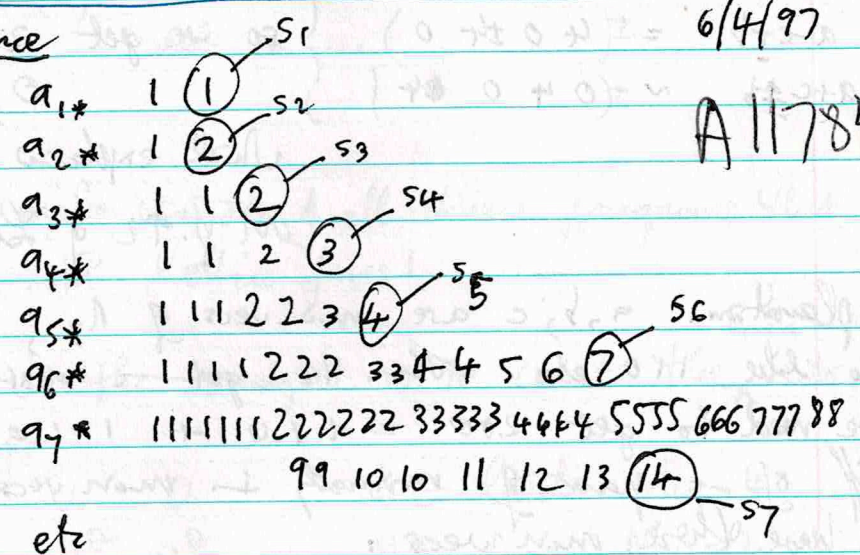


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Lerine's sequence

Define an array



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in which if row a_{6*} is $a_{61} \dots a_{6l}$ then row $7*$ is

$$\{ 1^{a_{6l}} 2^{a_{6,l-1}} 3^{a_{6,l-2}} \dots \}$$

The sequence $s(n)$ is the last term of a_{n*} and runs

- 1 2 2 3 4 7 14 42 a
- 213 2387 175450 139759600
- 9 10 11 12

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and I have 3 more terms beyond that. Note that it has the following properties:

(0) last term of $a_{n*} = s(n)$ (defn)

(1) $l(n) = \text{length of } n^{\text{th}} \text{ row} = s(n+1)$

(2) $\sum_{j=1}^{l(n)} a_{nj} = s(n+2)$

(3) $\sum_{j=1}^{l(n)} (l(n) + 1 - j) a_{nj} = s(n+3)$

(4) No of 1's in $a_{n*} = s(n-1)$ (for $n \geq 2$)

(5) No of 2's in a_{n*} (for $n \geq 5$) = $s(n-2)$

and by following down three rows we get

(6) Let $b_{n*} = \text{row } a_{n*} \text{ read backward } (1 \leq j \leq l(n))$

then

(a) $s(n+4) = \sum_{j=1}^{l(n)} j \left\{ \left(s_{n+2} - \sum_{r=1}^{j-1} b_{nr} \right) b_{nj} - \binom{b_{nj}}{2} \right\}$

The file seqs95 in bin does all this.

To use it: $a1 := [1, 1];$

$a2 := st(a1);$

$a3 := st(a2);$

sum1(a3);

sum2(a3);

sum3(a3);

Here's how to do it in \mathcal{S} (from CLM)
 (Colin Mallows)

Splus

> st <- function(v)

```
function(v)
{
  revv <- rev(v)
  a <- NULL
  for(i in 1:length(v))
    a <- c(a, rep(i, revv[i]))
}
```

