

# Lessons of Many-Body Quantum Mechanics for Circuit Design

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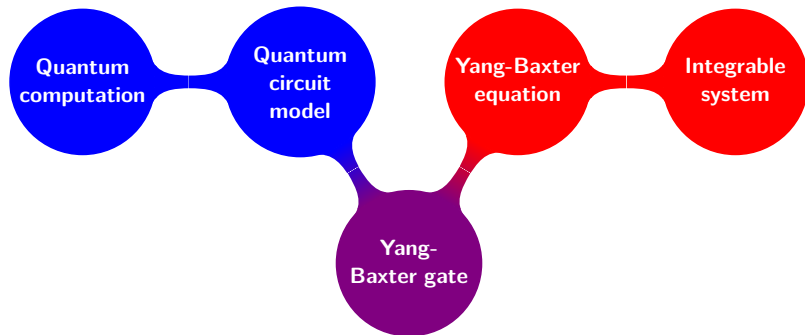
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# Outline

- Two distinguished classes of models in many-body QM. In free models the dynamics is 1-body reducible. In BA solvable the dynamics is 2-body reducible.
- General theories about the two-qubit gate
  - Gate-based decomposition
  - Pulse-based decomposition
- Braid gate and Yang-Baxter gate (YBG)
  - All families of braid gates
  - YBGs from the Yang-Baxterization
- Realization of YBGs on quantum computers
- Conclusions and outlooks

# Background

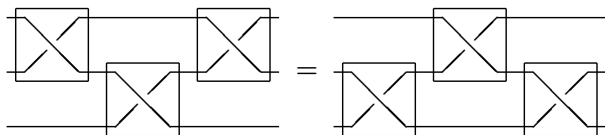
The **Yang-Baxter gate** is the key element connecting the study of **quantum computation** and **integrable system**.



# Yang-Baxter equation

- The **Yang-Baxter equation** (YBE) plays the central role in the study of integrable models.

$$\check{R}_{12}(\mu_1)\check{R}_{23}(\mu_1 + \mu_2)\check{R}_{12}(\mu_2) = \check{R}_{23}(\mu_2)\check{R}_{12}(\mu_1 + \mu_2)\check{R}_{23}(\mu_1).$$



- The Yang-Baxter equation reduces to the **braid group relation**, given by

$$\check{B}_{12}\check{B}_{23}\check{B}_{12} = \check{B}_{23}\check{B}_{12}\check{B}_{23}.$$

# Why Yang-Baxter gate?

- The braid gate is introduced because of the **topological quantum computing** (A. Kitaev, M. H. Freedman, Z.-H. Wang).
- The braid gate is introduced to study the relationship between the topological entanglement and quantum entanglement (S. J. Lomonaco Jr., L. H. Kauffman).
- Quantum simulation of integrable spin chain via the **Suzuki-Trotter decomposition** is realized by YBG. The simulation circuit is integrable (T. Prosen, I. L. Aleiner, Y. Miao).
- The algebraic Bethe ansatz can be applied to study the integrability of **quantum cellular automaton** (B. Pozsgay).
- The properties of many-body nonequilibrium dynamics can be probed from the integrable circuit.

# Trotter decomposition

- Consider the XXX model

$$H_{\text{XXX}} = \frac{J}{2} \sum_{j=1}^{2N} (1 + \vec{\sigma}_j \cdot \vec{\sigma}_{j+1}),$$

with the Pauli vector  $\vec{\sigma}_j$ .

- Consider the Trotter product formula

$$e^{-itH} = \lim_{n \rightarrow \infty} \left( e^{-\frac{itH_{\text{even}}}{n}} e^{-\frac{itH_{\text{odd}}}{n}} \right)^n.$$

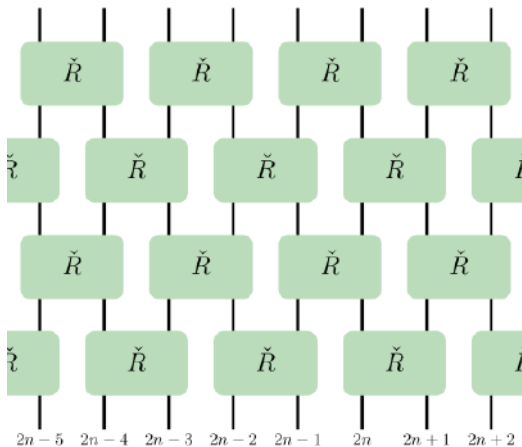
- Then for large  $n$ , the evolution of XXX model can be approximated by applying

$$U_{j,j+1} = \exp \left( -\frac{iJt}{2n} (1 + \vec{\sigma}_j \cdot \vec{\sigma}_{j+1}) \right)$$

on **even and odd sites** interchangeably.

# Trotter circuit

Trotter circuit with a nearest-neighbor interaction. The figure is from [M. Vanicat, L. Zadnik, T. Prosen, arXiv:1712.00431].



# Transfer matrix

- The Trotter circuit has the **transfer matrix**

$$T(\lambda) = \text{Tr}_0 \left( \overset{\leftarrow}{\prod}_{1 \leq j \leq N} R_{0j} \left( \lambda - (-1)^j \frac{\phi}{2} \right) \right),$$

with  $R(\lambda) = P\check{R}(\lambda)$ ,  $\phi = -tJ/n$ , and the right to left order of product [M. Vanicat, L. Zadnik, T. Prosen, arXiv:1712.00431].

