 TC D1
 SET #15#
 SC WADCA
 CTC D11
 SET LMT/R
 STC WADCB
 TC D9
 SET LMT/R
 STC WADCA
 TC D9
 TCA #1
 TC D#
 RMK
 TCA #7
 RMK
 SET O#0#00
 STC S#CON1
 OCT #43#
 RMK
 TC D2
 CON 2
 S#CON1
 S#CON2
 CON 11
 CON 8
 CON 11
 RMK
 TCT #00003#,1
 R#K
 OCT LSTCD
 SET 770001
 LMS #5147444163
 CTC # 1
 TCT R-ECOM
 SET 150001
 LMS #5147444163
 CTC # 1
 OCT #77
 TC #4
 TC EDL#R
 CTC #1
 SET R#FCOM#AS
 SC O 0151
 CTC # 1
 TCA L-C-T,R
 TC FFF
 RMK

TCODE
 06270
 06280
 06290
 06300
 06310
 06320
 06330
 06340
 06350
 06360
 06370
 06380
 06390
 06400
 06410
 06420
 06430
 06440
 06450
 06460
 06470

PROCESS T CODE
 SET # LINES TO J
 FOR VERT MARGIN
 IS PREVIOUS CODE
 J, H, E, O, T
 CHANGE ANY # P
 IN COMMENTS WORK C
 ARE/ TO 77 DCTAL C
 TO FAKE EDL#R S#R C
 INTO ENDING LINE AT C
 THAT POINT C
 EDIT LINE OF TEXT
 END FIELD INDIC SET
 REAL END OF
 COMMENT FIELD
 ADD TO LIVE COUNTER
 PREV CODE LCT JMBRT

NAME	AUTOFLO	MS	OP	A	B	SEQ #	TAG	INDEX	4000/501	SEG # 2	DATE	111665	OP	A ADDRESS	B ADDRESS	CODE	COMMENTS
013750	22	020010	00	003024		006400	T3	RK	4551				RK	4551			PAGE 016 COMMENTS CHANGE CODE TO PROCESS AS NOTE INSTEAD OF TEXT
013760	71	010540	00	000000		006500		RK	400002H,1				RK	400002H,1			J SET AM5 TO TEXT LINE IS5 TO LEFT OF FIELD AM7 TO NEXT LINE RECYCLE
013770	44	000153	00	017767		006510	T4	TC	400002H,1				TC	400002H,1			P REST SWITCH FPS IF RECORD IS N OR T CODE, SET INDICATOR IN CHAIN
014000	22	017757	05	777777		006530		TC	474				TC	474			C TABLE 30 PASS 3 KMD'S 1ST WCD OF CHAIN IS N OR Y
014110	44	000173	00	017767		006540		TC	T1				TC	T1			J END OF PASS 1 WRITE EF
014120	71	013660	00	000000		006550		RK	FIRST RECORD OF EA CH CHAIN				RK	FPS,5			P WRITE ED
014300	22	007247	00	007245		006570	FRC	OC	FPS,7				OC	FPS,5			P REWIND
014500	22	003074	00	014051		006580		OC	CODE				OC	1,1			P PUT 00 IN RMC PLUS 1 OF CHAIN TABLE BRING IN SEG 02
014600	31	020111	00	020112		006590		SET	770001				SET	1,5			P WRITE ED
014670	61	014100	00	007250		006610		LWS	1,7				LWS	000073			P REWIND
014100	25	004040	00	000133		006620		TCT	4LOC				TCT	1,5			P PUT 00 IN RMC PLUS 1 OF CHAIN TABLE BRING IN SEG 02
014110	22	020113	03	000004		006630		OC	405				OC	3			P WRITE ED
014120	71	007250	00	000000		006640		TC	FPSA				TC	1,5			P REWIND
014130	22	004055	00	014145		006650		RK	1,5				RK	1,5			P PUT 00 IN RMC PLUS 1 OF CHAIN TABLE BRING IN SEG 02
014140	12	014147	00	000072		006660	ENDP1	OC	T12				OC	000072			P WRITE ED
014150	22	004055	00	014165		006670		L4	1,7				L4	1,5			P REWIND
014160	12	014167	00	000073		006680		OC	T12				OC	000073			P PUT 00 IN RMC PLUS 1 OF CHAIN TABLE BRING IN SEG 02
014170	22	004055	00	014205		006690		L4	1,7				L4	1,5			P REWIND
014200	17	000000	00	000000		006710		RWD					RWD	33			P PUT 00 IN RMC PLUS 1 OF CHAIN TABLE BRING IN SEG 02
014210	25	004040	00	000133		006720		TCT	4LOC				TCT	1,5			P WRITE ED
014220	22	017741	03	000000		006730		OC	RWD				OC	3			P REWIND
014230	72	001167	00	600000		006740		RIS	START				RIS	1,5			P PUT 00 IN RMC PLUS 1 OF CHAIN TABLE BRING IN SEG 02
014240	24	014260	00	014277		006750		SET	001167				SET	1,5			P WRITE ED
014250	71	000340	00	000000		006760		SYC	2				SYC	1,5			P REWIND
014260	00	004170	00	000000		006770		TC	000340				TC	000073			P PUT 00 IN RMC PLUS 1 OF CHAIN TABLE BRING IN SEG 02
014270	00	000000	00	004600		006780		OO	START				OO	1,5			P WRITE ED
014280	00	000000	00	004600		006790		OO	03				OO	1,5			P REWIND
014300	25	00243	00	015330		006820		RK	STRASZ				RK	1,5			P PUT 00 IN RMC PLUS 1 OF CHAIN TABLE BRING IN SEG 02
014310	25	005140	00	000113		006830		RK	1,5				RK	1,5			P WRITE ED
014320	72	015470	00	000000		006840		RK	1,5				RK	1,5			P REWIND
014330	25	005140	00	014530		006850		RK	1,5				RK	1,5			P PUT 00 IN RMC PLUS 1 OF CHAIN TABLE BRING IN SEG 02
014340	25	005134	00	005140		006860		RK	1,5				RK	1,5			P WRITE ED
014350	25	014530	00	005134		006870		RK	1,5				RK	1,5			P REWIND
014360	25	005140	00	000113		006880		RK	1,5				RK	1,5			P PUT 00 IN RMC PLUS 1 OF CHAIN TABLE BRING IN SEG 02
014370	62	014370	00	014370		006890		RK	1,5				RK	1,5			P WRITE ED
014400	72	000000	10	600000		006920		SSH	10				SSH	10			P REWIND
014410	43	020114	00	020114		006930		SC	872H				SC	872H			P PUT 00 IN RMC PLUS 1 OF CHAIN TABLE BRING IN SEG 02
014420	61	014530	00	014530		006940		CTC	INP11				CTC	INP11			P WRITE ED
014430	22	010026	00	010025		006950		OC	FPS,5				OC	000000H,1			P REWIND
014440	45	004053	00	000043		006970		TC	SERVO,R				TC	FP4			P PUT 00 IN RMC PLUS 1 OF CHAIN TABLE BRING IN SEG 02
014450	61	006500	00	006500		006980		CTC	FP4				CTC	FP4			P WRITE ED
007000						006990		RIS	LOAD1				RIS	LOAD1			P PUT 00 IN RMC PLUS 1 OF CHAIN TABLE BRING IN SEG 02

NAME	AUTOFLOW	A	N	B	SEQ N	TAG	INDEX	4000R501	N	SEG C2	DATE	3116B5	PAGE 017	COMMENTS
MSN	OP						OP	A ADDRESS						
014460	72	001167	00	600000			SET	001167						
014470	24	014510	00	014527			STC	* 2						
014500	71	000340	00	000000			TC	000340						
014510	00	003000	00	000000			00	LOAD1						
014520	00	000001	11	004260			00	01						
014530	14	000000	00	000000			R/K							
014540	72	000011	10	600000			LRF	060C11						
014550	43	020002	00	020002			SET	060C11						
014560	01	014630	00	014670			SC	R2E*						
014570	72	008266	00	600000			CTC	I.P12A						
014600	24	000011	11	000017			SET	I.P12A						
014610	72	005275	00	600000			STC	000C11						
014620	24	000007	11	000007			SET	I.P12A						
014630	72	005304	00	600000			STC	000002						
014640	24	000075	11	000033			SET	I.P12A						
014650	36	005144	00	005257			STC	000C25						
014660	72	000C23	10	600000			SET	IPCOM,R						
014670	43	020115	00	020117			SC	R/K						
014700	61	015120	00	015120			CTC	I.P12C						
014710	72	005257	00	600000			SC	"R/K"						
014720	35	000102	11	000137			SET	IPCOM,R						
014730	61	014740	00	014750			SCR	00102						
014740	24	017760	00	017760			CTC	* 1						
014750	35	000055	11	000100			STC	R1H						
014760	61	014770	00	015000			SCR	000F55						
014770	24	017760	00	017760			CTC	* 1						
015000	35	000052	11	000053			SCR	R1H						
015010	61	015020	00	015030			SCR	000C52						
015020	24	017760	00	017760			CTC	* 1						
015030	35	000025	11	000050			STC	R1H						
015040	34	005260	00	005266			SCR	000C25						
015050	34	005276	00	005304			SCC	I.P12A						
015060	72	005257	00	600000			SET	IPCOM,R						
015070	43	017760	00	017760			SC	R1H						
015100	61	015150	00	015150			CTC	I.P12B						
015110	71	014310	00	000000			TC	I.P1A						
015120	72	005257	00	600000			R/K							
015130	35	000102	11	000137			SET	IPCOM,R						
015140	61	015150	00	014310			SCR	00102						
015150	72	005256	00	600000			CTC	* 1						
015160	43	020120	00	020121			SET	IPCOM,R=1						
015170	61	014310	00	014310			SC	R1C1A						
015200	22	005257	00	015211			CTC	I.P1A						
015210	72	770001	00	600000			SET	IPCOM,R						
015220	31	005750	00	005762			SET	770001						
015230	61	015240	00	015360			L/S	VALCO						
015240	22	005257	00	005267			CTC	* 1						
015250	34	005257	00	005257			SET	IPCOM,R						
015260	22	005267	00	015271			SCC	IPCOM,R						
015270	72	770001	00	600000			SET	IPCOM,R						
015300	31	020122	00	020123			SET	770001						
015310	61	015330	00	015340			L/S	"I,C"						
015320	22	020124	00	014321			CTC	I.P1X						
015330	71	777777	00	000000			CTC	R1H						
015340	22	020114	00	014320			TC	I.P1X						
015350	71	015330	00	000000			TC	I.P1X						

*NEXT J
 READ SOURCE RECORD P
 JS TAG FIELD BLANK D
 Z PICK UP TAG P
 FROM SOURCE INPUT C
 MOVE ASSY LINE # P
 TO STORAGE C
 MOVE A ADDRESS P
 TO STORAGE C
 IS THIS A R/K LINE D
 MOVE AND RIGHT P
 JUSTIFY ALL FIELDS C
 OF R/K LINE C
 WAS ENTIRE LINE D
 A NULL ONE C
 BACK FOR NEXT RCD J
 NOT REMARK LINE N
 MOVE COMMENTS P
 AND RIGHT JUSTIFY C
 LAST CHARACTER D
 SEPARATED BY C
 7 TO SPACES C
 LAST CHARACTER D
 A VALID FLOW CHART C
 CODE C
 MOVE TO CODE AREA F
 CLEAR LAST CHARACTER P
 TEST CODE FOR D
 D OF C C
 D'I.P1X,C01.P13A,NO# C
 SET SWITCH I.P1B P
 RESET SWITCH I.P1B E
 P F J

NAME	AUTOFLOW	MSW	OP	A	N	B	SEQ #	TAG	I-DEX	4000-501	N	SEG #2	DATE	111665	PAGE	CIB
015360	22	005257	00	015371			007540		OP	A ADDRESS						CON'TENTS
015370	72	770001	00	600000			007550	INP13	RK	ICOM,R						NOT FLOW CHART CODE
015400	31	005360	00	005372			007560		SET	770001						IS THIS A VALID
015410	61	015420	00	014310			007570		LMS	LICD						LABEL CODE
015420	22	005257	00	005340			007580		CTC	*						PRIVATE NEW RCD
015430	71	015640	00	600000			007600		TC	I-P15						STATE LABEL CODE
015440	22	000125	00	005267			007610		OC	T*						PICK UP ADDRESSES
015450	34	005257	00	005257			007620		SCP	ICOM,R						CHANGE CODE TO C
015460	71	015330	00	000000			007630		TC	I-P1X						CLEAR LABEL CODE
015470	25	000243	00	015630			007650	INP14	TCT	BSTP						PICK UP ADDRESSES IN
015500	72	020138	00	600000			007660		SET	*A TC CTC*						ABSENCE OF LABEL CODE
015510	43	000021	11	000023			007670		SC	010021						SET EXIT
015520	61	015540	00	015540			007680		CTC	* 2						IS CP CODE A CTC,
015530	71	015610	00	000000			007690		TC	I-P14A						TC, OR TALLY
015540	43	000021	11	000023			007700		SC	000021						
015550	61	015570	00	015570			007710		CTC	* 2						
015560	71	015610	00	000000			007720		TC	I-P14A						
015570	43	000021	11	000023			007730		SC	000021						
015600	61	015630	00	015630			007740	INP14	CTC	I-P14X						ASSUME Y CODE
015610	22	020137	00	005340			007750		OC	*A						PICK UP ADDRESSES
015620	71	015640	00	000000			007760		TC	I-P15						EXIT
015630	71	777777	00	000000			007770	INP14X	TC	777777						EXIT
015640	25	000243	00	015720			007780		RMK	SUBSTITUTE TO PICK UP						ADDRESSES
015650	72	005347	00	600000			007790	INP15	TCT	BSTP						SET EXIT
015660	24	000025	11	000033			007800		SET	WADCA,R						PICK UP A ADDRESS
015670	72	005356	00	600000			007810		SC	010025						
015700	24	000055	11	000063			007820		SET	WADCB,R						PICK UP B ADDRESS
015710	22	020114	00	014320			007830		SC	010055						
015720	71	777777	00	000000			007840	INP15X	OC	B2P						RESET SWITCH INP15
015730	25	000243	00	016170			007850		RMK							EXIT
015740	72	005257	00	600000			007860		RMK							MOVE COMMENT FROM
015750	35	005144	00	005257			007870	MVCOM	SET	IPCOM,R						INPUT TO W/A AND
015760	61	015770	00	016170			007880	MV1	SCR	IPCOM						ADJUST AMS
015770	73	000133	00	600000			007890		CTC	*						SET EXIT
016000	44	000133	00	017763			007920		STR	B3						RIGHT JUSTIFY
016010	72	000000	30	600000			007930		TC	B3						COMMENTS FIELD
016020	43	017760	00	017760			007940		SET	000000						FIELD HAVE ANYTHING
016030	61	016050	00	016050			007950		SC	B1A						STORE Y AS LHE
016040	71	016000	00	000000			007960		CTC	* 2						LOCATE FIRST
016050	25	015740	00	016070			007970		TC	B4						NONSPACE CHARACTER
016060	45	016073	00	000133			007980		TC	MV1						FROM LEFT END
016070	72	777777	50	600000			007990		TCT	MV1						
016100	73	016203	00	600000			008000	MV2	TCB	MV2J3						GET LENGTH OF FIELD
016110	45	016203	00	016207			008010		STR	MV3,R						AND SET Y TO RME
016120	61	016170	00	016130			008020		TCB	MV3,R						OF RECEIVING AREA
016130	24	000000	30	005257			008030		CTC	MVCOMX						WILL COMM W/A
016140	44	000153	00	016073			008040		SET	000000						OVERFLOW
016150	44	000153	00	017767			008050		TC	B5						MOVE TO COMM W/A
016160	34	777777	55	777777			008060		SET	000000						CREATE SPACE TO
016170	71	777777	60	000000			008070		TC	B5						RIGHT OF FIELD AND
							008080		SCC	777777						ADJUST AMS TO MSC
							008090		RMK							OF NEXT FIELD
							008100	MVCOMX	TC	777777						EXIT