reverse

Concatenate

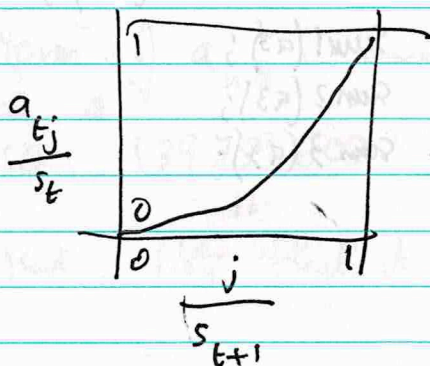
repeat

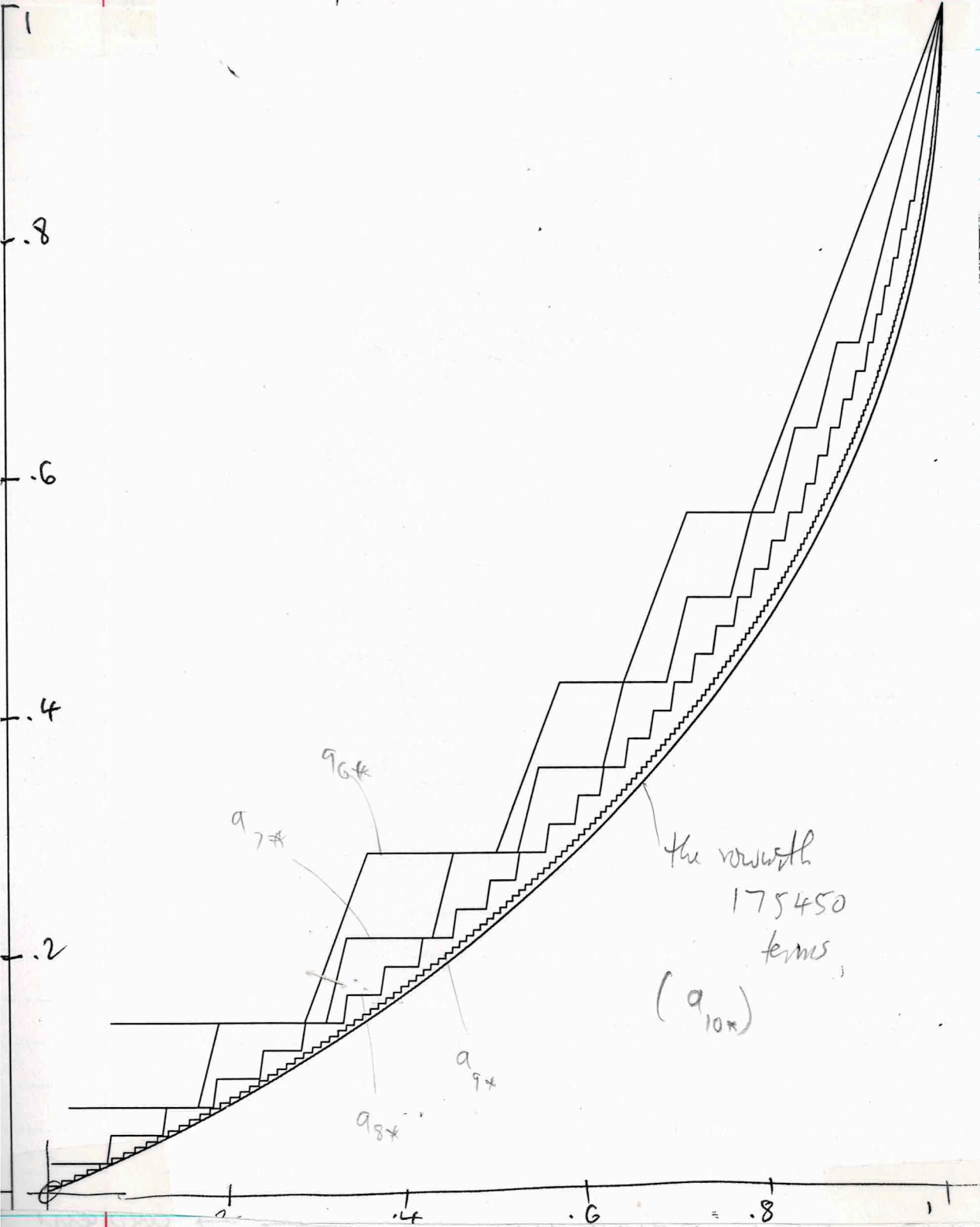
Then

```
v <- c(1,1) (means string "1 1")
for (j in 1:6)
{ v <- st(v) st(v)
  print(v)
}
```

Colin's analysis Look at n^{th} row of table, scale it to have length 1 & height 1. Then the successive rows appear to converge to a smooth curve - see opposite

Call this $g(x)$. Define





1 1 1 (1) 2 2 (2) 3 (3) ... k k (k) ...

look at last k in row t

Then

$$(1) \quad k = a(t, s_{t-1}) + \dots + a(t, s_{t-1} - k + 1)$$

$$\text{Now } g\left(\frac{j}{s_{t+1}}\right) \approx \frac{a(t, j)}{s_t}$$

Then from (1),

$$k = \sum_{j=s_{t-1}-k+1}^{s_{t-1}} a(t, j)$$

and this becomes

$$\frac{k}{s_{t+1} s_t} \approx \sum_{j=s_{t-1}-k+1}^{s_{t-1}} g\left(\frac{j}{s_{t+1}}\right) \frac{\Delta j}{s_{t+1}}$$

$$\approx \int_{s_{t-1}-k+1}^{s_{t-1}} g(x) dx$$

If done correctly, get $\int_{1 - \frac{k-1}{s}}^1 g(x) dx$

Then differentiate to get

$$(2) \quad g(1 - g(x)) g'(x) = \frac{s_{t+1}}{s_t s_{t-1}} \rightarrow \text{constant}$$

$$\therefore \log s_{t+1} = \log s_t + \log s_{t-1} + \cancel{c}$$

$$\therefore \log s_{t+1} + c = \log s_t + c + \log s_{t-1} + c$$

$\therefore A_t = (\log s_t + c)$ satisfies Fig. eqn

$$\therefore \log s_t \sim \phi^t$$

(2) looks a bit nicer if set $y = g(x)$, $x = h(y)$,
get $1-y = h\left(\frac{h'(y)}{h(y)}\right)$

— but still can't solve (1) or (2) exactly.

Q: what is soln to

$$g(1-g(x)) g'(x) = C$$

which is pos & monotonic in \square ?

Cont. p 164

Bjorn Boman on Levine's sequenceCont from page 155

$$\log S_n \sim c \phi^n$$

$$\text{Proof (1)} \quad L_n := \text{left of } n^{\text{th}} = S_{n+1} \quad ; \quad S_n = n^{\text{th}} \text{ term}$$

$$L_{n+1} = \sum_{n+1} = \text{sum of } n^{\text{th}} = S_{n+2}$$

$$\text{claim (1)} \quad L_{n+1} \leq L_n S_n \quad \checkmark \quad (\text{easy})$$

$$\text{claim (4)} \quad \sum_{n+1} \geq \underbrace{1+1+\dots+1}_{S_n} + \underbrace{2+2+\dots+2}_{S_n} + \dots + 2$$

$$\geq \frac{1}{2} \left(\frac{L_{n+1}}{S_n} \right)^2 S_n$$

$$\therefore S_{n+3} \geq \frac{1}{2} \frac{S_{n+2}^2}{S_n}$$

$$\frac{S_{n+3}}{2S_{n+2}} \geq \left(\frac{S_{n+2}}{2S_{n+1}} \right) \left(\frac{S_{n+1}}{2S_n} \right)$$

Hence by $\frac{S_{n+3}}{2S_{n+2}}$, satisfies Fib. recurrence

