

# Least area of triangles enclosing a circle with radius n

## 1) Examples for $n = 1, 2, 3, 4, 5$

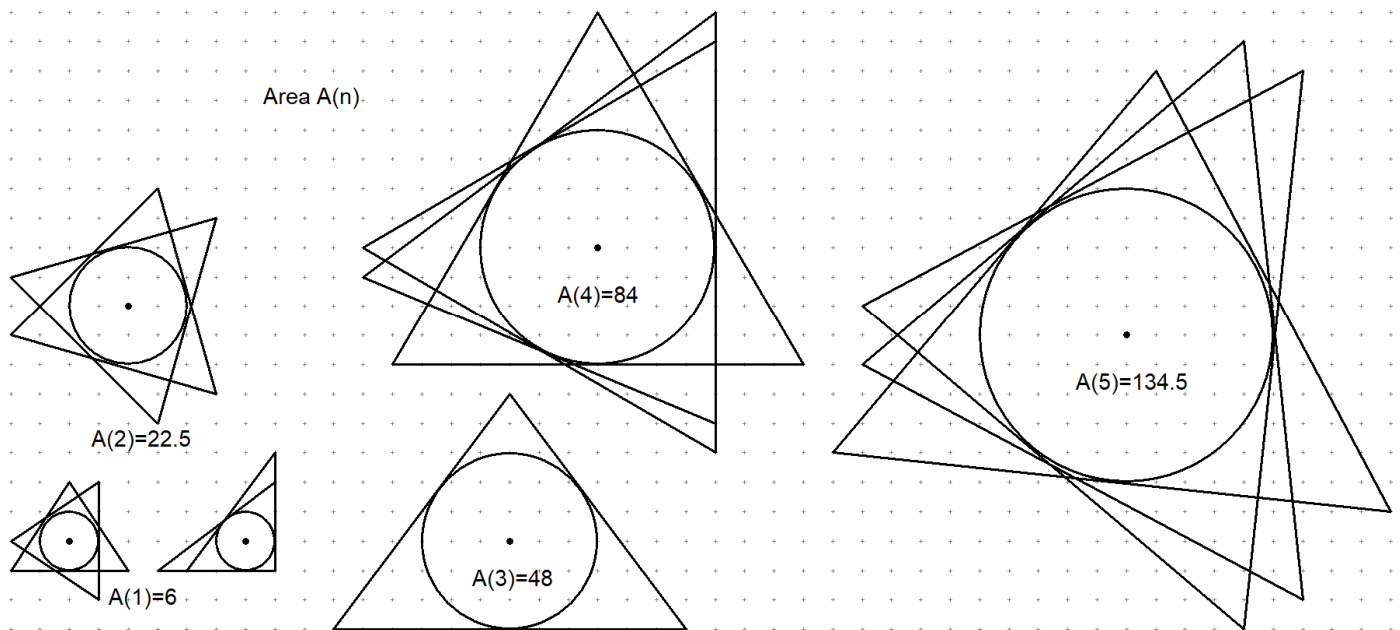


Fig. 1

## 2) Basic ideas of the algorithm

Definition: A triangle enclosing the circle with integer vertex coordinates is a grid triangle.  
 with three tangent edges is a tangent triangle.  
 without further restrictions is a general triangle.

Lemma:

The general triangle with a given vertex and the least area is the symmetrical tangent triangle whose given vertex is located on the symmetry axis. This is evident without proof.

Example  $n=5$  (fig. 2)

$A_{01}B_0C_0$  is an equilateral tangent triangle (height 15 and area  $75 \cdot \sqrt{3} \approx 129.9$ ). This is a lower limit for the least area of a grid triangle.

$A_{01}B_1C_1$  is a symmetrical grid triangle with the area 135.