NAME	AUTOFLOW	HSA	OP	A	N	B	SEC #	TAG	INDEX	40000501	OP	A	ADDRESS	SEG #	DATE	111665	OP	A	ADDRESS	COMMENTS	PAGE	019
016201	01	005101	01	010101					OP	DAR	3,1											
016204	01	005101	01	010101					OP	DAC	MACDHR											
016210	25	000243	00	016650					OP	TCT	SSP		PAREN									
016220	34	005117	00	005123					OP	SEC	HCLAB		HDLAB,R									
016230	34	005110	00	005116					OP	SEC	HDTAG		HDTAG,R									
016240	25	000173	00	003775					OP	TCT	ST		TEMP1									
016250	25	005130	00	000173					OP	TCT	RHECOM											
016260	72	000173	00	000000					OP	SET	OC0173		ST									
016270	43	000151	00	000153					OP	SC	OC0151		000153									
016300	61	016310	00	016320					OP	CTC	# 1		# 2									
016310	71	016660	00	000000					OP	TC	PAR3		ST									
016320	72	150001	00	000000					OP	SET	150001		ST									
016330	31	000001	57	000000					OP	LNS	OC0001		PAR3									
016340	61	016350	00	016660					OP	CTC	# 1		ST									
016350	25	000223	00	000113					OP	TCT	SSTA		ST									
016360	72	350001	00	000000					OP	SET	350001		ST									
016370	31	000001	51	000001					OP	LNS	OC0001		PAR4									
016400	61	016700	00	016700					OP	CTC	PAR4		ST									
016410	25	000223	00	000173					OP	TCT	SSTA		ST									
016420	22	020045	00	016643					OP	OCT	MR		PAR3,3									
016430	72	005123	00	000000					OP	SET	HCLAB,R		ST									
016440	24	000002	77	000006					OP	STC	OC0002		77 000006									
016450	72	150001	00	000000					OP	OCT	MR		ST									
016460	61	005117	00	005123					OP	SET	150001		ST									
016500	34	000001	20	005123					OP	LNS	HCLAB		HDLAB,R									
016510	72	005116	00	000000					OP	CTC	# 1		# 2									
016520	24	000001	55	000007					OP	SEC	OC0001		20 HDLAB,R									
016530	72	350001	00	000000					OP	SET	HDTAG,R		ST									
016540	31	005110	00	005116					OP	SET	350001		ST									
016550	61	016540	00	016570					OP	LNS	HCTAG		HDTAG,R									
016560	34	000001	20	005116					OP	CTC	# 1		# 2									
016570	34	000001	51	000001					OP	SEC	OC0001		20 HDTAG,R									
016600	25	000113	00	000153					OP	SEC	OC0000		ST									
016610	44	000153	00	017763					OP	TCT	ST		\$									
016620	22	017757	05	777777					OP	TCA	ST		\$									
016630	25	003775	00	000173					OP	OCT	TR47		\$									
016640	72	000077	00	100000					OP	TCT	TR47		ST									
016650	71	777777	00	000000					OP	SET	OC0077		\$PRI									
016660	22	020002	00	016643					OP	TC	777777		PAR2,3									
016670	71	016630	00	000000					OP	OCT	MR		PAR2,3									
016700	22	017760	00	016643					OP	TC	PAR2-1		PAR2,3									
016710	22	020140	01	000001					OP	OCT	MR		PAR2,3									
016720	71	016510	00	000000					OP	OCT	MR		ST									
016730	25	000243	00	017050					OP	CHK	EXIT		CH'EXIT									
016740	25	004040	00	000133					OP	TCT	CH'LC		\$									
016750	25	004050	03	000003					OP	TCT	SPD		\$									
016760	25	005174	03	000007					OP	TCT	LCV		\$									

EXTRACT TAG AND
 LABEL FROM
 COMMENTS FIELD,
 PRZ-LABEL PRESENT
 PRZ=NO TAG
 PRN=NO LABEL
 SET EXIT
 CLEAR HOLD AREAS
 PRESERVE ANY
 RHE COMM TO ANY
 COMMENT FIELD
 EXHAUSTED

LOCATE RIGHT END
 OF ASTERISK FIELD
 RIGHT END FOUND
 RHE TO ANY
 LOCATE COMMA
 WITHIN FIELD
 HAS COMMA FOUND

SET SIGNAL FOR PRP
 MOVE LABEL TO
 HOLD AREA, LEFT
 JUSTIFY, AND
 CLEAR TO RIGHT

MOVE TAG TO
 HOLD AREA, LEFT
 JUSTIFY, AND
 CLEAR TO RIGHT

CLEAR ASTERISK FIELD
 ADJUST ANS TO NEW
 LMS OF COMM FIELD
 ISS TO HSC COMM
 RESTORE ANY
 SET PRIS
 EXIT
 NO ASTERISK FIELD
 SET SIGNAL FOR PRZ
 GO LABEL IN FIELD
 SET SIGNAL FOR PRN
 SET COMMA OVER STAR

ENTER NEW CHAIN
 IN CHAIN TABLE
 SET EXIT
 PICK UP LSC OF TABLE
 SET # TO TABLE
 # LINES TO TABLE

NAME	AUTOFLD	OP	A	V	B	SEC #	TAG	INDEX	4000	501	N	SEG #2	DATE	111665	PAGE 020	COMMENTS
016770	22	017761	03	000000		006190	OP A ADDRESS					03	000000		SET IND TO GO	
017000	22	017761	03	000014		006700	OCT #01					03	000014		1ST RCD INDIC OF P	
017010	44	004043	00	020027		006710	R/K					000010#1			NEXT CHAIN TO NOP C	
017020	44	004037	00	017763		006720	TCA C-LOG,R					000001#1			ADVANCE LSC OF TABLE P	
017030	25	020043	00	005124		006730	TCA N-CHG,R					LMCT			ADD TO CHAI' COUNTER P	
017040	22	007246	00	007245		006740	TCT #00000#1					FP#15			RESET LINE CTR P	
017050	71	777777	00	000000		006750	OCT FP#6					777777			SET SWITCH FPS P	
						006760	CHEATX								EXIT F	
						006770	R/K								EDIT LINE T	
						006780	R/K								SURROUTINE=MOVE A C	
						006790	R/K								LINE FROM I/P TO C	
						006800	R/K								O/P. R ADDRESS C	
						006810	R/K								OF TC CONTAINS C	
						006820	R/K								NO OF CHAR IN LINE C	
						006830	R/K								* PRZ=END OF FIELD C	
						006840	R/K								PRZ=NOT END FIELD C	
						006850	R/K								PRZ=WORD TOC RIG C	
						006860	R/K								STP TO AH1 P	
						006870	R/K								PICK UP LINE LENGTH P	
						006880	R/K								AM3 TO MAX RIGHT END P	
						006890	R/K								END MESSAGE WITHIN D	
						006900	R/K								LIMITS OF LINE C	
						006910	R/K								LAST LINE OF COMMENT N	
						006920	R/K								COMPUTE TRUE LENGTH P	
						006930	R/K								OF FINAL LINE C	
						006940	R/K								SET RIGHT AND LEFT P	
						006950	R/K								LIMITS OF LINE C	
						006960	R/K								MOVE LINE TO O/P P	
						006970	R/K								SET SIGNAL FOR PRZ C	
						006980	R/K								ADD TO LINE CTR P	
						006990	R/K								SET PRI INDICATOR P	
						007000	R/K								EXIT F	
						007010	R/K								NOT LAST LINE P	
						007020	R/K								SET LOCATION OF P	
						007030	R/K								RIGHTMOST CHARACTER C	
						007040	R/K								IS CHARACTER A SPACE D	
						007050	R/K								ISS OVERLAYS SPACE P	
						007060	R/K								SET AM3 AND AM7 TO P	
						007070	R/K								TO LINE BOUNDARIES C	
						007080	R/K								MOVE LINE TO O/P P	
						007090	R/K								SET SIGNAL FOR PRN P	
						007100	R/K								SHIFT LEFT BY 1 CHAR J	
						007110	R/K								SET PRP *ED04 J	
						007120	R/K								SKTP EXTRA SPACES TO RIGHT OF PAREN FIELD IN COMMENTS SET EXIT	
						007130	R/K								CURRENT LINE A SPACE	
						007140	R/K								ADVANCE LME BY ONE	

NAME	AUTOFLOW	MS#	OP	A	N	B	SEG #	TAG	INDEX	4000501	OP A ADDRESS	SEG 02 DATE 111665	N B ADDRESS	PAGE 021
		017510	71	017450	00	000000	009270							COMMENTS
		017520	22	017757	05	777777	009280	SKSP2	TC	SKSP1				RECYCLE
		017530	71	777777	00	000000	009290	SKSPX	TC	777777		65 777777		PUT IN FIELD MARK
		017540					009300		BAR	9A				EXIT
		017760							CON	22,0				FOR PATCHES
		017770							CON	22,0				
		017756							CON	01				
		017757							CON	01				
		017760							CON	01				
		017761							CON	03,1				
		017764							CON	01				
		017765							CON	03,1				
		017770							CON	04				
		017775							CON	03,1				
		020000							CON	02				
		020002							CON	01				
		020003							CON	01				
		020004							CON	01				
		020005							CON	01				
		020006							CON	01				
		020007							CON	01				
		020010							CON	01				
		020011							CON	01				
		020012							CON	01				
		020013							CON	01				
		020014							CON	01				
		020015							CON	01				
		020021							CON	03,1				
		020025							CON	03,1				
		020030							CON	01				
		020031							CON	03,1				
		020032							CON	03,1				
		020033							CON	03,1				
		020041							CON	03,1				
		020044							CON	01				
		020045							CON	01				
		020051							CON	03,1				
		020055							CON	03,1				
		020060							CON	02				
		020065							CON	03,1				
		020070							CON	02				
		020075							CON	03,1				
		020101							CON	03,1				
		020104							CON	03,1				
		020111							CON	02				
		020113							CON	01				
		020114							CON	01				
		020115							CON	03				
		020120							CON	02				
		020122							CON	02				
		020124							CON	01				
		020125							CON	01				

NAME	AUTOFLOW	MS#	OP	A	N	B	SEG #	TAG	INDEX	4000 501	OP A ADDRESS	SEG 02 DATE 111665	N B ADDRESS	PAGE 022
		020126						#0053	CON	09				COMMENTS
		020137						#0054	CON	01				
		020140						#0055	CON	01				
		020150	00	000000	00	000002			CON	0,1				

NAME	AUTOFLOW	MSH	OP	A	N	B	SEC 4	TAG	OP	A	ADDRESS	INDEX	40000501	SEC 03	DATE	111665	COMMENTS	PAGE 023	
004170	25	000243	00	004240			009310	LOCTAG	TCT	65TP									
004200	42	000000	30	600001			009320		SET	000000									
004210	43	004207	00	004207			009330		SC	0117									
004220	61	004250	00	004250			009340		CTC	LOC1									
004230	72	000001	00	100000			009350		SET	000001									
004240	71	777777	00	000000			009400	LOCTAG	TC	777777									
004250	24	004012	00	600000			009410	LOC1	SET	TEMP3,R									
004260	24	000000	33	000006			009420		STC	000000									
004270	25	004020	00	004001			009430		TCT	NUMTAG									
004300	25	004030	00	000153			009440		TCT	TOR									
004310	56	004340	00	004001			009450	LOC2	TA	LCC3									
004320	72	000004	00	100000			009510		SET	000004									
004330	71	004240	00	000000			009520		TC	LOCTAGX									
004340	72	000007	50	600000			009530	LOC3	SET	000007									
004350	43	004004	00	004012			009540		SC	TEMP3,R									
004360	61	004410	00	004410			009550		CTC	LCC4									
004370	72	000002	00	100000			009560		RMK										
004400	71	004240	00	000000			009570		SET	000002									
004410	44	000153	00	004017			009580		TC	LOCTAGX									
004420	71	004310	00	000000			009590	LOC4	TA	RCC3									
004430							009600		TC	LCC2									
004440							009610	LOAD4	CON	0,10									
004450	01	023209	01	010101			009620	PAC4	CON	4									
004460	01	023510	01	010101			009630	BCC4	CON	4									
004470	01	000001	01	010101			009640	STOR3	DAC	IPAR1									
004480	01	000001	01	010101			009650	STOR4	DAC	IPAR2									
004490	01	000000	01	010101			009660	COLVO	DAC	000001									
004500	01	000001	01	010101			009670	TP05	DAC	000000									
004510	01	000001	01	010101			009680	TS0	DAR	3,1									
004520	01	000001	01	010101			009690	LCC1	DAC	0-0001									
004530	01	000001	01	010101			009700	LCC2	DAC	000001									
004540	01	777777	01	010101			009710	LCC3	DAC	777777									
004550	01	000006	21	010101			009720	DECA0	RMK										
004560	01	000006	21	010101			009730	STOR5	R-K										
004570	01	000006	21	010101			009740	STOR6	DAR	3,1									
004580	01	000006	21	010101			009750	STOR7	DAR	3,1									
004590	01	000006	21	010101			009760	STOR8	DAR	3,1									
004600	01	000006	21	010101			009770	TP01	DAR	3,1									
004610	01	000006	21	010101			009780	SP5T1	DAR	3,1									
004620	01	000006	21	010101			009790	EV00	CON	1									
004630	01	000006	21	010101			009800	PSRCD	DAR	1,00									
004640	22	004055	00	004615			009810	STPAS2	R-K										
004650	17	000000	00	000000			009820		CON										

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USED AS STORAGE
 TO READ RYN
 COLUMN #
 SEC # OF LAST REC
 DEST TAG SEC #
 ADDRESS COL1
 ADDRESS COL2
 TEMP STORAGE TO
 HOLD ADDRESS
 USED IN LOGGING
 COLUMN 2 ADDRESS
 STORAGE FOR INFO
 RELATING TO E GAPS

STORAGE FOR TP05

40 IF EVERY COL
 PSELOC COUNT REC
 BEGIN PASS 2
 RELAND INPUT AND
 OUTPUT TAPES

00

0 115



NAME	AUTOFLOW	A	N	B	SEG #	TAG	INDEX	4000-501	N	SEG #3	DATE	111665	PAGE 024	COMMENTS
004620	22	004054	CC	004635	009580		OP							
004630	17	000000	CC	000000	009590		OCT T13							
004640	25	004044	CC	000113	009500		RWD							
004650	22	0111672	CC	004554	009510		TCT CH1							
004660	22	0111673	CC	004570	009520		RK #76#							
004670	22	0111674	CC	004555	009930		OCT #75#							
004700	75	000000	CC	010000	009940		OCT #67#							
004710	22	004055	CC	004725	009960		RVK							
004720	15	003010	CC	000000	009970		CSC							
004730	62	004730	CC	004730	009980		OCT T12							
004740	72	003010	CC	000000	010000		SPF OFAR							
004750	43	0111675	CC	001675	010010		SS4 #							
004760	61	000240	CC	005000	010020		SET DPAB							
004770	71	007240	CC	000000	010030		SC #72#							
005000	25	003054	CC	004484	010040		CIC 000240							
005010	72	0111676	CC	000000	010050		TC ENDR2							
005020	43	003024	CC	003024	010060		TCT OPS0							
005030	61	005050	CC	005050	010070		SET #4#							
005040	71	006730	CC	000000	010080		SC CODE							
005050	25	004474	CC	004504	010090		CIC # 2							
005060	44	004477	CC	003013	010100		TC SP40							
005070	72	004477	CC	000000	010110		TCT LNC1							
005100	43	004115	CC	004117	010120		TCA LNC1,R							
005110	61	005120	CC	005500	010130		SET LNC1,R							
005120	72	011877	CC	000000	010140		SC EPC0,R=2							
005130	43	003024	CC	003024	010150		CIC # 1							
005140	61	005140	CC	005140	010160		SET #							
005150	71	005270	CC	000000	010170		SC CODE							
005160	51	004447	CC	011701	010180		CIC # 2							
005170	71	010300	CC	000000	010190		TC SP40							
005200	25	004104	CC	000133	010200		TCT LNC1							
005210	71	004170	CC	000000	010210		TCA LNC1,R							
005220	61	005270	CC	005270	010220		SET LNC1,R							
005230	72	000017	CC	000000	010230		SC EPC0,R=2							
005240	26	003020	CC	003023	010240		CIC # 1							
005250	25	004504	CC	000053	010250		SET #							
005260	22	004463	CC	000020	010260		TC SP4							
005270	25	004504	CC	003017	010270		TC LOCTAG							
005300	22	004463	CC	003014	010280		CIC #2							
005310	22	003024	CC	005321	010290		TC SP4							
005320	72	000001	CC	000000	010300		SET 000017							
005330	31	011702	CC	000000	010310		SET BOX0							
005340	61	005390	CC	005360	010320		TC LNC3							
005350	71	007050	CC	000000	010330		OCT COLNO,R							
005360	72	011705	CC	000000	010340		TCT LNC3							
005370	43	003024	CC	003024	010350		OCT COLNO,R							
005400	61	005420	CC	005420	010360		OCT CODE							
005410	71	006000	CC	000000	010370		SET 000001							
005420	71	010360	CC	000000	010380		OCT #							
005430	22	003024	CC	005441	010390		LNS "DJS"							
005440	72	770001	CC	000000	010400		CIC # 1							
005450	31	011706	CC	011707	010410		TC D1							
005460	61	005470	CC	004710	010420		SET #4#							
005470	71	005410	CC	000000	010430		LNS #514#							
					010440		TC SP12							
					010450		RVK							