There are two symmetrical tangent triangles  $A_2B_2C_2$  and  $A_3B_3C_3$  with the same area 135. Their heights are  $h_2 = 8.05$  and  $h_3 = 13.00$ . The area of any symmetrical tangent triangle with height  $h$  and  $h_2 < h < h_3$  is less than 135. The vertex on the symmetry axis of the triangle must be located within the figure  $A_2D_2D_3A_3$ . Because of the symmetry of the grid, the range can be reduced to  $x \geq 0$  and  $x \leq y$ . As a consequence of the lemma, one vertex of the grid triangle with the least area must be located within the figure  $A_2D_2D_3A_3$ . The algorithm yields  $A_4B_4C_4$  as one of the grid triangles with the least area 134.5.

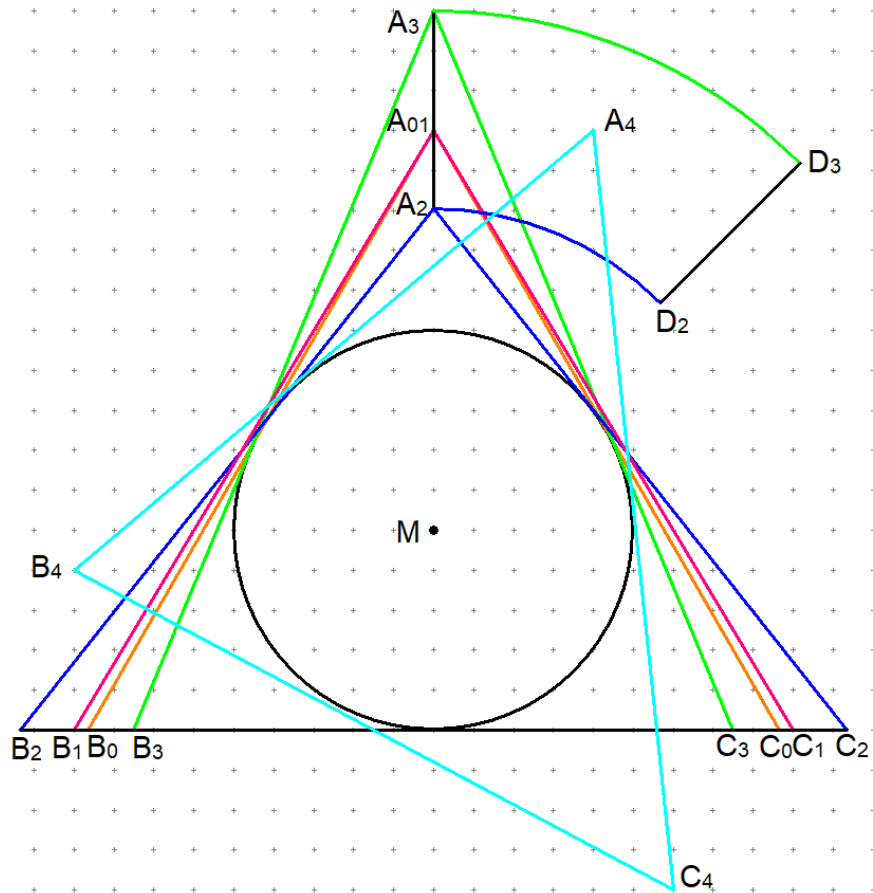


Fig. 2

3) Details of the algorithm for  $n = 5$

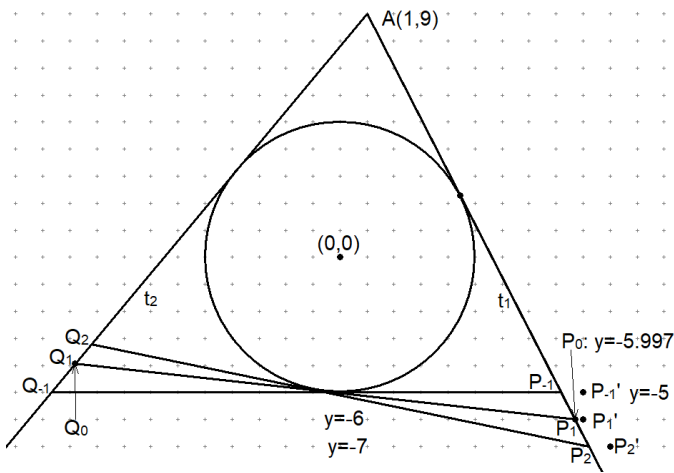


Fig. 3a

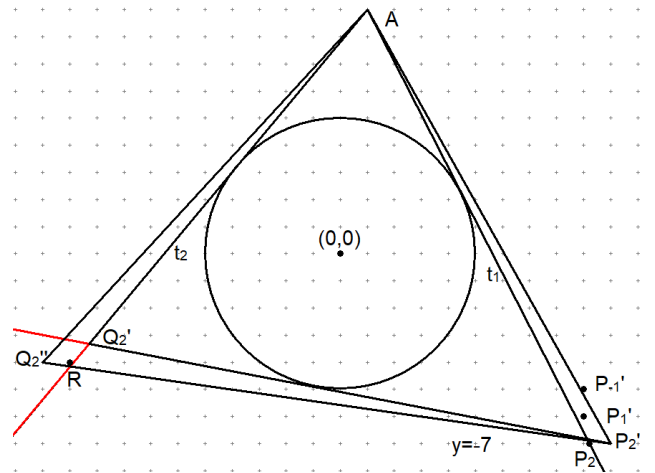


Fig. 3b

a) Select a vertex within  $A_2D_2D_3A_3$  in fig. 2, for example  $A(1,9)$  in fig. 3a.

b) Generate the symmetrical tangent triangle  $AP_0Q_0$  (clockwise, see annotation below). Its area is necessarily less than 135.

- c) Generate another “neighbored” tangent triangle  $AP_kQ_k$ ,  $k = -1,1,2$  for example, such that  $y$  in  $P_k(x, y)$  is an integer. Note that the difference between  $AP_0Q_0$  and  $AP_1Q_1$  is extremely small (“by accident”).
- d) Select  $P_k'(x', y)$  with  $x' = [x]$  (fig. 3a).
- e) The procedure is continued with  $P_2'(10, -7)$ , for example.  $Q_2'$  is the intersection of  $t_2$  and the tangent through  $P_2'$  (fig. 3b).