SET AW1 TO START

OF CHAIN TABLE

SET UP PSEUDO

CON RECORD

X 1 2ND POSITION

HEA'S PSEUDC CONN

BEGIN MAIN LOOP

READ RECORD

END FILE

STORE SEG #

B CODE

SET SYMBOL ADDRESS

ADD TO LINE CTR

ROOM IN COLUMN

FOR THIS RECORD

T CODE

INCREASE BOX #

MOVE BOX # TO O/P

ADDR OF TAG TO AM3

FIND TAG IN TABLE

ENTRY FOUND IN TABLE

PUT BOX # IN

TAG TABLE

PUT P/C ADDRESS

IN TAG TABLE

PUT P/C ADDRESS

IN RECORD

D OR J CODE

DECISION

WRITE RECORD TO O/P

J OR E CODE

IF NOT A JUMP, GO

NAME	AUTOFLOW	AS	OP	A	B	SEP 4	TAG	OP	A ADDRESS	4000501	SEP 03	DATE	111665	R	B ADDRESS	COMMENTS	PAGE 026
006150	72	004523	00	600000		011140	R4K	OP	A ADDRESS							LOCATE, PROCESS	
006160	43	004515	00	004517		011150	R4K	TCT	ALHT							CONING FOR COL 2.	
006170	61	006200	00	006270		011160	R4K	TCT	ALHT							AM3 HAS MSC OF	
006200	25	004100	00	000133		011170	R4K	TCT	ALHT							DEST TAG#47 HAS	
006210	22	011725	03	777777		011180	R4K	TCT	ALHT							LOC IN CHAIN TABLE	
006220	25	004521	00	004525		011190	R4K	TCT	ALHT							STORE SEQ # OF	
006230	25	004074	00	000133		011200	R4K	TCT	ALHT							LAST RECORD READ	
006240	22	011726	03	777777		011210	R4K	TCT	ALHT							WRITE DECISION RCD	
006250	25	004531	00	000173		011220	R4K	TCT	ALHT							TO OUTPUT FILE	
006260	71	006440	00	000000		011230	R4K	TCT	ALHT							TAPE INSTRUCTION	
006270	25	004074	00	000133		011240	R4K	TCT	ALHT							UNWIND GAPS TO	
006300	22	011725	03	777777		011250	R4K	TCT	ALHT							REACH CHAIN	
006310	25	004515	00	004525		011260	R4K	TCT	ALHT							SET INDICATOR IN	
006320	25	004100	00	000133		011270	R4K	TCT	ALHT							TABLE THAT THIS	
006330	22	011726	03	777777		011280	R4K	TCT	ALHT							CHAIN IS PROCESSED	
006340	71	006440	00	000000		011290	R4K	TCT	ALHT							STORE MAIN COLUMN #	
006350	25	004100	00	000133		011300	R4K	TCT	ALHT							SET COL # TO	
006360	71	010470	00	000000		011310	R4K	TCT	ALHT							EITHER 2 OR 4	
006370	61	005420	00	005420		011320	R4K	TCT	ALHT							IN MAIN COLUMN	
006400	25	004515	00	004525		011330	R4K	TCT	ALHT							SET ADDRESS FOR	
006410	71	006320	00	000000		011340	R4K	TCT	ALHT							EVEN COLUMN TO 3	
006420	25	004521	00	004525		011350	R4K	TCT	ALHT								
006430	71	006230	00	000000		011360	R4K	TCT	ALHT								
006440	25	004444	00	004535		011370	R4K	TCT	ALHT								
006450	71	010300	00	000000		011380	R4K	TCT	ALHT								
006460	25	004525	00	006500		011390	R4K	TCT	ALHT								
006470	22	004055	00	006505		011400	R4K	TCT	ALHT								
006500	06	777777	00	000000		011410	R4K	TCT	ALHT								
006510	22	011727	07	000000		011420	R4K	TCT	ALHT								
006520	25	004460	00	004541		011430	R4K	TCT	ALHT								
006530	44	004463	00	011733		011440	R4K	TCT	ALHT								
006540	25	004474	00	004545		011450	R4K	TCT	ALHT								
006550	45	004477	00	004513		011460	R4K	TCT	ALHT								

1ST TAG NOT SUITABLE FOR COL 2, TRY 2ND MSC 2ND TAG TO AM3
 DETERMINE STATUS OF 2ND TAG SUITABLE FOR ASSIGNMENT TO COL 2
 USE TAG 2 FOR COL 2 SET # GAPS
 USE TAG 2 FOR COL 2 SET # OF GAPS
 LOCATE, PROCESS CONING FOR COL 2. AM3 HAS MSC OF DEST TAG#47 HAS LOC IN CHAIN TABLE STORE SEQ # OF LAST RECORD READ WRITE DECISION RCD TO OUTPUT FILE PUT # GAPS INTO TAPE INSTRUCTION UNWIND GAPS TO REACH CHAIN SET INDICATOR IN TABLE THAT THIS CHAIN IS PROCESSED STORE MAIN COLUMN # SET COL # TO EITHER 2 OR 4 STORE CURRENT ADDR. IN MAIN COLUMN SET ADDRESS FOR EVEN COLUMN TO 3

NAME	AUTOFLOW	HSN	TP	A	N	B	SEQ #	TAG	INDEX	4000:503	N. B. ADDRESS	SEG #3	DATE	111665	PAGE #2B	COMMENTS
007120	41	000000	55	000000			012190	BA	010000		55 010000					OF CROSS REFERENCES
007130	72	000000	50	000000			012200	SET	010000		50 ST					FIRST REFERENCE
007140	43	011742	00	011742			012210	SC	R11F		50 R11F					TO THIS POINT
007150	61	007160	00	007200			012220	CTC	* 1		DJ2					YES/DJ2/NO*
007160	72	000033	50	600000			012230	SET	010033		50 ST					BOX # OF CURRENT P
007170	24	003020	00	003023			012240	STC	BOX#0		53					RECORD TO TABLE
007200	25	004100	00	000133			012250	TCT	ALMT		DJ3					SET FOR END TAG
007210	66	007070	00	007230			012260	TA	DJ1							BOTH TAGS PROCESSED
007220	71	005360	00	000000			012270	TC	SF5							RETURN TO MAIN PATH
007230	01	000000	01	010101			012280	DAC	010000							TALLY FIELD
007240	72	000000	00	000000			012290	RMK								END OF PASS 2
007250	72	004550	00	007265			012300	IGN			* 1.5					IF EVEN COLUMN
007260	72	000240	00	000000			012310	OCT	EV00							ILLOGICAL HALT
007270	72	011742	00	600000			012320	SET	010240		ST					BYPASS WRITING FINAL
007300	43	004445	00	004446			012330	SET	R11F		ST					EF IF BOX # IS 00
007310	61	007320	00	007340			012340	CTC	* 1		* 3					WRITE EF FOR LAST
007320	22	004056	00	007335			012350	OCT	TT3		* 1.5					-PAGE OF CHART
007330	12	011675	00	770000			012360	LW	R72H		770000					REWIND OUTPUT
007340	22	004056	00	007355			012370	OCT	TT3		* 1.5					
007350	12	007357	00	000073			012380	LW	* 7		000073					
007360	22	004056	00	007375			012390	OCT	TT3		* 1.5					
007370	17	000000	00	000000			012400	R-D			* 1.5					REWIND INPUT
007400	22	004055	00	007415			012410	R-D	TT2							READ SEQ 04
007410	17	000000	00	000000			012420	RWD	LOAD4		04 STPASS					
007410	17	000000	00	000000			012430	RIS	LOAD4		ST					
007410	17	000000	00	000000			012440	SET	001167		* 3.7					
007420	72	001167	00	600000			012450	STC	* 2							
007430	24	007450	00	007467			012460	TC	010340							
007440	71	000340	00	000000			012470	OO	LOAD4							
007450	00	004430	00	000000			012480	OO	04							
007460	00	000004	00	007450			012490	RMK			STPASS					
007470	25	000243	00	007710			012500	RMK								
007500	25	004504	00	004563			012510	RMK								
007510	22	004463	00	004560			012520	RMK								
007520	72	011743	00	600000			012530	ENDCOL			ENDCOLX					
007530	43	003024	00	003024			012540	TCT	STTP		PSRCD17					
007540	61	007560	00	007560			012550	TCT	LNC3		PSRCD14					
007550	71	007720	00	600000			012560	OCT	COLNO,R		ST					
007560	72	011724	00	600000			012570	SET	* P		CODE					
007570	43	004443	00	004463			012580	CTC	* 2		* 2					
007600	61	007720	00	007610			012590	TC	ENDCOL2							
007610	72	011733	00	600000			012600	SC	COLNO,R		COLNO,R					
007620	43	004501	00	004503			012610	CTC	ENDCOL2		* 1					
007630	61	010010	00	010010			012620	SET	RC00001,R,1		ST					
007640	44	004463	00	011733			012630	SC	L1C2,R-2		LNC2,R					
007650	71	010040	00	000000			012640	CTC	ENDCOL3		ENDCOL3					
007660	25	011733	00	000474			012650	TCA	COLNO,R		RC00001,R,1					
007670	25	004474	00	004500			012660	TCT	RCN		LNC1					
007700	25	004474	00	004504			012670	TCT	LNC1		LNC2					
007710	71	777777	00	000000			012680	ENDCOLX	TC 777777		LNC3					
007720	25	011733	00	000460			012690	RMK								
007730	51	004443	00	011701			012700	ENDCOL2	TC RC00001,R,1		COLNO					
007740	21	011746	00	004446			012710	DA	PAC1,R		W4248					
								IT	W4232W		BXC1,2					

NAME	AUTOFLO.	LSN	OP	A	B	SEQ #	TAG	I'DIX	4000-501	A ADDRESS	SEG #	DATE	111665	PAGE #30	COMMENTS
010510	61	011170	OC	011150						SCOLL1					RESULT OF TABLE
010520	72	000017	5C	600000						50 ST					SEARCH FOUND SCOLL3
010530	43	011742	OC	011742						NOH					FOUND SCOLL1, NOT AG
010540	61	011150	OC	011150						SCOLL1					IS TAG ALREADY
010550	25	000013	50	004471						50 T50					ASSIGNED
010560	72	000003	1C	600000						10 ST					GET SEQ # OF TAG
010570	43	004471	OC	004473						T50,R					FROM CHAIN TABLE
010600	61	010620	OC	011150						SCOLL1					IS TAG LOCATED
010610	71	011150	OC	000000						SCOLL1					IF CURRENT CHAIN
010620	25	000113	OC	000173						ST					GET CHAIN OF
010630	44	000173	OC	011717						NO00010R,1					DESTINATION TAG
010640	72	000003	7C	600000						70 ST					ADV. TO NEXT ENTRY
010650	43	004471	OC	004473						T50,R					IF CHAIN TABLE
010660	61	010630	OC	01067C						ST					TAG PART OF THIS
010670	25	777773	7C	004001						70 TEMP					CHAIN TABLE ENTRY
010700	44	004003	OC	011733						ST					GET START OF THIS
010710	72	004003	OC	600000						NO00001R,1					CHAIN FROM TABLE
010720	43	004471	OC	004473						T50,R					IS TAG THE START
010730	61	011210	OC	011170						SCOLL3					OF A CHAIN
010740	25	004504	OC	004001						TEMP					IS COLUMN 2 FREE
010750	44	004003	OC	011757						NO00007R,1					AT THIS POINT
010760	45	004003	OC	004503						LHC2,R					ENOUGH ROOM IN COL 2
010770	61	011000	OC	011150						SCOLL1					TO HOLD ENTIRE CHAIN
011000	25	004474	OC	004001						TEMP2					
011010	45	004003	OC	004513						DECAD,R					
011020	44	004003	OC	000007						NO 000007					
011030	45	004003	OC	011737						EPCON,R					ENTIRE CHAIN
011040	45	004003	OC	004117						NO 00002R,1					
011050	61	011150	OC	011060						70 STOR5					COMPUTE # GAPS
011060	25	777773	7C	004515						TEMP2					TO REACH CHAIN
011070	45	004517	OC	004467						TPOS,R					# GAPS TO REACH
011100	25	004517	OC	004001						NO00144R,1					CHAIN TOO GREAT
011110	45	004003	OC	011763						STOR5					
011120	61	011170	OC	011130						TEMP2					
011130	72	000002	OC	100000						TEMP2					CHAIN OK FOR ADJ COL
011140	71	777777	OC	000000						TEMP2					SET PRZ
011150	72	000001	OC	100000						TEMP2					EXIT
011160	71	011140	OC	000000						TEMP2					CHAIN NOT SUITED
011170	22	011125	OC	777777						TEMP2					FOR ADJACENT COLUMN
011200	71	011150	OC	000000						TEMP2					SET PRN
011210	25	004001	OC	011300						TEMP2					TO EXIT
011220	44	011303	OC	011733						TEMP2					SET INDICATOR
011230	45	011303	OC	004473						TEMP2					FOR CONNECTOR
011240	61	011170	OC	011170						TEMP2					
011250	22	000004	7C	011265						TEMP2					
011260	72	010740	OC	000000						TEMP2					
011270	71	011170	OC	000000						TEMP2					
011300	01	000000	OC	010101						TEMP2					

GET CHAIN OF DESTINATION TAG
 ADV. TO NEXT ENTRY IF CHAIN TABLE
 TAG PART OF THIS CHAIN TABLE ENTRY
 GET START OF THIS CHAIN FROM TABLE
 IS TAG THE START OF A CHAIN
 IS COLUMN 2 FREE AT THIS POINT
 ENOUGH ROOM IN COL 2 TO HOLD ENTIRE CHAIN
 ENTIRE CHAIN
 COMPUTE # GAPS TO REACH CHAIN # GAPS TO REACH CHAIN TOO GREAT
 CHAIN OK FOR ADJ COL SET PRZ
 EXIT CHAIN NOT SUITED FOR ADJACENT COLUMN SET PRN
 TO EXIT SET INDICATOR FOR CONNECTOR
 IS TAG SECOND RECORD OF CHAIN
 NO IS 1ST RCD N OR T NO SCOL6, YES NO GOOD FOR COL 2 WORK AREA

NAME	AUTOFLOW	OP	A	B	SEQ #	TAG	OP	A ADDRESS	INDEX	4000-501	N	SEQ #	DATE	111665	PAGE	032
MSV							SEG OF				LOAD#					COMMENTS
004430					013690		DAR 1									PASS THREE
004600					013690	HDR1	DAR 120,200									USED FOR ST MESS
004600					013910		SLC H'R1									MAIN HEADER
004607					013920	HDP11	CON 7									
004655					013930		CON 1									
					013940		CON 1									
					013950	HDR12	CON 17									
					013960		CON 17									
004730					013970	HDR13	SLC H'R1,88									
					013980		CON 4									
004757					013990	HDR14	SLC H'R1,R=8									
					014000		CON 4									
004770					014010		SLC H'R1,R 1									
005000					014020	HDR2	CON 2									
005010					014030		DAR 120,200									
005170					014040	HUR21	SLC HLR2,6									
005172					014050		CON 6									
005200					014060		SLC H'R2,R 1									
005220					014070		CON 2									
005240					014080		DAR 1									
					014090	HDR3	DAR 120,200									
					014100	HDR31	SLC H'R3,16									
					014110		CON 16									
					014120		CON 14									
					014130		CON 14									
					014140		SLC H'R3,76									
					014150		CON 16									
					014160		CON 14									
					014170		CON 1									
005334					014180	PRAR	DAR 120,200									
005370					014190		RMK									
005372					014200		RMK									
005400					014210		CON 1									
005570					014220	ADDLAY	DAC LAYOUT									
005574					014230		RMK									
005600					014240	FCON1	CON 12									
005614					014250	FCON2	CON 9									
005625					014260		RMK									
005626					014270	FCON3	CON 1									
005630					014280	ONE	CON 2									
005634					014290	CTZR	DAC 00000									
005640					014300	CT5X	DAC 000006									
005644					014310	CTON	DAC 00001									
					014320	LNTH	DAC 000170									
					014330		RMK									
005651					014340	TALI	DAR 3,1									

COMMENTS
 PASS THREE
 USED FOR ST MESS
 MAIN HEADER

SUB HEADER

USED FOR ST MESS
 T OF C HEADER

AREA FROM WHICH
 ACTUAL PRINT OR
 WRITE IS DONE

MISC OF PAGE
 LAYOUT AREA

OF PROCESS BOX

CHAR PER LINE
 ON PAGE LAYOUT
 MISC TALLIES

NAME	AUTOFLOW	A	N	B	SEQ #	TAG	OP	A ADDRESS	INDEX	40000501	M	B ADDRESS	SEC 04	DATE	111668	PAGE	033	COMMENTS	
005629					014390	TAL2	DAR	3/1											
005661					014360	TAL3	DAR	3/1											
005665					014370	TAL4	DAR	3/1											
005670					014380	VTCON	CON	1											
	50																		
005671					014400	HZCON	CON	12											VERTICAL LINE
005705					014410		CON	17											
005724					014420	CNLW	DAC	000016											
005730					014440	NMBL	DAC	000170											
005734					014470	PNO	EOL	010000											
005742					014480	PCMES1	CON	6											
005746					014490	PCMES2	CON	4											
005752					014500	CRC1	CON	7											
005761					014510		CON	7											
005770					014520		CON	7											
005777					014530		CON	7											
006006					014540		CON	7											
006020					014550	CRC2	DAC	000005											
006024					014560	CRC3	DAC	000007											
006030					014570	NCON1	CON	17											
006044					151515151515151515		CON	5											
006051					1515151515		CON	1											
006052					15		CON	4											
006056					55566344		CON	7											
006065					14141414141414		CON	7											
006071					47405363		CON	4											
006100					014640	SCON	CON	7											
006107					014650		CON	7											
006116					014660		CON	7											
006125					014670		CON	7											
006134					014680		CON	7											
006143					014690		CON	7											
006152					014700		CON	7											
006161					014710		CON	7											
006170					014720		CON	7											
006177					014730		CON	7											
006206					014740		CON	7											
006215					014750		CON	7											
006224					014760		CON	7											
006233					014770		CON	7											
006242					014780		CON	7											
006251					014790		CON	7											
006260					014800		CON	7											
006267					014810		CON	7											
006276					014820		CON	7											
					014830		CON	7											