- f) Check the areas  $|AP_2Q_2| = 131.7$  (fig. 3a) and  $|AP_2'Q_2'| = 137.7 > 135$  (fig. 3b). The area  $|AP_2'Q_2''|$  of any grid triangle is even greater. Therefore, no minimum grid triangle with vertices A and  $P_2'$  exists. So try another point  $P_k$  and stop the search (above and below the point  $P_0$ ) when  $|AP_kQ_k| \geq 135$  and  $|AP_k'Q_k'| \geq 135$ .
- g) For demonstrating the algorithm in fig. 3b, let us assume  $|AP_2'Q_2'| < 135$ . Any point  $Q_2''$  generating a grid triangle  $AP_2Q_2''$  is located within the (red) angle between the tangents through  $P_2'$  and A, for example  $Q_2''(-11, -4)$ . Obviously,  $Q_2'' = R(-10, -4)$  would be a better choice generating a smaller area. The general way of finding the best choice is evident and will not be discussed here.
- h) By selecting  $P_1$  instead of  $P_2$ , the triangle  $(1,9)(9, -6) (-10, -4)$  with the least area 134.5 can be generated. In fig. 3b however, this would be a bad demonstration of the algorithm because the difference between the tangent triangle and the grid triangle is very small.

Annotation:

The tangent  $t_1$  is always steep (slope  $|m_1| > 1.39$ ), whereas  $t_2$  has a wide range of slopes ( $0.15 < m_2 < 3.2$ ). Starting with  $t_2$  instead of  $t_1$  (anticlockwise) would require case distinctions in 4c) such that either  $x$  or  $y$  is an integer in  $P_k(x, y)$ .