CHARACTER FROM
 LEFT SIDE OF COL
 # TO CENTERLINE
 # LINES PER PAGE
 OFF-LINE BREAKPNT

LINES TO CIRCLE
 LENGTH OF LINE

LINES OF HALT

NAME	AUTOFLOW	CP	A	B	SEC #	TAG	INDEX	ADDRESS	DATE	SEC #	ADDRESS	COMMENTS	PAGE
008305					014440	OP	4000*501		111665			PAGE 034	
008314					014450	CON						COMMENTS	
					014460	DAC						# LINES=1 1/2	
					014470	RMK						SUBROUTINE SYMBOL	P
					014480	RMK						OUTPUT ROUTINE	T
					014490	RMK						SUBROUTINE TO	C
					014510	RMK						PRINT OR WRITE	C
					014920	RMK						OUTPUT, IF RPO	C
					014930	RMK						IS ON, WRITE	C
					014940	RMK						TO OFF-LINE TAPE	C
					014950	RMK						SET EXIT	P
					014960	RMK						MOVE PAGE # TO HDR	P
					015000	RMK						ADD 1 TO PAGE #	P
					015010	RMK						OFF-LINE	D
					015020	RMK						WRITE MAIN AND	P
					015030	RMK						SUB HEADERS	C
					015040	RMK						A ADD VARIABLE	P
					015050	RMK						SET TALLY TO	P
					015060	RMK						# LINES ON PAGE	C
					015070	RMK						INITIALIZE AM3	C
					015080	RMK						ADJUST TALLY TO	P
					015090	RMK						AVOID PRINTING ANY	P
					015100	RMK						BLANK LINES AT	C
					015110	RMK						BOTTOM OF PAGE	C
					015120	RMK						SET AM3 TO BOTTOM	C
					015130	RMK						LINE OF PAGE	C
					015140	RMK						RESET TALLY TO	P
					015150	RMK						# LINES ON PAGE	P
					015160	RMK						TEST LINE FOR	P
					015170	RMK						ALL SPACES	P
					015180	RMK						REDUCE TALLY	P
					015190	RMK						THIS IS BLANK PAGE	P
					015200	RMK						SET FOR LINE ABOVE	P
					015210	RMK						RECYCLE	P
					015220	RMK						INITIALIZE AM3	P
					015230	RMK						ADJUST TALLY	P
					015240	RMK						IS AN EXTRA P/C	P
					015250	RMK						REQUIRED TO GIVE	P
					015260	RMK						A TFO-PAGE CHART	P
					015270	RMK						SET EXTRA P/C SW	P
					015280	RMK						MOVE LINE TO	P
					015290	RMK						OUTPUT AREA	P
					015300	RMK						OFF-LINE	C
					015310	RMK						WRITE LINE TAPE	D
					015320	RMK						TO PRINT TAPE	P
					015330	RMK						SET FOR NEXT LINE	P
					015340	RMK						END OF PAGE	D
					015350	RMK						CLEAR PAGE AREA	P
					015360	RMK						PRINT OR WRITE	P
					015370	RMK						PAGE CHANGE	P
					015380	RMK						FALL THRU IF EXTRA	C
					015390	RMK						P/C IS REQUIRED	C
					015400	RMK						TO GIVE THO-PAGE	C
					015410	RMK						CHART	C

NAME	AUTOPLOM	MSM	OP	A	M	B	SEG #	TAG	INDEX	40000501	N	B	DATE	311665	PAGE	035	COMMENTS
007050	71	007000	00	000000			015420	ENPG3A	OP A ADDRESS								
007060	71	007000	00	000000			015430	ENPG4	TC ENPG3A								
007070	02	004000	00	000000			015440	ENPG5	TC 777777								
007100	03	000000	00	000000			015450	PA 000002	PR MOR1								
007110	02	003200	00	000000			015460	ENPG5A	PR MOR2								
007120	03	000002	00	000000			015470	PA 000002	PA 000002								
007130	71	006440	00	000000			015480	ENPG6	TC ENPG1								
007140	75	000000	00	000000			015490	ENPG6	PR PRAR								
007150	02	003400	00	000000			015500	PA 000001	PA 000001								
007160	03	000001	00	000000			015510	PA 000001	PA 000001								
007170	75	000000	00	010000			015520	C5G ENP93	C5G ENP93								
007200	71	006750	00	000000			015530	ENPG7	TC ENP93								
007210	75	000000	00	000000			015540	ENPG7	C5G PA 020000								
007220	03	020000	00	000000			015550	ENPG7	C5G PA 020000								
007230	75	000000	00	010000			015560	SPACES	TC ENPG3B								
007240	71	007030	00	000000			015580	PCPAL	DAR 12000								
007250							015590	LOADS	DAC 000075								
007440	01	000075	01	010101			015600	LOADS	RMK DAR 8.10								
007450	36	007250	00	007437			015610	LOADS	SLC LOADS								
007460	72	004620	00	600000			015620	LOADS	RMK								
007470	24	004154	00	034165			015630	STPASC	RMK								
007500	72	004661	00	600000			015640	STPASC	RMK								
007510	24	000470	00	004155			015650	STPASC	RMK								
007520	22	000271	00	004735			015660	STPASC	SET MOR1, R 10								
007530	22	000272	00	004736			015670	STPASC	STC INDEX								
007540	22	011326	00	004737			015680	STPASC	SET MOR1, R 13								
007550	22	000273	00	004740			015690	STPASC	STC TITLE								
007560	22	000274	00	004741			018700	STPASC	OCT DC0271								
007570	22	011326	00	004742			018710	STPASC	OCT DC0272								
007600	22	000275	00	004743			018720	STPASC	OCT R22								
007610	22	000274	00	004744			018730	STPASC	OCT DC0273								
007620	71	007720	00	010000			018740	STPASC	OCT DC0274								
007630	22	011327	00	004577			018750	STPASC	OCT DC0275								
007640	22	011327	00	004577			018760	STPASC	TC S131								
007650	22	011327	00	005177			018770	STPASC	OCT R768								
007660	22	011327	00	005377			018780	STPASC	OCT R769								
007670	22	004057	00	007705			018790	STPASC	OCT R770								
007700	12	005742	00	000000			015600	ST31	TC S132								
007710	71	007730	00	000000			015610	ST32	TC S132								
007720	03	020000	00	000000			015620	ST32	PA 020000								
007730	22	004056	00	007745			015630	ST32	OCT T13								
007740	17	000000	00	000000			015640	ST32	RMD								
007750	36	023000	00	004027			015650	ST32	SET L1YFR								
007760	25	004020	00	005665			015660	ST32	RMK								
007770	25	004030	00	0050133			015670	ST32	RMK								
010000	25	005574	00	000173			015680	ST32	TCT N1TAG								
							015690	ST32	TCT I 1								
							015700	ST32	RMK								
							015710	ST32	TCT A, DLAY								
							015720	ST32	RPF								

EXIT
PRINT MAIN AND
SUB HEADERS
A ADD VARIABLE
OPEN GATE
PRINT LINE AND
PAPER ADVANCE
CLOSE GATE
PAGE CHANGE
AREA TO TEST SPACES
CONSTANT FOR TEST
OF EXTRA P/C
LOAD P/SEG 5
PRODUCE TABLE
OF CONTENTS
SET SPACES CONSTANT
PUT INDEX SERIAL #
IN MAIN HEADER
PUT TITLE IN
MAIN HEADER
MOVE DATE FROM
STANDARD LOCATION
TO HEADER
OFF-ONE
PUT IN START
MESSAGES
WRITE P/C MESSAGE
PAPER ADVANCE P
AND I/P TAPE
CLEAR PAGE AREA P
BEGIN CODING TO
PRINT TABLE OF
CONTENTS
SET TALLY TO N TAG
SET AMS TO 'SC
OF TAG TABLE
INITIALIZE FOR
NEW PAGE
SET AMS TO TGP
OF COLUMN 1

000000
010000
000000
010000
SPACES, R
ST INDEX R
ST
TITLE R
MOR1, R 2
MOR1, R 3
MOR1, R 4
MOR1, R 5
MOR1, R 6
MOR1, R 7
MOR1, R 8
MOR1, R 9
BPO
MOR1, R 1
MOR2, R 1
MOR3, R 1
PRAR, R 1
* 1,5
* 1,5
LAYOUT, R
TALA
35
17

NAME	AUTOFLO.	MS.	OP	A	B	SEQ. #	TAG	OP	A	ADDRESS	IN	SEG. #	DATE	111665	PAGE 036	COMMENTS.
010010	22	010257	00	010255		016000		OP	A <td>ADDRESS</td> <td></td> <td></td> <td></td> <td></td> <td></td> <td></td>	ADDRESS						
010020	25	004114	00	005655		016010		OP	A <td>ADDRESS</td> <td></td> <td></td> <td></td> <td></td> <td></td> <td></td>	ADDRESS						
010030	66	010050	00	005665		016020	TBCN3	TCT	EPCN							
010040	71	010550	00	000000		016030	TBCN4	TA	T	CN5						
010050	72	000001	50	600000		016040	TBCN5	TC	T	CN13						
010060	43	011330	00	011330		016050		SET	0	0001						
010070	61	010110	00	010110		016060		SC	R77							
010100	71	010340	00	000000		016070		CTC	2							
010110	72	000026	70	600000		016080		TC	T	CN10						
010120	24	000001	55	000007		016090	TBCN6	SEY	0	0026						
010130	72	000032	70	600000		016100		STP	0	0001						
010140	24	000014	55	000015		016110		SET	0	0032						
010150	72	000037	70	600000		016120		STC	0	0014						
010160	24	000016	55	000017		016130		SET	0	0037						
010170	72	000052	70	600000		016140		STC	0	0016						
010200	24	000024	55	000027		016150		SET	0	0052						
010210	22	000010	57	000044		016160		STC	0	0024						
010220	44	000173	00	005647		016170	TBCN7	OCT	0	0010						
010230	44	000153	00	004017		016180		TCA	87							
010240	66	010030	00	005655		016190		RNK								
010250	72	010320	00	004000		016200		TCA	65							
010260	25	005574	00	000173		016210		RNK								
010270	44	000173	00	011333		016220		TA	T	CN4						
010300	22	010256	00	010255		016230		RNK								
010310	71	010070	00	000000		016240		RNK								
010320	71	006370	00	000000		016250		RNK								
010330	71	010000	00	000000		016260	TBCN8	SET	T	CN9						
010340	72	005657	00	600000		016270		TC	ADL							
010350	43	011334	00	011336		016280		TCA	87							
010360	61	010330	00	010370		016290		OCT	T	CN8,6						
010370	72	005657	00	600000		016300	TBCN9	TC	T	CN3						
010400	43	004115	00	004117		016310		TC	E	P						
010410	61	010430	00	010430		016320		RNK								
010420	71	010450	00	000000		016330		RNK								
010430	44	000173	00	005647		016340		TC	T	CN2						
010440	66	010450	00	005657		016350		RNK								
010450	72	000014	70	600000		016360	TBCN10	SET	T	AL2,R						
010460	24	010520	00	010524		016370		SC	E	CN0,6						
010470	72	000045	70	600000		016380		CTC	T	CN12						
010500	24	000002	55	000027		016390		SET	T	AL2,R						
010510	71	010220	00	000000		016400		SC	E	CN0,R,Z						
010520	44	005667	00	011343		016410		CTC	2							
010530	44	005667	00	011343		016420		TC	2							
010540	71	010250	00	000000		016430		TCA	87							
010550	71	006320	00	000000		016440		RNK								
010560	71	006320	00	000000		016450		TA	6	1						
						016460		SET	0	00014						
						016470		STC	T	CN11						
						016480		SET	0	00045						
						016490		STC	0	00002						
						016500	TBCN11	TC	T	CN7						
						016510		CON	5							
						016520	TBCN12	TCA	T	AL4,R						
						016530		RNK								
						016540		TC	T	CN8						
						016550		RNK								
						016560	TBCN13	TC	E	P						

COMMENTS.
 TO FALL THROUGH
 SET LINE TALLY
 END OF TAG TABLE
 IS THIS ENTRY
 A CHART TITLE
 MOVE TAG TO PAGE
 MOVE PAGE, POX, AND
 AND ASSEMBLY LINE
 NUMBER TO PAGE
 ADJUST AM7 TO.
 NEXT LINE OR PAGE
 ADJUST AM5 TO
 NEXT TABLE ENTRY
 END OF COLUMN
 SWITCH BELOW IS
 SET TO TFR IF
 THIS IS COLUMN 2
 *TBCN9,TFR
 SET AM7 TO TOP
 OF COLUMN T-D
 SET SWITCH TO TFR
 PRINT OR WRITE
 ENTIRE PAGE
 START NEW PAGE
 CHART TITLE ENTRY T
 MORE THAN 6 LINES
 -LEFT IN COLUMN
 IF THIS IS TOP
 OF COLUMN, DO
 NOT LEAVE
 BLANK LINE
 ADVANCE AM7 TO
 LEAVE BLANK LINE
 MOVE CHART TITLE
 TO PAGE
 ADD 1 TO TAG TALLY
 FORCE END COLUMN
 END TABLE OF CONTENTS
 PRINT OR WRITE

NAME	AUTOFLOW	MSB	OP	A	N	B	SEG #	TAG	OP	A ADDRESS	INDEX	4000501	N	B ADDRESS	SEG 04 DATE	111665	PAGE 037	COMMENTS
010560	21	011247	00	005740			016570	RK	#74232324R									FINAL PAGE
010570	34	004665	00	004705			016580	IT										SET PAGE NUMBER
010580	72	004704	00	000000			016590	SC	HRI12									BACK TO ONE
010600	24	004704	00	000000			016600	SET	HRI12,R#1									CHANGE TITLE OF
010610	24	010710	00	010725			016610	SIC	TCN14									MATH HEADER
010620	25	010730	00	006430			016620	TCT	TCN15									MODIFY ENPG SUBRYN
010630	25	010734	00	007110			016630	TCY	TCN15									TO PRINT DIFFERENT
							016640	RK										SUBHEADER
							016650	RIS	LOAD5									LOAD SEGMENT 5
010640	72	001167	00	000000			016660	SET	0:1167									
010650	24	010670	00	010707			016670	STC	0:2									
010660	71	000340	00	000000			016680	TC	000340									
010670	00	007450	00	000000			016690	OO	LOAD5									
010700	00	000005	00	007450			016700	OO	05									
010710							016660	RK										
							016670	CON	14									
010730	01	004777	01	010101			016680	DAC	HRI2,R#1									
010734	01	005000	01	010101			016690	DAC	HRI2									
010740							016700	DAR	ZCO									
011250								CON	ZCO									
011300								CON	ZCO									
								CON	ZCO									
011324								CON	01									
011327								CON	01									
011330								CON	01									
011331								CON	03,R									
011334								CON	03									
011341								CON	03,R									
011344								CON	04									
011350	00	000000	00	000004				CON	04									

J

etp1e

FOR PATCHES

0000000000000000000000

22
76
77
000074
000006
000001
74232324
00000000000000000004

501 FLOW CHART

NAME	AUTOFLOW	A	N	SEQ #	TAG	INDEX	40000501	N	DATE	111655	PAGE	009
MSH	OP					OP	A ADDRESS	SC	CODE		T	CODE
010310	43	003024	00	003024	017280	SC	CODE	PC	2			
010320	61	010340	00	010340	017290	TC	2	TC	2			
010330	71	014670	00	000000	017300	TC	DGD	TC	2			
010340	72	032045	00	000000	017310	SET	MS#	TC	2			
010350	43	003024	00	003024	017320	SC	CODE	TC	2			
010360	61	010400	00	010400	017330	TC	TCB	TC	2			
010370	71	012500	00	000000	017340	RK		TC	2			
010400	01	000000	00	000000	017350	RK		TC	2			
010410	71	007450	00	000000	017360	PKS		TC	2			
010420	45	000173	00	005647	017370	TC	TP1	TC	2			
010430	72	777774	70	600000	017380	RK	97	TC	2			
010440	24	003025	00	003033	017400	RK		TC	2			
010450	22	003022	07	000011	017410	SET	777774	TC	2			
010460	22	003023	07	000012	017420	STC	TAG	TC	2			
010470	44	000173	00	005647	017430	SET	00XND/2	TC	2			
010500	72	000012	70	600000	017440	OCT	00XND/3	TC	2			
010510	24	005600	00	005624	017450	TCA	97	TC	2			
010520	44	000173	00	005647	017460	SET	000012	TC	2			
010530	22	005625	07	000012	017480	TCA	97	TC	2			
010540	22	005625	07	777756	017490	RK	PCN3	TC	2			
010550	71	013710	00	000000	017500	OCT	PCN3	TC	2			
010560	61	010530	00	010530	017510	TC	MVLN	TC	2			
010570	72	000012	70	600000	017520	CTE	PRD1	TC	2			
010600	24	005600	00	005624	017530	SET	000012	TC	2			
010610	71	014320	00	000000	017540	STC	PCN1	TC	2			
010620	71	007450	00	000000	017550	RK	PCN2/R	TC	2			
010630	22	003022	07	000005	017560	RK		TC	2			
010640	22	003023	07	000006	017570	TC	VTLN	TC	2			
010650	72	777773	70	600000	017580	RK		TC	2			
010660	35	003025	00	003033	017590	TC	TP1	TC	2			
010670	71	013370	00	000000	017600	RK		TC	2			
010700	25	004074	00	000133	017610	JCD	00XND/2	TC	2			
010710	71	004170	00	000000	017620	OCT	00XND/3	TC	2			
010720	61	011000	00	011000	017630	SET	777773	TC	2			
010730	22	005014	57	777774	017640	SCR	TAG	TC	2			
010740	22	000015	57	777777	017650	TC	GNCR	TC	2			
010750	22	022046	07	000000	017660	TCT	ALMT	TC	2			
010760	22	005016	57	000001	017670	RK		TC	2			
010770	22	000017	57	000002	017680	TC	LICTAG	TC	2			
011000	44	000173	00	005647	017690	CTC	JCD1	TC	2			
011010	44	000173	00	005647	017700	OCT	000014	TC	2			
011020	72	000013	70	600000	017710	OCT	000015	TC	2			
011030	24	003035	00	003043	017720	OCT	000001	TC	2			
011040	71	007450	00	000000	017730	OCT	000002	TC	2			
011050	72	777773	70	600000	017740	TC	97	TC	2			
011060	35	003025	00	003033	017750	JCD1		TC	2			
011070	22	003022	07	000005	017760	TC	97	TC	2			
011100	22	003023	07	000006	017770	SET	000013	TC	2			
011110	71	013370	00	000000	017780	STC	L T	TC	2			
					017790	TC	T-1	TC	2			
					017800	RK		TC	2			
					017810	SET	777773	TC	2			
					017820	SCR	T-1	TC	2			
					017830	OCT	000005	TC	2			
					017840	OCT	000006	TC	2			
					017850	TC	97	TC	2			

PAGE 009 COMMENTS

Z D
 ILLOGICAL CONDITION
 ILLEGAL CODE
 RESTART
 PROCESS P. CODE
 BACK UP AM7 TO
 LINE ABOVE
 PUT TAB ON LINE
 ABOVE SYMBOL
 PUT BOX # ON LINE
 ABOVE SYMBOL
 PUT UPPER LINE
 ON PROCESS BOX
 ON PAGE LAYOUT
 INCREMENT TO
 GET TO NEXT LINE
 PUT IN SIDES OF
 PROCESS BOX
 LINE OF TEXT TO PAGE
 END OF COMMENT FIELD
 PUT LOWER LINE
 ON PROCESS BOX
 ON PAGE LAYOUT
 PUT IN VERTICAL
 CONNECT LINE
 GET NEW RECORD
 PROCESS J CODE
 PUT BOX NUMBER
 ON PAGE
 PUT TAB OF JUMP
 ON PAGE
 DRAW CIRCLE ON PAGE
 MSC OF DESTINATION
 TAG TO AM3
 FINE TAG IN TABLE
 TAG FOUND IN TABLE
 PUT PAGE AND BOX #
 OF DESTINATION TAG
 INSIDE CIRCLE
 ADJUST AM7 TO
 BOTTOM OF CIRCLE
 PUT DESTINATION
 TAG ON PAGE
 GET NEW RECORD
 PROCESS E. CODE
 PUT TAG OF EXIT
 ON PAGE
 PUT BOX NUMBER
 ON PAGE
 LFA - CIRCLE ON PAGE