#### 4) Formulas (used in fig. 3ab)

a) Slope of the tangent  $t_i$ :  $m_i = \frac{x_0 \cdot y_0 \pm n \cdot \sqrt{x_0^2 + y_0^2 - n^2}}{x_0^2 - n^2}$ ,  $i = 1,2$  for  $x_0 \neq n$

For  $x_0 = n$ :  $m_2 = \frac{y_0^2 - n^2}{2 \cdot n \cdot y_0}$ .  $m_1$  is not defined: Set  $P_k' = P_k$  in fig. 3a.

b)  $P_2(x, y)$  in fig. 3a:  $y = -n \cdot \frac{(r+n)}{y_0}$  for  $x_0 = n$

$$y = \frac{((x_0 \cdot y_0 - m_1 \cdot (n \cdot r + x_0^2)))}{m_1 \cdot y_0 + x_0}$$
 for  $x_0 \neq n$  with  $r = \sqrt{x_0^2 + y_0^2}$

c) Intersection of  $t_1$  and  $y = -k$ :  $x_1 = x_0 - \frac{y_0 + k}{m_1}$

d) Area of the triangle  $(x_0, y_0)(x_1, y_1)(x_2, y_2)$  :

$$\frac{1}{2} \cdot ((x_2 - x_0) \cdot (y_1 - y_0) - (y_2 - y_0) \cdot (x_1 - x_0))$$

5) Appendix

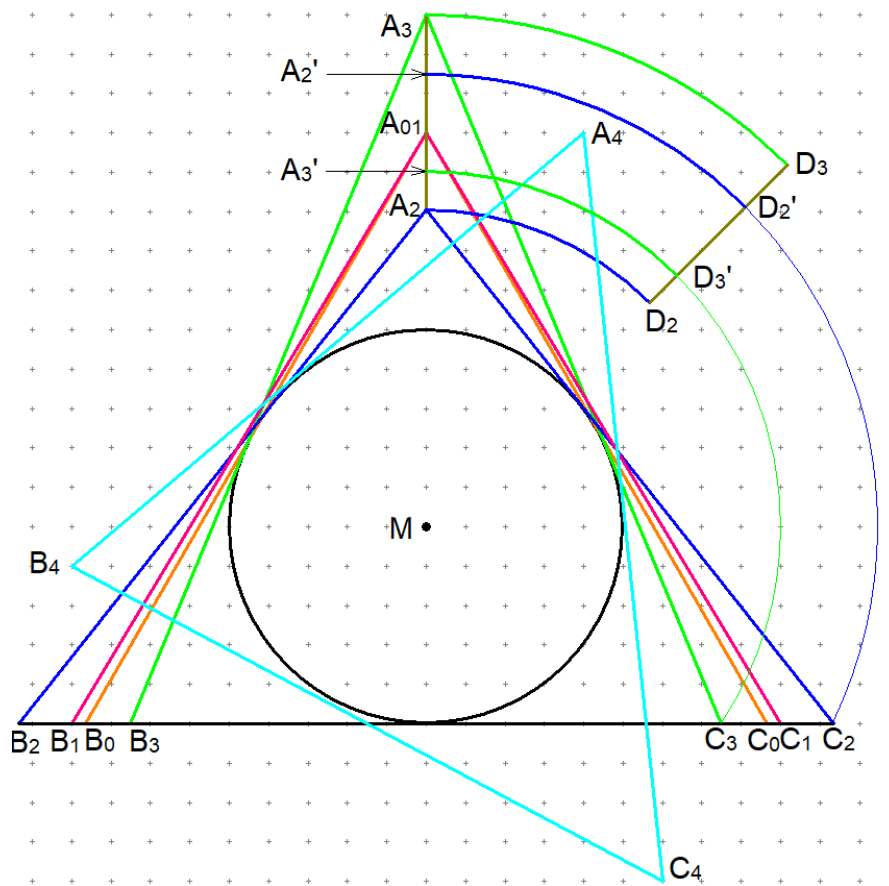


Fig. 4

Fig. 4 is an extension of fig. 2. The circular arcs show  $|MC_2| - |MC_3| < |MA_3| - |MA_2|$ :  
 The distances of baseline vertices from M vary less than those of vertices on the symmetry axis.

Conjecture: There is a minimum grid triangle with a vertex located in  $A_3'D_3'D_2'A_2'$ .  
 The conjecture has been tested for  $n \leq 200$ .