SEE OF DATE 111655

INDEX 40000501

OP A ADDRESS

SC CODE

TC 2

ST

CODE

PCD

JCD1

93

70 ST

70 ST

PCD1

70 ST

PCN2/R

07 000005

07 000006

70 ST

70 ST

JCD1

57 777774

57 777777

07 000000

57 000001

57 000002

70 ST

LMT/R

70 ST

LMT/R

70 ST

TAG/R

07 000005

07 000006

```

NAME AUTOFLD*  NAME A B  SF# TAG  FID: X  40007501  SEG 'S' DATE 111665  PAGE 040
RSH  OP  A  B  OP  A ADDRESS  N B ADDRESS  N B ADDRESS  COMMENTS
011120 72 000002 70 600000  SET 0 0002  70 ST  #4475063#  70 ST #4475063#  INSIDE CIRCLE
011130 24 022047 00 022052  STC #4475063#  OCT BKN0#2  07 000007  GET NEW RECORD
011140 71 007450 00 000000  TC TPI  OCT BKN0#3  07 000010  PROCESS B CODE
011150 25 003060 00 000133  RMK  SET TAG  70 ST  TAG#R#1  PUT TAB ON PAGE
011160 44 000133 00 004073  TCA #7  TC BCD1  70 ST  TAG#R#1  AM7 TO NEXT LINE
011170 45 000133 00 004067  TCA #3  TC BCD2  70 ST  TAG#R#1  PUT IN TOP LINE
011200 61 011220 00 011210  TCS #3  TC BCD1  70 ST  TAG#R#1  OF NOTE SYMBOL
011210 71 011220 00 000000  TC BCD2  TC BCD1  70 ST  TAG#R#1  AM7 TO NEXT LINE
011220 72 005020 30 600000  SET HDR2#R 3  30 ST  HDR2#R 3  PUT IN SIDE OF
011230 24 003064 03 003064  STC CR#V  03 COM#  07 000010  NOTE SYMBOL
011240 34 005020 00 005020  SCC HDR2#R 3  SCC HDR2#R 3  07 000010  LINE OF TEXT TO PAGE
011250 34 005021 30 003167  TC TPI  TC TPI  07 000010  END OF COMMENT FIELD
011260 71 007450 00 000000  SCC HDR2#R 3  SCC HDR2#R 3  07 000010  PUT IN BOTTOM LINE
011270 34 005020 00 005167  TC BCD1  TC BCD1  07 000010  OF NOTE SYMBOL
011300 71 011260 00 000000  RMK  TCA #7  07 000010  ADJUST AM7 BACK TO
011310 44 000173 00 022057  RMK  TCA #7  07 000010  COLUMN CENTERLINE
011320 45 000173 00 003647  RMK  TCS #7  07 000010  CONNECT LINE
011330 72 000003 70 600000  RMK  SET 0 0003  70 ST  TAG#R#1  GET NEW RECORD
011340 24 006052 00 006055  STC N0#4  OCT BKN0#2  07 000007  PROCESS M CODE
011350 22 003022 07 000007  RMK  OCT BKN0#3  07 000010  SET AM7 TO LINE
011360 22 003023 07 000010  RMK  SET TAG  70 ST  TAG#R#1  ABOVE SYMBOL
011370 72 777775 70 600000  TCA #7  TC BCD1  70 ST  TAG#R#1  PUT TAB ON PAGE
011400 24 003025 00 003032  RMK  SET 000010  07 000010  ADJUST AM7 TO TOP
011410 44 000173 00 003647  TCA #7  TC BCD2  70 ST  TAG#R#1  LINE OF SYMBOL
011420 72 000010 70 600000  SET MCD1  07 000010  PUT IN TOP LINE
011430 24 006030 00 006050  RMK  TCA #7  07 000010  OF NOTE SYMBOL
011440 44 000173 00 003647  TCA #7  OCT MCD1  07 000010  AM7 TO NEXT LINE
011450 22 006051 07 000010  RMK  OCT MCD1  07 000010  PUT IN SIDE OF
011460 22 006051 07 777770  RMK  OCT MCD1  07 777770  NOTE SYMBOL
011470 71 013710 00 000000  TC MVLN  07 000010  LINE OF TEXT TO PAGE
011500 61 011450 00 011450  CTC MCD1  07 000010  END OF COMMENT FIELD
011510 72 000010 70 600000  SET 000010  70 ST  TAG#R#1  PUT IN BOTTOM LINE
011520 24 006030 00 006050  RMK  TCS #7  07 000010  OF NOTE SYMBOL
011530 45 000173 00 022057  RMK  TC TPI  07 000010  ADJUST AM7 BACK TO
011540 71 014320 00 000000  TC VLEN  07 000010  COLUMN CENTERLINE
011550 45 000173 00 022057  RMK  TCS #7  07 000010  CONNECT LINE
011560 71 007450 00 000000  TC TPI  07 000010  GET NEW RECORD
011570 45 000173 00 003647  RMK  TCS #7  07 000010  PROCESS M CODE
011600 22 003022 07 000005  RMK  OCT BKN0#2  07 000005  SET AM7 TO LINE
011610 22 003023 07 000006  RMK  OCT BKN0#3  07 000006  ABOVE SYMBOL
011620 72 777773 70 600000  SET 777773  70 ST  TAG#R  PUT TAB ON PAGE
011630 35 003025 00 003033  SCR TAG  07 000010  ADJUST AM7 TO TOP
011640 44 000173 00 003647  TCA #7  LNTM,R  LINE OF SYMBOL
011650 72 000003 70 600000  RMK  SET 000003  70 ST  HDR2#R 3  PUT IN TOP LINE
011660 24 006056 00 006064  STC MCD1  LNTM,R  OF MALT SYMBOL
011670 44 000173 00 003647  TCA #7  LNTM,R  ADJUST AM7 TO
011700 22 022060 07 777774  RMK  OCT BCD#  07 777774  NEXT LINE
011710 22 022061 07 000004  RMK  OCT BCD#  07 000004  PUT IN ENDS OF
018430  RMK  OCT BCD#  07 000004  MALT SYMBOL

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NAME	MSN	OP	A	B	INDEX	40000501	SEG 05	DATE	111665	COMMENTS	PAGE 042
					OP	A ADDRESS	N	B ADDRESS			
	012500	25	004044	00	012553					PROCESS Y CODE	T
	012510	45	012553	00	022033		TC02,3			SET START OF	P
	012520	25	000173	00	012643		#000001#,1			COMMENTS FIELD	C
	012530	45	000173	00	022077		TC04,3			SAVE AMT.	P
	012540	25	000173	00	012720		#000016#,1			ADJUST AMT TO	C
	012550	27	777777	00	012720		TXLST			FIRST LINE OF COL	P
	012560	25	000223	00	012553		TXLST			SET MSC OF NEXT	C
	012570	34	000000	77	000000		TC02,3			LINE TO BE MOVED	P
	012600	41	012630	00	012630		.77 000000			MOVE ONE LINE OF	C
	012610	44	000173	00	005647		TC03			TEXT TO PAGE	P
	012620	71	012540	00	000000		LNTH,R			STORE STA AS LME	C
	012630	45	012643	00	005647		LNTH,R			OF NEXT LINE	C
	012640	72	777777	00	600000		LNTH,R			CLEAR ISS FROM	P
	012650	43	022100	00	022100		TC03			FRONT OF LINE	D
	012660	41	007450	00	007450		LNTH,R			ANY MORE LINES	P
	012670	44	000173	00	022077		LNTH,R			ADJUST AMT TO	C
	012700	71	014320	00	000000		TP1			NEXT LINE ON PAGE	J
	012710	71	007450	00	000000		#000016#,1			RECYCLE	C
	012720						0177777701000000			IS THIS AN ADJACENT	D
	012730	25	004070	00	000133					COLUMN	C
	012740	44	000133	00	022103					DROP A VERTICAL LINE	P
	012750	71	013540	00	000000					TO START OF CHAIN	C
	012760	71	013370	00	000000					J	J
	012770	72	013113	00	600000					PROCESS PSEUDO	T
	013000	26	003020	00	003023		43			CONNECTOR	C
	013010	22	013110	07	777776		#000004#,1			SET AMT TO ADDRESS	P
	013020	22	013111	07	777777					FIELD WITHIN RCD	C
	013030	22	022046	07	000000					COMPUTE LOCATION OF	P
	013040	22	022104	00	013111					SVNOL ON PAGE	C
	013050	51	013114	00	022106					DRAW CIRCLE ON PAGE	S
	013060	22	013112	07	000001					MOVE PAGE AND	P
	013070	22	013113	07	000002					BOX # TO W/A	C
	013100	71	007450	00	000000					MOVE PAGE # TO	P
	013110									INSIDE CIRCLE	C
	013120	22	004056	00	013135					ADD ONE TO	P
	013130	17	000000	00	000000					BOX NUMBER	C
	013140	71	013220	00	040000					MOVE BOX # TO	P
	013150	72	001167	00	600000					INSIDE CIRCLE	C
	013160	24	013200	00	013217					GET NEW RECORD	J
	013170	71	000340	00	000000					END OF PASS 3	T
	013200	00	004170	00	000000					RND T/P TAPE	P
	013210	00	000007	00	004170					CROSS REFERENCE	D
										LISTING REQUESTED	C
										BRING IN SEGMENT 7	P
										OF CROSS REFERENCE	J
										LISTING	N

OP RETURN HERE AT END

NAME	AUTOFLO	MSM	OP	A	N	B	SEG #	TAG	INDEX	4000/501	N	SEG 05	DATE	111665	COMMENTS	PAGE 043
013220	71	013240	00	010000			019530	ENDP30	OP A ADDRESS						OFF LINE LISTING	D
013230	22	004057	00	013243			019540	OCY T14	TC E:DP32						WRITE PAGE P	
013240	12	013360	00	000000			019550	LW P:CHMS	OCY T14						CHANGE MESSAGE C	
013250	22	004057	00	013263			019560	LH P:7	OCY T14						WRITE END FILE P	
013260	12	013247	00	000072			019570	RI5 LIADI	LH P:7						TC OUTPUT TAPE P	
013270	72	001167	00	000000			019580	SET 001167	RI5 LIADI						CALL IN CONTROL	P
013300	24	013320	00	013337			019590	STP P:2	SET 001167							
013310	71	000340	00	000000			019600	TC 000340	STP P:2							
013320	00	003000	00	000000			019610	OO LUAC1	TC 000340							
013330	00	000001	11	004260			019620	OO OI	OO LUAC1							
013340	03	000000	00	000000			019630	RMK	OO OI							
013350	71	013270	00	000000			019640	RMK	RMK							
013360							019650	RMK	RMK							
013370	25	000243	00	013460			019660	RMK	RMK							
013400	25	005020	00	002651			019670	RMK	RMK							
013410	25	022067	00	000133			019680	RMK	RMK							
013420	66	013470	00	005651			019690	RMK	RMK							
013430	45	000173	00	005647			019700	GENCR	TCT 8:STP							
013440	45	000173	00	005647			019710	GENCR	TCT 8:STP							
013450	45	000173	00	005647			019720	GENCR	TCT 8:STP							
013460	71	777777	00	000000			019730	GENCR	TCT 8:STP							
013470	72	000003	70	600000			019740	GENCR	TCT 8:STP							
013500	24	005752	33	005760			019750	GENCR	TCT 8:STP							
013510	44	000133	00	008027			019760	GENCR	TCT 8:STP							
013520	44	000173	00	005647			019770	GENCR	TCT 8:STP							
013530	71	013470	00	000000			019780	GENCR	TCT 8:STP							
013540	25	000243	00	013640			019790	GENCR	TCT 8:STP							
013550	25	000003	30	005651			019800	GENCR	TCT 8:STP							
013560	25	005574	00	000173			019810	GENCR	TCT 8:STP							
013570	44	000173	00	005647			019820	GENCR	TCT 8:STP							
013600	66	013650	00	005651			019830	GENCR	TCT 8:STP							
013610	22	000000	30	005653			019840	GENCR	TCT 8:STP							
013620	45	000173	00	022113			019850	GENCR	TCT 8:STP							
013630	66	013670	00	005651			019860	GENCR	TCT 8:STP							
013640	71	777777	00	000000			019870	GENCR	TCT 8:STP							

*NEXT J
 PAGE CHANGE P
 J
 PASS 3 SUBROUTINES
 CREATE CIRCLE ON
 PAGE LAYOUT
 ANY POINTS TO
 MIDPOINT OF TOP
 OF CIRCLE=CHANGED
 TO CENTER OF CIRCLE
 SET EXIT
 8 LINES TO TALLY
 INITIALIZE AM3
 ALL LINES MOVED
 BACK OFF AM7 TO
 CENTER OF CIRCLE
 EXIT
 MOVE FORM LINE TO
 PAGE LAYOUT
 SET ADDRESSES FOR
 NEXT LINE
 RECYCLE
 LOCATE FROM AN
 ADDRESS, PROPER
 POSITION ON PAGE
 AM3 HAS MSC OF
 COL #ADDR FIELD
 AM7 GETS ANSWER
 SET EXIT
 LINE # TO STORAGE
 SET AM7 TO START OF
 PAGE LAYOUT AREA
 EXTRA TOP LINE
 TALLY ON LINE #
 ZEROSGET3#
 AM3 NON HAS MSC
 OF LINE ON LAYOUT
 NOW ADJUST AM7 TO
 START OF COLUMN
 COLUMN # TO TALLY
 TALLY ON COLUMN #
 ZEROSGET4#
 EXIT

NAME	AUTOFIELD	RSV	OP	A	B	SEQ	J	TAG	OP	A ADDRESS	4000	500	REP	OS	DATE	111665	COMMENTS	PAGE	044
013650	44	001173	00	005647		020650		GCT3	TCA	\$7			\$1				STEP AM7 BY 120		
013660	71	013600	00	000000		020670		GCT1	TCA	\$7			\$1				STEP AM7 BY 30		
013670	44	001173	00	022117		020690		GCT3	TCA	\$7			\$1				RMV COMMENT LINE		
013700	71	013630	00	000000		020690		GCT2	TCA	\$7			\$1				FRM I/P RCL TO		
						020700		RK									PAGE LAYOUT, AM7		
						020710		RK									SET TO CENTER		
						020720		RK									OF AREA ON PAGE		
						020730		RK									AM1 IS MSC=1 OF		
						020740		RK									LINE TO BE MOVED		
						020750		RK									AM7 ADJUSTED TO		
						020760		RK									TEXT LINE AT EXIT		
						020770		RK									SET EXIT		
						020780		RK									RCL LGTH TO AM3		
						020790		RK									LOCATE NEXT ISS		
						020800		RK									TO COMM FIELD		
						020810		RK									RHE LINE TO AM3		
						020820		RK									LAST LINE IN FIELD		
						020830		RK									NO MVLNBYTES		
						020840		RK									SET FOR PRN		
						020850		RK									SET LENGTH OF LINE		
						020860		RK									NULL LINE		
						020870		RK									Z STORE AM3		
						020880		RK									DIVIDE LENGTH BY Z		
						020890		RK											
						020900		RK											
						020910		RK											
						020920		RK											
						020930		RK											
						020940		RK											
						020950		RK											
						020960		RK											
						020970		RK											
						020980		RK											
						020990		RK											
						021000		RK											
						021010		RK											
						021020		RK											
						021030		RK											
						021040		RK											
						021050		RK											
						021060		RK											
						021070		RK											
						021080		RK											
						021090		RK											
						021100		RK											
						021110		RK											
						021120		RK											
						021130		RK											
						021140		RK											
						021150		RK											
						021160		RK											
						021170		RK											
						021180		RK											
						021190		RK											
						021200		RK											
						021210		RK											
						021220		RK											
						021230		RK											
						021240		RK											
						021250		RK											
						021260		RK											
						021270		RK											
						021280		RK											
						021290		RK											
						021300		RK											
						021310		RK											
						021320		RK											
						021330		RK											
						021340		RK											
						021350		RK											
						021360		RK											
						021370		RK											
						021380		RK											
						021390		RK											
						021400		RK											
						021410		RK											
						021420		RK											
						021430		RK											
						021440		RK											
						021450		RK											
						021460		RK											
						021470		RK											
						021480		RK											
						021490		RK											
						021500		RK											
						021510		RK											
						021520		RK											
						021530		RK											
						021540		RK											
						021550		RK											
						021560		RK											
						021570		RK											
						021580		RK											
						021590		RK											
						021600		RK											
						021610		RK											
						021620		RK											
						021630		RK											
						021640		RK											
						021650		RK											
						021660		RK											
						021670		RK											

NAME	AUTOFLOW	A	N	B	SEQ #	TAG	OP	A ADDRESS	INDX	40000201	SEG 05	DATE	111685	N	B ADDRESS	COMMENTS
014320	25	000243	00	014370	020640	VTLN	RMK				VTLNX					PAGE 045
014330	25	022057	00	005651	020650	VTLN	TCT	857P			TALI					ABOVE START
014340	44	000173	00	005647	020660	VTLN	TCT	8C00002#1			LNTA,R					SET TALLY
014350	22	005670	07	000000	020670	VTLN	TCA	87			07	000000				ADVANCE ONE LINE
014360	66	014340	00	005651	020680	VTLN	OCT	VTCN			TALI					PUT IN VERY CHAR
014370	71	777777	00	000000	020690	VTLN	TA	VTLN1								EXIT
014401					020700	VTLN	TC	777777								FOLLOWING AREAS
014404					020710	RMK	RMK									USED IN PROCESSING
014405					020720	RMK	RMK									DECISIONS
014406					020730	TAGTA	DAR	3,1								TALLY IN DEC LEG
014411					020740	UNDS	DAR	1								UNDER TAG INDICAT
014415					020750	RSS	DAR	1								RIGHT SIDE SIGNAL
014415					020760	LSS	DAR	1								LEFT SIDE SIGNAL
014420					020770	SV7	DAR	3,1								SAVE AREA-AM7
014430					020780	VTAL	DAR	3,1								MISC TALLY
014437					020790	TPXX	DAR	8,0								TEMP STORAGE
014446					020800	DCN	DAR	8,0								
014451					020810	DCN	CON	7								
014460					020820	CON	CON	7								
014472					020830	CON	CON	3								
014501					020840	CON	CON	7								
014513					020850	CON	CON	7								
014522					020860	CON	CON	7								
014531					020870	CON	CON	7								
014534					020880	CON	CON	3								
014543					020890	CON	CON	7								
014552					020900	CON	CON	7								
014555					020910	CON	CON	3								
014564					020920	CON	CON	7								
014573					020930	CON	CON	7								
014576					020940	CON	CON	7								
014605					020950	CON	CON	3								
014614					020960	CON	CON	7								
014617					020970	CON	CON	7								
014626					020980	CON	CON	3								
014635					020990	CON	CON	7								
014640					021000	CON	CON	7								
014647					021010	CON	CON	3								
014656					021020	CON	CON	7								
014664	01	000010	01	010101	021030	CON	CON	7								
014670	25	000173	00	014411	021040	CON	CON	3								
014700	25	022067	00	000133	021050	CON	CON	7								
014710	25	014664	00	005651	021060	CON	CON	7								
014720	72	000010	70	600000	021070	CON	CON	7								
014730	24	014430	33	014450	021080	CON	CON	7								
014740	44	000133	00	022123	021090	CON	CON	7								
014750	44	000173	00	005647	021100	CON	CON	3								
014760	66	014720	00	005651	021110	CON	CON	7								
014770	72	777777	70	600000	021120	CON	CON	7								
015000	35	003076	00	003102	021130	CON	CON	7								
015010	45	000173	00	005647	021140	CON	CON	7								

LINES=1 TO FORM
 DECISION PROCESSING A
 PROCESS D CODE T
 PRESERVE AM7
 INITIALIZE AM3
 # LINES TO TALLY P
 MOV. LINE OF FORM P
 TO PAGE LAYOUT AREA C
 STEP ADDRESSES
 TO TEXT LINE C
 ALL FORM LINES MOVED C
 PUT LABEL ON P
 EQU'L FRAC. C
 ADJUST AM7 TO LAST P

01010101010101
 # 01150101010101
 # 010101
 # 01010101010115
 # 01010115010101
 # 010101
 # 010101150101
 # 010101011501
 # 010101011501
 # 010101
 # 01011501010101
 # 01010101010101
 # 150101
 # 010115
 # 010101010101
 # 010101010101
 # 15010101150101
 # 010101
 # 01011501010101
 # 010101010101
 # 010101
 # 01010101010101
 # 01010101010101
 # 010101
 # 01150101010101
 # 010101

SV7
 83
 TALI
 70 \$T
 33 DCO#16
 # LNTA,R
 TALI
 70 \$T
 COM#14
 LNTA,R

NAME	AUTOFLOW	HSH	OP	A	N	B	SEQ #	TAG	INDEX	40009501	N	SEG 05	DATE	111645	PAGE 047	COMMENTS	
015540	22	014403	00	015535			021790		RMK							DROP CONNECTOR RIGHT C	
015550	72	016200	00	000000			021800	DCD7	RMK			1-5				YES=DCD1,400	
							021810		SET	DCD11		000000				RIGHT CONNECTOR N	
015560	22	022031	00	016035			021820		RMK							SET DEST TAG SWITCH P	
015570	44	000173	00	022127			021830		SET	RMK		DCD8B5				SIDE OF DEC SYMBOL C	
							021840		TCA	87		#000013#P1				PUT LABEL ON P	
015600	72	000062	70	600000			021850		RMK		70	8T				PAGE C	
015610	24	003071	00	003075			021860		SET	COM#9		LNTH#R				ADJUST AMT TO C	
015620	44	000173	00	005647			021870		STC	COM#5						VERTICAL MIDDLE C	
							021880		TCA	87						PUT IN HORIZONTAL C	
015630	72	000000	70	600000			021890		RMK		70	8T				LINE TO RIGHT C	
015640	24	005671	00	005673			021900		SET	HZCON#2		RSS				INDICATOR TO ON C	
015650	22	022130	00	014405			021910		STC	HZCON#2						CODING BELOW IS C	
							021920		OCT	W408						FOR CONNECTORS C	
							021930		RMK							PUT IN VERTICAL C	
015660	44	000173	00	005647			021940		RMK							LINE TO CONNECTOR P	
015670	22	005670	07	000000			021950		RMK							TOP LINE OF CONN C	
015700	44	000173	00	005647			021960	DCD8	TCA	VTC#N		LNTH#R				DRAW CIRCLE ON PAGE S	
							021970		OCT	VTC#N		07	000000			UNDER INDICATOR ON D	
							021980		TCA	87		LNTH#R				NO=DCD8A#YES# C	
015710	71	013370	00	000000			022000		RMK			1-5				PUT PAGE AND BOX C	
015720	22	014404	00	015735			022010		TC	GENR						INSIDE CIRCLE C	
015730	72	016010	00	000000			022020		OCT	U#D5		57	777776			AM5 SET TO ENTRY C	
015740	22	000014	57	777776			022030		SET	DCD8A						IN TAG TABLE BY C	
015750	22	000015	57	777777			022040		OCT	00015		07	000000			LOC TAG SUBROUTINE C	
015760	22	022046	07	000000			022050		OCT	R164		57	000002			DROP AMT TO BOTTOM P	
015770	22	000016	57	000001			022060		OCT	00016		LNTH#R				CF CIRCLE C	
015780	22	000017	57	000002			022070		OCT	00017		LNTH#R				LEFT-RIGHT SIDE SWITCH C	
016000	22	000017	57	000002			022080	DCD8A	TCA	87		LNTH#R				MOVE DESTINATION TAG P	
016010	44	000173	00	005647			022090		RMK							TO SIDE OF CONNECTOR C	
016020	44	000173	00	005647			022100		TCA	87							
016030	72	016070	00	000000			022110	DCD8B	SET	4							
016040	72	777775	70	600000			022120		SET	777775							
016050	35	003045	00	003053			022130		SCR	R-T							
016060	71	016110	00	000000			022140		TC	D#D9							
016070	72	000011	70	600000			022150		SET	0CD#11							
016100	24	003045	00	003053			022160		STC	R#T							
016110	66	016130	00	014401			022170	DCD9	TA	DCD10							
							022180		RMK								
016120	71	007450	00	000000			022190		TC	TP1							
							022200		RMK								
016130	72	003053	00	600000			022210	DCD10	SET	R#T#R							
016140	26	003034	00	003043			022220		STC	L#T#C							
016150	72	003075	00	600000			022230		SET	COM#9							
016160	24	003064	00	003070			022240		STC	COM#9							
016170	71	015410	00	000000			022250		TC	D#D4							
							022260		RMK								
016200	22	022130	00	016035			022270	DCD11	OCT	H#C#2		DCD8B5					
016210	45	000173	00	022127			022280		TCA	87		#000013#P1					
							022290		RMK								
016220	72	010002	70	600000			022300		SET	00002							
016230	24	003071	00	003075			022310		STC	COM#5		70	8T				
016240	44	000173	00	005647			022320		TCA	87		LNTH#R					
							022330		RMK								
016250	72	000002	70	600000			022340		SET	0 0002							
016260	24	005671	00	005673			022350		STC	H#C#4							
							022360		RMK								

HAVE BOTH TAG FIELDS D
 BEE: PROCESSED C
 GET NEXT RECORD J
 PROCESS 2ND TAG N
 MOVE TAG AND LABEL P
 OF 2ND BRANCH C
 TO OVERLAY FIELDS C
 OF FIRST BRANCH C
 LEFT CONNECTOR J
 SET DEST TAG SWITCH N
 SET AMT TO LEFT P
 SIDE OF DEC SYMBOL C
 PUT LABEL P
 ON PAGE C
 ADJUST AMT TO C
 VERTICAL MIDDLE C
 PUT IN HORIZONTAL P
 LINE TO LEFT C
 GENERATE CONNECTOR A

NAME	AUTOFLOW	WH	OP	A	N	B	SEG N	TAG	INDEX	ADDRESS	N	SEG 05	DATE	111665	PAGE 069	COMMENTS
016750	45	014417	00	022133			022950	OP	A	ADDRESS						DEC SYMBOL AND TOP OF DESTINATION SYN
016760	45	014417	00	022057			022960	RK								ADJUST COUNT TO TWO LINES ABOVE
016770	44	000173	00	005647			022970	RK								TOP OF DEST SYMBOL AM7 TO NEXT LINE
017000	66	017260	00	014415			022980	RK								IS LINE COMPLETE BACK UP ONE LINE
017010	45	000173	00	005647			023000	RK								AND PUT IN LETTER V FOR DOWN ARROW
017020	22	022134	07	000000			023010	RK								ON LINES MON TO 80TH UP AND D
017030	44	000173	00	005647			023020	RK								SET LEFT SIDE INDICATOR TO ON
017040	22	022130	00	014406			023030	RK								IS THERE A TOP FROM CONN IN MAY
017050	72	000017	70	600000			023100	RK								PUT IN HORIZ LINE BACK TO CENTER
017060	43	022046	00	022046			023110	RK								IS THERE A STAR IN THE WAY
017070	61	017130	00	017130			023120	RK								PRINT STAR IN WAY SET SM TO PUT IN
017100	72	000013	70	600000			023130	RK								PUT IN HORIZ LINE BACK TO CENTER
017110	24	005671	00	005704			023140	RK								IS THERE A STAR IN THE WAY
017120	71	016110	00	000000			023150	RK								PRINT STAR IN WAY SET SM TO PUT IN
017130	72	150001	00	600000			023160	RK								PUT IN HORIZ LINE BACK TO CENTER
017140	31	000000	77	000016			023170	RK								IS THERE A STAR IN THE WAY
017150	61	017160	00	017170			023180	RK								PRINT STAR IN WAY SET SM TO PUT IN
017160	22	017217	00	017215			023190	RK								PUT IN HORIZ LINE BACK TO CENTER
017170	72	000016	70	600000			023200	RK								IS THERE A STAR IN THE WAY
017200	24	005671	70	600000			023210	RK								PRINT STAR IN WAY SET SM TO PUT IN
017210	72	017250	00	404000			023220	RK								PUT IN HORIZ LINE BACK TO CENTER
017220	72	017216	00	017215			023230	RK								IS THERE A STAR IN THE WAY
017230	72	000013	70	600000			023240	RK								PRINT STAR IN WAY SET SM TO PUT IN
017240	24	022135	00	022137			023250	RK								PUT IN HORIZ LINE BACK TO CENTER
017250	71	016110	00	000000			023260	RK								IS THERE A STAR IN THE WAY
017260	22	005670	07	000000			023270	RK								PRINT STAR IN WAY SET SM TO PUT IN
017270	71	016770	00	000000			023280	RK								PUT IN HORIZ LINE BACK TO CENTER
017300	25	003017	00	014415			023290	RK								IS THERE A STAR IN THE WAY
017310	45	014417	05	000023			023300	RK								PRINT STAR IN WAY SET SM TO PUT IN
017320	44	014417	00	022133			023310	RK								PUT IN HORIZ LINE BACK TO CENTER
017330	25	014415	00	017621			023320	RK								IS THERE A STAR IN THE WAY
017340	25	000173	00	000133			023330	RK								PRINT STAR IN WAY SET SM TO PUT IN
017350	72	777761	30	600000			023340	RK								PUT IN HORIZ LINE BACK TO CENTER
017360	43	022035	00	022035			023350	RK								IS THERE A STAR IN THE WAY
017370	61	015540	00	015540			023360	RK								PRINT STAR IN WAY SET SM TO PUT IN
017400	66	017560	00	014415			023370	RK								PUT IN HORIZ LINE BACK TO CENTER
017410	25	017621	00	014415			023380	RK								IS THERE A STAR IN THE WAY
017420	72	777767	70	600000			023390	RK								PRINT STAR IN WAY SET SM TO PUT IN
017430	35	003071	00	003075			023400	RK								PUT IN HORIZ LINE BACK TO CENTER
017440	44	000173	00	005647			023410	RK								IS THERE A STAR IN THE WAY
017450	45	000173	00	022113			023420	RK								PRINT STAR IN WAY SET SM TO PUT IN
017460	72	000006	70	600000			023430	RK								PUT IN HORIZ LINE BACK TO CENTER
017470	24	005671	00	005677			023440	RK								IS THERE A STAR IN THE WAY
017500	45	000173	00	005647			023450	RK								PRINT STAR IN WAY SET SM TO PUT IN
017510	66	017600	00	014415			023460	RK								PUT IN HORIZ LINE BACK TO CENTER
017520	44	000173	00	005647			023470	RK								IS THERE A STAR IN THE WAY

NAME	AUTOFLOW	LINE NO	OP	A ADDRESS	SEG #	DATE	PAGE	COMMENTS
017530	22	022130	07	000000		111665	50	PUT IN LETTER A FOR AN UP ARROW C
017540	45	005173	00	005647				AP3 TO LINE ABOVE
017550	71	017040	00	000000				PUT IN ELEMENT OF VERTICAL LINE
017560	45	005133	00	005647				TALLY STORAGE
017570	71	017350	00	000000				SET TO SCAN VERT LINE FOR UP ARROW C
017600	22	005670	07	000000				STEP DOWN 1 LINE COMPARE CHARACTER TO OCTAL 50 C
017610	71	017500	00	000000				RECYCLE
017621								IS VERT LINE GOING TO SAME DESTINATION
017630	25	003017	00	020053				PUT LABEL ON PAGE
017640	44	020053	00	022143				CONNECT DEC SYMBOL TO VERTICAL LINE
017650	25	005173	00	000133				FROM CONNECTORS
017660	44	020053	00	022033				LOCATE ENTRY FOR TAG IN TABLE
017670	44	005133	00	005647				FIND ENTRY IN TABLE CROSS REFERENCE FOR THIS SYMBOL
017700	72	022144	00	000000				CROSS-REF PRESENT
017710	43	777761	33	777761				BACK OFF TO TWO LINES ABOVE SYMBOL CUR AT TOP ADJ CHAIN
017720	61	017740	00	015540				PAGE # OF RECORD TO D PAGE # OF REF
017730	71	017660	00	000000				EXTRACT BOX # OF REF FROM BOX # OF THIS RECORD
017740	72	000023	50	600000				COMPARE DIFF VS 2 2&FRIA, MORE PRG, 10 PRN # REDUNDANT
017750	43	020051	00	020053				
017760	61	015540	00	015540				
017770	72	777767	70	600000				
020000	35	003071	00	003075				
020010	44	005173	00	005647				
020020	72	777767	70	600000				
020030	24	005671	00	005676				
020040	71	016110	00	000000				
020051								
020060	25	004104	00	000133				
020070	71	004170	00	000000				
020100	61	007770	00	007770				
020110	72	000033	50	600000				
020120	43	022035	00	022035				
020130	61	020150	00	020150				
020140	71	007770	00	000000				
020150	45	000173	00	005647				
020160	45	000173	00	005647				
020170	72	000031	50	600000				
020200	43	003020	00	003021				
020210	61	020370	00	020370				
020220	21	022150	00	021123				
020230	72	021122	00	000000				
020240	24	003022	00	003023				
020250	72	021126	00	600000				
020260	24	000032	55	000033				
020270	52	021123	00	021126				
020300	61	020370	00	020370				
020310	52	021123	00	022152				
020320	61	020370	00	020360				
020330	72	777470	70	600000				

NAME	AUT	FILE	F	CP	B	SEQ	TAG	OP	A	4000	501	SEG	DATE	111665	COMMENTS
024700						024700		OP	A	4000	501	SEG	DATE	111665	COMMENTS
024710						024710		SEC	07						
024720						024720		MEM	D						
024730						024730		R/K							
024740						024740		R/K	R, D/C:	TAG	TABLE	TD			
024750						024750		R/K	#, #0X#	FIRST	REF#				
024760						024760		TCT	N#TAG						
024770						024770		TCS	TEMP#						
024780						024780		TCT	N	000001#					
024790						024790		TCT	T-1						
024800						024800		R/K							
024810						024810		TCT	RT#SC						
024820						024820		R/K							
024830						024830		SET	0:0001						
024840						024840		SC	R76#						
024850						024850		CTC	# 1						
024860						024860		SET	0:0031						
024870						024870		CTC	C2A						
024880						024880		SET	0:0000						
024890						024890		SET	0:0007						
024900						024900		SET	0:0001						
024910						024910		SET	0:0013						
024920						024920		SET	0:0014						
024930						024930		TCA	N#TAG#R						
024940						024940		TCA	#3						
024950						024950		R/K							
024960						024960		TCA	#1						
024970						024970		R/K							
024980						024980		TA	C2	CONSTRUCTION	OF	NEW			
024990						024990		R/K	SORT	TAG	TABLE	INTO			
025000						025000		OCT	R72#						
025010						025010		TCT	RT#SC						
025020						025020		TCT	N#TAG						
025030						025030		TCS	TEMP#						
025040						025040		CTC	# 1						
025050						025050		SET	0:0013						
025060						025060		SC	0:0034						
025070						025070		CTC	C3B						
025080						025080		OCT	R71#						
025090						025090		SET	0:0000						
025100						025100		SET	0:0000						
025110						025110		SET	0:0023						
025120						025120		SET	0:0023						
025130						025130		SET	0:0024						
025140						025140		SET	0:0047						
025150						025150		SET	0:0047						
025160						025160		TA	C3D						
025170						025170		R/K							
025180						025180		SET	C3						
025190						025190		R/K							
025200						025200		TC	C4						
025210						025210		TCA	#7						
025220						025220		R/K							
025230						025230		TC	C3A						
025240						025240		DAR	ZU#0						
025250						025250		DAR	ZU						
025260						025260		OCT	T13						
025270						025270									

CROSS REFERENCE LIST P
CONTAINING TAG, PAGE Y
REFERENCE, TO NUMBER C
SET TALLY TO NUMER C
OF TABLE ENTRIES-1 P
SET N TAGS CNTRY TO G P
SET A#1 TO -SC P
OF OLD TABLE P
SET A#3 TO MSC OF P
NEW TABLE C
THIS ENTRY A Z
CHART TITLE C
THIS ENTRY HAVE A
CROSS REFERENCE C
MOVE TAG, PAGE, AND P
BOX # TO NEW TABLE C

ADD 1 TO B TAGS CNTR P
ADJUST A#3 TO NEXT P
ENTRY IN NEW TABLE C
ADJUST A#1 TO NEXT P
ENTRY IN OLD TABLE C
END OF TABLE D
BOX # N N
SET INDICATOR TO 72 P
SET A#1 TO MSC TABLE P
SET TALLY TO B TAGS C
MINUS ONE P
IF LESS THAN 1, END D
COMPARE ENTRY VS C
NEXT ENTRY C
HIGH#C3B,LOW# C

ILLOGICAL P
SET INDICATOR TO 71 P
REVERSE ENTRIES C
IN TABLE C
END OF ITERATION D
IS TABLE IN D
PROPER SEQUENCE Y
SORT IS FINISHED N
ADVANCE TO NEXT J
ENTRY IN TABLE P
TEMPORARY STORAGE C
AVAILABLE IF EXPANSION J
REMAINING ON TAPE P

END OF ITERATION D
IS TABLE IN D
PROPER SEQUENCE Y
SORT IS FINISHED N
ADVANCE TO NEXT J
ENTRY IN TABLE P
TEMPORARY STORAGE C
AVAILABLE IF EXPANSION J
REMAINING ON TAPE P

NAME	AUTOFLOW	A	L	R	SEQ #	TAG	1 ST DIX	4000:501	2 ^D P ADDRESS	SEG 17	DATE	111665	COMMENTS	PAGE 056
005530	72	000001	50	600000	025450		OP	A ADDRESS	50 ST				IS BRANCH TAG	D
005540	43	011016	00	011017	025460		SC	W 101#	R010#				ALL SPACES	C
005550	61	005570	00	005570	025470		CTC	C12	C12				EXIT	E
005560	71	777777	00	000000	025480	C11	TC	777777	ST				MOVE TAG TO	F
005570	72	005513	00	600000	025490	C12	SET	F 1G5#R	55 000004				WORK AREA	C
005580	26	000000	52	000006	025510		TC	F1G5					FINP TAG IN TABLE	P
005610	71	00533#	00	000000	025520		CTC	C11	C11				TAG FOUND	S
005620	61	005540	00	005560	025530		TCA	67	DF1#R				ADJUST AM7 TO FIRST	P
005630	44	000173	00	007473	025540		R#K						REF FOR THIS TAG	C
005640	72	000000	70	600000	025550	C13	SET	010000	70 ST				IS FIRST REFERENCE	D
005650	43	011016	00	011016	025560		SC	W 1#	E01#				AREA BLANK	C
005660	61	005720	00	005720	025570		CTC	C14	C14				FILL IN FIRST	P
005670	72	000003	70	600000	025580		SET	010003	70 ST				REFERENCE BOX	C
005700	26	003020	00	003023	025590		STT	B#X#0	BOX#0,R				REFERENCE THIS	J
005710	71	005540	00	000000	026100	C14	CTC	C11	ST				LAST REFERENCE	V
005720	22	000004	70	005735	026130		SET	C15	REFERENCE FOR THIS				TAG	P
005740	36	000000	33	000007	026140		R#K	HAVE REACHED LAST	33 000007				SET UP NEW LINK	P
005750	72	000003	30	600000	026150		SET	010000	30 ST				FOR THIS REFERENCE	C
005760	26	003020	00	003023	026160		STT	B#X#0	BOX#0,R				PUT ADDRESS OF NEW	P
005770	25	000133	07	000007	026170		TCT	83	07 000007				ENTRY IN OLD ENTRY	C
006000	22	011020	07	000004	026180	C20	CTC	RTLR	C7 000004				INCREMENT AM3 FOR	P
006010	44	000133	00	007463	026190		TCA	83	RTLR				NEXT ENTRY	C
006020	71	005560	00	000000	026110	C15	CTC	C11					SET ADDRESS OF NEXT	J
006030	25	000007	70	000173	026120		TCT	010007	70 ST				REFERENCE IN AM7	P
006040	71	003025	01	010101	026140		TC	C14					N PRINTING OF TABLE	T
006054	01	003045	01	010101	026150	C16	DAC	L#T					REWIND TAPE	P
006060	22	004056	00	006075	026160	C17	DAC	R#T					SET TALLY TO NUMBER	P
006070	17	000000	00	000000	026170		CTC	T#3					OF TABLE ENTRIES=1	C
006100	25	004020	00	003775	026180		R#D	N#TAG	TEMPI				M#C OF TABLE TO AM3	P
006110	45	003777	00	010773	026190		TC	TEMP#R	R000001#1				CLEAR PRINT LINE	P
006120	25	007464	00	000133	026200		TCT	RTM#C	83				SET AM7 TO START OF	P
006130	36	010200	00	010367	026210		SET	CPL#R	CPL#R				ENTRY FOR THIS TAG	C
006140	25	000133	00	000173	026220	C21	R#K	START PROCESSING FOR	ST				SET AM3 TO M#C	P
006150	25	010374	00	000153	026230		TCT	81	85				MOVE PAGE # BOX #	P
006160	72	000002	50	600000	026240		R#K	C#L#N1					AND TAG TO PRINT LINE	C
006170	24	000010	77	000011	026250		SET	010002	50 ST				SET AM7 TO 1ST REF	P
006200	72	000006	50	600000	026260		STC	000010	77 000011				ADJUST AM3 TO FIRST	P
006210	24	000012	77	000013	026270		STC	000006	50 ST				REF ON PRINT LINE	C
006220	72	000021	50	600000	026280		SET	010021	77 000013				PRINT LINE COMPLETED	D
006230	24	000001	77	000007	026290		STC	000001	50 ST				PUT OUT LINE	S
006240	44	000173	00	007473	026300		TCA	87	77 000007				SET AM7 TO 1ST REF	P
006250	44	000153	00	011023	026310		R#K		DF1#R				ADJUST AM3 TO FIRST	P
006260	66	006310	00	006710	026320	C22	TCA	85	R000030#1				REF ON PRINT LINE	C
006270	71	006540	00	000000	026330		TA	C23	GUT4				PRINT LINE COMPLETED	D
006300	71	006540	00	000000	026340		TC	C22					PUT OUT LINE	S
006310	22	000000	75	000000	026410	C23	CTC	010000	75 000000				MOVE THIS REFERENCE	P
006320	22	000001	75	000001	026420		CTC	000001	75 000001				. TO OUTPUT AREA	C

NAME	AUTOFLOW	MSH	OP	A	N	B	SEQ	INDFX	40000501	SEG	07	DATE	111649	PAGE	087
006330	22	011024	05	000002			006330	OP	A ADDRESS						
006340	22	000002	75	000003			026440	OCY	816#	N	B ADDRESS				
006350	22	000003	75	000004			026450	OCY	000002		05	000002			
006360	44	000153	01	011027			026460	OCY	000003		75	000003			
006370	22	000004	70	006405			026470	TCA	95		000010	006405			
006400	72	006440	00	000000			026480	RMK			'0	0 1 5			
006410	71	006500	00	000000			026490	SET	C24		01	01			
006420	46	006640	00	003775			026500	TC	OUT		TEMPI				
006430	71	007210	00	000000			026510	TA	C25						
006440	25	000007	70	000173			026520	RMK							
006450	71	006260	00	000000			026530	TC	ENDCC		70	07			
006460	44	000133	00	007453			026540	TCY	000007		RTL,R				
006470	71	006140	00	000000			026550	RMK							
006500	25	000243	00	006450			026560	TC	C21						
006510	25	010374	00	000153			026570	RMK							
006520	44	000153	00	011023			026580	TCY	85TP						
006530	72	000001	50	000000			026590	TCY	CPLN1						
006540	43	011001	00	011001			026600	TCA	95						
006550	61	006630	00	006560			026610	SET	000001		50	01			
006560	46	006600	00	006720			026620	SC	82#		002#				
006570	71	006730	00	000000			026630	TC	OUT2		01				
006600	71	006640	00	010000			026640	TA	OUT1						
006610	22	004057	00	008625			026650	TC	C2P						
006620	12	010177	00	000000			026660	TC	OUT3						
006630	36	010200	00	010367			026670	OCY	T4		010000				
006640	25	006714	00	006710			026680	OCY	T4		0 1 5				
006650	71	777777	00	000000			026690	LW	CPLN#-1						
006660	02	010200	00	000000			026700	SET	CPLN		CPLN,R				
006670	03	000002	00	000000			026710	TC	0-T4		OUT4				
006680	71	006630	00	000000			026720	PR	CPLN						
006690	01	000014	01	010101			026730	TC	OUT2						
006700	01	000014	01	010101			026740	DAC	000014						
006710	01	000014	01	010101			026750	DAC	00014						
006720	01	000000	01	010101			026760	DAC	000000						
006724	01	000031	01	010101			026770	DAC	000031						
006730	25	000243	00	007100			026780	RMK							
006740	51	007204	00	011031			026790	TCY	85TP		CHPX				
006750	72	007746	00	000000			026800	DA	C-P,R		#742#				
006760	24	007202	00	007203			026810	SET	H-C12,6						
006770	32	007745	00	007766			026820	STC	C-P,R-2						
006780	71	007110	00	010000			026830	Z5	H-C12,5		CHP,R-1				
006790	12	007170	00	007025			026840	TC	C-P2		H-C12,6				
006800	22	004057	00	007045			026850	OCY	T4		010000				
006810	12	007577	00	000000			026860	LW	C-P3		0 1 5				
006820	22	004057	00	007065			026870	LA	H-C1,1-1						
006830	12	007777	00	000000			026880	OCY	T4		0 1 5				
006840	25	006774	00	006720			026890	LA	H-C2,1-1						
006850	71	777777	00	000000			026900	TCY	0-T4		OUT5				
006860	03	000000	00	000000			026910	TCY	0-T4						
006870	02	007800	00	000000			026920	PA	070700						
006880	03	000002	00	000000			026930	PR	H-C1						
006890	02	010100	00	000000			026940	PA	0 0002						
006900	02	010100	00	000000			026950	PR	H C						

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 COMMENTS
 ADVANCE AMS TO NEXT
 O/P AREA FIELD
 IS THERE ANOTHER
 REFERENCE
 PUT OUT LINE
 WAS THIS LAST TAB
 END OF LISTING
 PICK UP ADDRESS OF
 NEXT REFERENCE
 ADVANCE AMS TO NEXT
 TAB ENTRY IN TABLE
 NEXT LINE
 SET EXIT
 RESET AMS TO FIRST
 O/P AREA FIELD
 ANY REFERENCES
 ON THIS LINE
 END OF PAGE
 CHANGE PAGE ROUTINE
 B.P. O ON
 WRITE LINE TO
 OPERLINE LIST TAPE
 CLEAR PRINT LINE
 RESET TALLY
 EXIT
 PRINT LINE
 TALLY-ON ENTRIES ON LINE
 INITIALIZER FOR ABOVE
 TALLY-ON LINES ON PAGE
 INITIALIZER FOR ABOVE
 END OF PAGE ROUTINE
 SET EXIT
 ADD 1 TO PAGE #
 MOVE TO HEADER
 B.P. O ON
 WRITE PAGE CHANGE
 MESSAGE
 WRITE MAIN AND
 SUB HEADERS
 4.
 RESET LINE TALLY
 EXIT
 O/P LINE PAGE CHANGE
 PRINT TALLY AND
 SUB HEADERS

NAME	AUTOFLOW	HSV	CP	F	A	B	INDEX	40000501	N	B	DATE	111665	PAGE	050	COMMENTS
007150	03	000002	00	000000	CC	000000	OP	A	ADDRESS						
007160	71	007070	00	000000	CC	000000	PA	000002							
007170					CHP3		CP	0001							
007200					CHP4		CON	0005							
007210	71	007300	00	000000	ENDCC		RNK								REFERENCE LISTING P
007220	72	011030	00	000000	ENDCC		TC	E	0000						PAGE CHANGE
007230	42	007203	00	007203	SET	000000	SET	E	0000						IF AN ENTRY OF PAGE P
007240	22	007203	00	007255	BS	000000	BS	C	0000						WERE WRITTEN FOR THE C
007250	75	000000	00	000000	CON	000000	CON	C	0000						CROSS REF LISTING C
007260	65	007270	00	007300	CON	000000	CON	0	0000						GIVE AN EXTRA P/C
007270	71	007300	00	000000	TC	000000	SSG	0	0000						FOR EDITING PURPOSES C
007300	75	000000	00	010000	CON	000000	RIS	L	0000						BRING BACK SEGMENT S P
007310	72	001167	00	000000	SET	000000	SET	0	0000						
007320	24	007300	00	007357	STR	000000	STR	0	0000						
007330	71	000300	00	000000	TC	000000	TC	0	0000						
007340	00	007450	00	000000	OC	000000	OC	L	0000						
007350	00	000005	00	013220	OC	000000	OC	0	0000						
007360	25	000203	00	007440	ENDCC		RMK								ENDP300*
007370	71	007430	00	010000	TC	000000	TC	0	0000						PAGE CHANGE
007400	22	006057	00	007415	OC	000000	OC	T	0000						SUBROUTINE
007410	12	007170	00	000000	CON	000000	CON	0	0000						
007420	71	007440	00	000000	TC	000000	TC	0	0000						
007430	03	000000	00	000000	ENDCC		PA	0	0000						
007440	71	777777	00	000000	TC	777777	TC	7	777777						LENGTH NEW TABLE ENTRY
007450	01	000074	01	010101	DAC	000074	DAC	0	000074						LENGTH OLD TABLE ENTRY
007454	01	000034	01	010101	DAC	000034	DAC	0	000034						LENGTH OF EXTRA ENTRIES
007460	01	000010	01	010101	DAC	000010	DAC	0	000010						MISC OF TABLE
011050					DAR	000000	DAR	0	000000						
007464	01	011050	01	010101	DAC	000014	DAC	0	000014						
007470	01	000014	01	010101	DAR	100200	DAR	1	100200						MAIN HEADER
007474					SAC	000001	SAC	0	000001						
007600					CON	000007	CON	0	000007						
007607					CON	000001	CON	0	000001						
007662					CON	000015	CON	1	000015						
007701					CON	000008	CON	0	000008						
007732					SAC	000000	SAC	0	000000						
007760					SAC	000011	SAC	0	000011						
007770					SAC	000012	SAC	0	000012						
007772					SAC	000002	SAC	0	000002						
010000					DAR	100200	DAR	1	100200						
007777					SAC	000001	SAC	0	000001						
					CON	000001	CON	0	000001						

What is claimed is:

For use in a system for automatically controlling a computer having a storage, a processor, an output unit, and control apparatus for controlling the operation of said processor, storage and output unit to perform sequences of operations on blocks of data in the form of coded digital signals stored in said storage and representing successive instructions of various types of a computer program including process, unconditional transfer and conditional branch instructions; means for directing the operation of said control apparatus to process said data blocks sequentially and to produce a record of a flow chart representative of said program, the method comprising:

1. processing said data blocks to identify the types of instructions represented thereby and to establish flow chart symbols therefor;

2. allocating successive ones of said symbols forming a main flow and branch sequences of the computer program as arrays in sections of each of successive flow chart pages;

3. identifying branch sequences of the computer program associated with branch instructions of said main flow sequences and initiating said allocating of successive symbols of said branch sequences to other sections of the associated flow chart pages to form a symbol array with said main flow array; and producing a record of said symbol arrays in successive flow chart pages in accordance with the allocation thereof including producing chart indications of the relationships of branch and main flow symbols.

2. The directing means having the method as recited in claim 1 wherein the identifying of said branch sequences includes identifying sub-branch sequences of a computer program that are associated with branch instructions of said branch sequences and allocating successive symbols of said sub-branch sequences to other sections of the associated flow chart page to form a symbol array with the branch symbol array.

3. The directing means having the method as recited in claim 2 wherein said allocating includes forming clusters of main flow and branch sequences with each cluster being formed as a sequence of main flow symbols and available sequences of branch and sub-branch symbols associated with branch instructions of said main flow and branch sequences, respectively, and allocating said clusters successively to said chart pages.

4. The directing means having the method as recited in claim 1 and further comprising:

a. determining the amount of page space required for each of said symbols;

b. identifying chains of said symbols occurring between successive transfer instructions and determining the amount of page space required for the symbols of each chain;

c. and wherein said branch sequence identifying includes determining in connection with each branch instruction of the main flow whether the page space required for the symbols of the associated branch chain fits within the space of an unallocated section of the associated chart page.

5. The directing means having the method as recited in claim 2 and further comprising:

a. determining the amount of page space required for each of said symbols;

b. identifying chains of said symbols occurring between successive transfer instructions and determining the amount of page space required for the symbols of each chain;

c. wherein said branch sequence identifying further includes determining in connection with each branch instruction of the main flow and branch chains whether the page space required for the symbols of the associated branch and sub-branch chain, respec-

tively, fits within the space of an unallocated section of the associated chart page.

6. The directing means having the method as recited in claim 4 wherein said allocating includes, upon allocating a branch instruction symbol, determining the fit of a branch chain.

7. The directing means having the method as recited in claim 6 wherein said branch symbol allocating is performed upon the allocating of a branch instruction of the main flow sequence, and said main flow symbol allocating continues upon the completion of the branch chain allocation.

8. The directing means having the method as recited in claim 6 wherein said branch symbol allocating is performed upon the completion of a page section of main flow allocating.

9. The directing means having the method as recited in claim 6 wherein said data block processing includes scanning a predetermined field of each of said data blocks to establish different flow chart symbols therefor or to reject said data block, respectively, in accordance with different coded signal groups in said field, and extracting coded signal portions from said field for certain types of instructions; and wherein said record producing includes means for producing indications on said chart representative of the extracted signal portions.

10. For use in a system for automatically controlling a computer having a storage, a processor, an output unit, and control apparatus for controlling the operation of said processor, storage and output unit to perform sequences of operations on blocks of data in the form of coded digital signals in a certain format stored in said storage and representing successive instructions of various types of a computer program including process, unconditional transfer and conditional branch instructions; a method for directing the operation of said control apparatus to process said data blocks sequentially and to produce a record of a flow chart representative of said program, said method comprising:

a. processing said data blocks successively including scanning predetermined fields of each of said data blocks to establish a flow chart symbol therefor or to reject said data block, respectively, in accordance with different coded signal groups in said fields;

b. allocating successive ones of said symbols forming a sequence of the computer program as an array and allocating each of successive ones of said sequences as an array in a section of each of successive flow chart pages;

c. and producing a record of said symbol arrays in successive flow chart pages in accordance with the allocation thereof.

11. The method as recited in claim 10 wherein said scanning of said data block field includes establishing different flow chart symbols in accordance with different coded signal groups in said field.

12. The method as recited in claim 10 wherein said scanning of said data block field includes extracting from said data block field a coded signal portion representing an instructional commentary and extracting therefrom a coded signal portion representing a supplementary commentary to append it to the extracted portion of a preceding data block, respectively, in accordance with different coded signal groups in said field; and wherein said record producing includes producing symbols with representations of said commentary portions set forth therein and of variable size in accordance with the lengths of the commentary portions.

13. The method as recited in claim 10 wherein said scanning of said data block field includes identifying the types of instructions represented by said data blocks and extracting from said data block field for transfer and branch instructions; a coded signal portion identifying the transferee instruction; and wherein said record producing includes producing from said identifying portions

indications of the relationships of the transferor and transferee symbols.

14. The method as recited in claim 10 wherein said scanning of said data block field includes identifying the types of instructions represented by said data blocks and extracting from said data block field for branch instructions a coded signal portion identifying the operational conditions of the branch instruction; and wherein said record producing includes producing from said identifying portion labels of the branch conditions at the associated branch instruction symbol.

15. For use in a system for automatically controlling a computer having a storage, a processor, an output unit, and control apparatus for controlling the operation of said processor, storage and output unit to perform sequences of operations on blocks of data in the form of coded digital signals stored in said storage and representing successive instructions of various types of a computer program including process, unconditional transfer and conditional branch instructions; the method of directing the operation of said control apparatus to process said data blocks sequentially and to produce a record of a two-dimensional flow chart representative of said program, said method comprising:

processing said data blocks to identify the types of instructions represented thereby and to establish flow chart symbols therefor;

forming sequences of successive related symbols and allocating some of said sequences as arrays in sections of successive flow chart pages;

identifying branch sequences associated with branch instructions of others of said sequences and allocating said branch sequences to other sections of the associated flow chart pages to form two-dimensional arrays;

and producing a record of said symbol arrays in successive flow chart pages in accordance with the allocation thereof including producing chart indications of the relationships of branch sequences to the associated branch instruction symbols.

16. The method as recited in claim 15 wherein said sequence forming and allocating allocates the symbols of said sequences in an order corresponding to that of said data blocks and said branch sequence identifying and allocating includes allocating said branch sequence symbols subsequent to the allocation of and on the same page as the associated branch instruction.

17. In a data processing system for producing flow charts of a computer program and having a storage, a processor, an input and an output unit, and control apparatus for controlling the operation of said processor, storage and input and output units to perform sequences of operations on input blocks of data in the form of coded combinations of digital signals stored in said storage and representing successive operations of various types employed in a computer program, including conditional branch instructions;

said control apparatus being operative with said processor and storage for determining the size of page space to display flow chart symbolic representations of said data blocks and of one-dimensional arrays of said symbolic representations in relation to subdivisions of a flow chart page, and operative with said processor and storage for allocating successive ones of said symbolic representations of a one-dimensional array thereof to sections of a flow chart page, and operative with said processor and output unit for producing an output record of pages of said one-dimensional flow chart arrays;

control means for directing the operation of said data processing system to produce an output record of a two dimensional flow chart representative of said computer program, the method of said flow chart control means comprising:

initiating operation of said allocating to allocate

one-dimensional branch sequence arrays having symbolic representations for data blocks referenced by said branch instruction data blocks of parent sequences to unallocated sections of the same pages as the parent one-dimensional arrays that include said branch instruction representations;

producing an output record of successive flow chart pages in accordance with said allocations of said one-dimensional sequence arrays and said one-dimensional branch sequence arrays with chart indications of the relationships of the symbolic representations of the branch sequence arrays to those of the associated branch data block representations in the parent arrays on the same pages; whereby clusters of one-dimensional arrays and associated one-dimensional branch arrays are formed each for display on a page to provide a two-dimensional flow chart display.

18. The method as recited in claim 17 wherein said initiating of said allocating includes initiating allocating of one-dimensional sub-branch sequence arrays having symbolic representations for data blocks referenced by said branch instruction data blocks of parent branch sequences to unallocated sections of the same pages as the parent branch arrays that include said branch instruction representations; and producing records of said pages in accordance with the allocations of said sub-branch arrays and in relation to parent arrays on the same pages; whereby clusters may be formed with sub-branch arrays.

19. For use in a data processing system having a storage, a processor, an input and an output unit, and control apparatus for controlling the operation of said processor, storage and input and output units to perform sequences of operations on input blocks of data in the form of coded combinations of digital signals stored in said storage and representing successive operations of various types of a program including branch operations,

an automatic control method for directing the operation of said data processing system to process said data blocks sequentially and to produce a record of a two-dimensional flow chart representative of said program for reproduction on a rectangular page of a known size in a plurality of similar rectangular columns, said automatic control method comprising:

processing said data blocks to establish flow chart symbolic representations thereof for display within a width corresponding to that of one of said columns, including processing sequences of said data blocks of varying lengths to form one-dimensional arrays of said symbolic representations for display within said column width;

allocating to the columns of said flow chart pages said one-dimensional arrays for parent sequences of said program, including identifying branch ones of said one-dimensional arrays for data block sequences having operations referenced by branch operations of said parent sequences and having lengths that fit within unallocated portions of said columns on the same pages as said parent arrays and allocating the identified branch arrays to said unallocated column portions; and producing an output record of successive flow chart pages in accordance with the allocations of said parent and branch arrays with chart indications of the relations of associated parent and branch arrays on the same pages; whereby clusters of parent and branch arrays are formed for a two-dimensional flow chart display.

20. An automatic control method for a data processing system as recited in claim 19, wherein said allocating includes identifying sub-branch ones of said one-dimensional arrays for data block sequences having operations

erenced by branch operations of said branch sequences having lengths that fit within unallocated portions of l columns on the same pages as said branch arrays and locating the identified sub-branch arrays to said unallocated column portions; and said output record producing includes producing records of said pages in accordance with the allocations of said sub-branch arrays with chart indications of the relations of associated branch sub-branch arrays on the same pages.

1. An automatic control method as recited in claim 1 wherein said data block sequence processing to form n-dimensional arrays includes determining the column length of each of said arrays in terms of lines of a page array for display of the symbolic representations thereof and determining the line length of each of said symbolic representations from the associated data blocks.
2. An automatic control method as recited in claim 1 wherein said data block sequence processing further includes storing in said storage a table of identifiers for l data block sequences including an identifier for the l array lengths determined for the associated one-dimensional arrays; and said allocating includes determining unallocated column lengths, locating within said storage table said identifiers of branch sequences and comparing the column lengths thereof with the unallocated lengths.
3. An automatic control method as recited in claim 1 wherein said sequence processing includes identifying each of said data blocks representative of operations of a certain type and generating said sequence identifiers for each group of sequential data blocks between successive ones of said certain type, whereby data block sequences of varied lengths are automatically formed in accordance with the character of the program.
4. An automatic control method as recited in claim 1 wherein said allocating and said column length coming include allocating a portion of a one-dimensional array to one column and a remaining portion of the array to a succeeding column, and generating and allocating a symbolic representation of a connective relation between l array portions, whereby an array that exceeds the unallocated column length may be broken in portions for display in separate columns.
5. An automatic control method as recited in claim 1 wherein said table storing includes storing a table of identifiers of data blocks in said branch sequences referred to by branch operations of said parent sequences; and allocating further includes supplying to said tag table identifiers of page allocations of the symbolic representations in said branch arrays for the tag data blocks; and said output record producing includes generating and allocating symbolic representations of connective relations between the associated referring and tag symbolic representations in the parent and branch arrays.
6. An automatic control method as recited in claim 1 wherein said connective relation generating includes generating connector line symbols where the referring l tag symbolic representations are on the same page l in suitable relation, and generating and allocating connector symbols at the parent referring symbolic representations and recording thereat identifiers of the page allocations of the associated tag symbolic representations in branch arrays.
7. An automatic control method as recited in claim 1 wherein said allocating further includes supplying to l tag table identifiers of page allocations of parent referring symbolic representations associated with tag symbolic representations; and said output record producing further includes producing a record for each tag data block of page allocations of all associated referring symbolic representations.
8. For use in a data processing system having a

storage, a processor, an input and an output unit, and control apparatus for controlling the operation of said processor, storage and input and output units to perform sequences of operations on input blocks of data in the form of coded combinations of digital signals stored in said storage and representing successive operations of various types of a program including branch operations; an automatic control system for directing the operation of said data processing system to process said data blocks sequentially and to produce a record of a two dimensional flow chart representative of said program for reproduction on a page of a known size, the method of said automatic control system comprising:

processing said data blocks to identify the types of operations represented thereby and to establish flow chart symbols therefor;

identifying sequences of said data blocks and storing in said storage a table of identifiers of said data block sequences, said identifying including determining the size of a one-dimensional array of the symbols for each of said data block sequences in relation to subdivisions of a flow chart page;

allocating successive ones of said symbols for parent and branch sequences of the program as parent and branch-sequence arrays, respectively, in sections of each of successive flow chart pages, including locating within said storage table branch sequences of the program associated with branch operations of said parent sequences and of sizes that fit within unallocated sections of the associated pages and initiating the operation of said allocating means to allocate successive symbols of the associated branch sequence arrays to other sections of the associated flow chart pages to form a two-dimensional array with said parent array;

producing a record of said symbol arrays in successive flow chart pages in accordance with the allocations thereof including producing chart indications of the relationships of branch and parent symbols;

whereby flow chart records are produced with parent and branch arrays on the same pages and corresponding to parent and branch sequences of data blocks in different sequential locations in said program.

29. The method of an automatic control system for a data processing system as recited in claim 28, wherein said sequence identifying includes generating said sequence identifiers for each group of sequential data blocks between successive ones of a certain type of operation, whereby data block sequences are automatically formed in accordance with the program.

30. In a method of controlling a data processing system to produce flow charts of a computer program, said system having a storage, a processor, an input and an output unit, and control apparatus for controlling the operation of said processor, storage and input and output units to perform sequences of operations on input blocks of data in the form of coded combinations of digital signals stored in said storage and representing successive instructions of various types of a computer program, including transfer instructions;

wherein the size of page space for reproducing flow chart symbolic representations of said data blocks and of sequences of said symbolic representations is determined in relation to subdivisions of said flow chart page, successive ones of said symbolic representations of any array thereof are allocated to sections of said flow chart pages, and a record of said pages of arrays is produced by means of said output unit;

the method of directing the operation of said data

processing system to produce a record of a two-dimensional flow chart representative of said computer program, said method comprising:

forming chain sequences of said data blocks;
 allocating chain sequences of said data block representations to certain sections of said pages;
 allocating chain sequences of said data block representations associated with transfer instructions of previously allocated chain sequences to unallocated sections of the same pages;
 and producing a record of successive flow chart pages in accordance with said chain sequence allocations with chart indications of the relationships of the transfer chain sequences to the associated transfer instruction data block representations.

31. The method set forth in claim 30, wherein said step of forming chain sequences includes determining the page space size thereof in relation to said page subdivisions; and wherein said step of allocating transfer chain sequences includes the step of first determining whether the page sizes thereof are within those of the unallocated page sections.

32. The method set forth in claim 30, wherein said transfer instructions of the previously allocated chain sequences are conditional transfer instructions; and wherein said step of forming chain sequences includes forming chain sequences between successive unconditional transfer instructions.

33. The method set forth in claim 32, wherein said step of forming chain sequences includes forming a plurality of shorter chain sequences from a longer chain sequence; and identifying on said pages the sequential relationships of the shorter chain sequences.

34. A data processor method of automatically producing two-dimensional flow charts of operations of a computer program, each of said program operations being represented by an input block of data in the form of coded combinations of digital signals; said method being performed with a digital computer and comprising:

forming signal records of sequence chains of flow chart symbols, each of said chains being representative of a different sequence of said data blocks;
 allocating some of said sequence chains together with associated branch path ones of said sequence chains to portions of the same flow chart pages;
 and identifying interrelationships of said sequence chains and the associated branch path sequences.

35. A data processor method as set forth in claim 34, wherein said step of chain allocating includes the step of allocating sub-branch path sequence chains to portions of the same flow chart pages to which associated branch path sequence chains are allocated.

36. A data processor method of automatically producing two-dimensional flow charts of operations of a computer program, each of said program operations being represented by an input block of data in the form of coded combinations of digital signals; said method being performed with a digital computer and comprising:

forming signal records of sequence chains of flow chart symbols, each of said chains being representative of a different sequence of said data blocks;
 allocating clusters of interrelated ones of said sequence chains to flow chart pages; and identifying interrelationships of said sequence chains in each of said clusters.

37. A data processor method as set forth in claim 36, wherein said cluster allocating step includes allocating a cluster to each flow chart page; and assembling clusters from said sequence chains associated in the same data block sequence chains and in branch and sub-branch sequence chains.

38. A data processor method as set forth in claim 37, wherein said cluster assembling step includes forming a

plurality of shorter sequence chains from a longer sequence chain.

39. A data processor method as set forth in claim 38, wherein said cluster assembling step includes assembling a cluster from at least one sequence chain and at least one associated branch sequence chain of combined page space within the space size of said chart pages.

40. A data processor method as recited in claim 36, wherein said program operations include transfer operations; and the step of forming records of sequence chains includes forming sequence chains as a sequence of flow chart symbols between successive transfer operations.

41. A data processor method as recited in claim 36, and further comprising producing from said signal records a display of said flow chart symbols on successive rectangular pages each comprising a plurality of rectangular columns with each column being of a width suitable for display of various ones of said symbols;

and wherein said step of allocating clusters of sequence chains includes allocating branch path ones of said interrelated sequence chains having lengths that fit within unallocated portions of said columns on the same pages.

42. A data processor method as recited in claim 37, and further comprising producing from said signal records a display of said flow chart symbols on successive rectangular pages each comprising a plurality of rectangular columns with each column being of a width suitable for display of various ones of said symbols;

and wherein said step of assembling clusters includes incorporating in a cluster a sequence chain and associated branch and sub-branch chains having lengths that fit within portions of said columns on the same page.

43. A data processor method as recited in claim 34, wherein said program operations include transfer operations; and the step of forming records of sequence chains includes forming sequence chains as a sequence of flow chart symbols between successive transfer operations.

44. A data processor method as recited in claim 34, and further comprising producing from said signal records a display of said flow chart symbols on successive rectangular pages each comprising a plurality of rectangular columns with each column being of a width suitable for display of various ones of said symbols;

and wherein said step of allocating sequence chains includes allocating those of said associated branch path sequence chains having lengths that fit within unallocated portions of said columns on the same pages.

45. A data processor method as recited in claim 35 and further comprising producing from said signal records a display of said flow chart symbols on successive rectangular pages each comprising a plurality of rectangular columns with each column being of a width suitable for display of various ones of said symbols;

said step of allocating branch and sub-branch sequence chains including the allocating of those of said associated branch and sub-branch chains having lengths that fit within unallocated portions of said columns on the same page.

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