- The evolution of the whole circuit is given by  $\mathcal{U}^n(\phi)$ . One layer of evolution is given by

$$\mathcal{U}(\phi) = T \left( -\frac{\phi}{2} \right)^{-1} T \left( \frac{\phi}{2} \right)$$



# Conserved charges

- The circuit has the **conserved charges**

$$Q_n^\pm = \frac{d^n}{d\lambda^n} \log T(\lambda) \Big|_{\lambda=\pm\frac{\phi}{2}}.$$

- For  $n = 1$ , we have

$$Q_1^+ = \sum_{n=1}^{N/2} q_{2n-2, 2n-1, 2n}^{[1,+]}, \quad Q_1^- = \sum_{n=1}^{N/2} q_{2n-1, 2n, 2n+1}^{[1,-]}.$$

with the local charge density

$$q_{1,2,3}^{[1,\pm]} = \frac{i}{2(1+\phi^2)} \left( \vec{\sigma}_1 \cdot \vec{\sigma}_2 + \vec{\sigma}_2 \cdot \vec{\sigma}_3 + \phi^2 \vec{\sigma}_1 \cdot \vec{\sigma}_3 \mp \phi \vec{\sigma}_1 \cdot (\vec{\sigma}_2 \times \vec{\sigma}_3) \right).$$

- At Trotter limit  $\phi \rightarrow 0$ , the charge density returns back to the Hamiltonian density.

# Experiments

Integrable circuit can be well characterized. It is practical to firstly realize the integrable circuit on computers. Then consider the **integrable-breaking** dynamics.

- Google studied the bound state in the XXZ integrable circuit. Then showed that the resilience of the bound states to **integrability breaking** [A. Morvan et al., arXiv:2206.05254].
- The conserved charges have been measured on quantum computers. Due to noises, the conserved charges decay through time [K. Maruyoshi, et al., arXiv:2208.00576].
- The spin transport of integrable circuits is studied on quantum computers. The results show the Kardar-Parisi-Zhang scaling [N. Keenan, et al., arXiv:2208.12243].

# Universal quantum gate

- In classical reversible computation, three-bit gate is required in order to do any computation.
- Remarkably, **two-qubit gates are universal** for quantum computation [D. P. Divincenzo, arXiv:cond-mat/9407022].
- Any two-qubit gate, which can generate quantum entanglement, combined with single-qubit gates, is universal, known as Brylinski's theorem [J. L. Brylinski, R. Brylinski, quant-ph/0108062].
- Common two-qubit gates include

$$\text{CNOT} = |0\rangle\langle 0| \otimes \mathbb{1}_2 + |1\rangle\langle 1| \otimes \sigma_x;$$

$$R_{zz}(\theta) = e^{-\frac{i}{2}\theta\sigma_z \otimes \sigma_z}.$$

- Many-qubit gates.  
Tofoli gates satisfy the tetrahedron equations [multi-dimensional generalization of Yang-Baxter] <https://arxiv.org/pdf/2405.16477>

# Nonlocal parameters

- Any two-qubit gate  $U \in SU(4)$  can be decomposed into [B. Kraus, J. I. Cirac, arXiv:quant-ph/0011050]

$$U = (V_1 \otimes V_2) e^{\frac{i}{2}(a_1 \sigma_x \otimes \sigma_x + a_2 \sigma_y \otimes \sigma_y + a_3 \sigma_z \otimes \sigma_z)} (V_3 \otimes V_4),$$

with  $V_k \in SU(2)$  ( $k = 1, 2, 3, 4$ ). Only the parameters  $[a_1, a_2, a_3]$ , called the **nonlocal parameters**, determine the entangling properties of  $U$ .

- Consider the average entropy generated by  $U$ , called the **entangling power** [P. Zanardi, et al., arXiv:quant-ph/0005031],

$$e_p(U) = \overline{E(U|\psi_1\rangle \otimes |\psi_2\rangle)}_{|\psi_1\rangle \otimes |\psi_2\rangle},$$

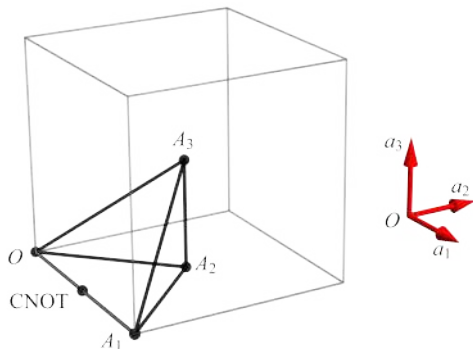
with the linear entropy of two-qubit state  $E(|\Psi\rangle) = 1 - \text{Tr}_1 \rho^2$  and  $\rho = \text{Tr}_2 |\Psi\rangle\langle\Psi|$ .

- The entangling power of two-qubit gate has the analytical form [S. Balakrishnan, R. Sankaranarayanan, arXiv:1005.2467]

$$e_p(U) = \frac{2}{9} (1 - \cos^2 a_1 \cos^2 a_2 \cos^2 a_3 - \sin^2 a_1 \sin^2 a_2 \sin^2 a_3).$$

# Geometric representation

The two-qubit gate has the geometric picture given by the Weyl chamber with  $A_1 = [\pi, 0, 0]$ ,  $A_2 = [\pi/2, \pi/2, 0]$ , and  $A_3 = [\pi/2, \pi/2, \pi/2]$ , called the **two-qubit tetrahedron** [J. Zhang, et al., arXiv:quant-ph/0209120].



# Gate decomposition

- Theorem 1** Any two-qubit gate  $[U] = [a_1, a_2, a_3]$  can be constructed from minimal  $n$  ( $n \geq 3$ ) applications of  $R_{zz}(\theta)$  gate, if the nonlocal parameters satisfies  $0 \leq a_1 + a_2 + a_3 \leq n\theta$  or  $a_1 - a_2 - a_3 \geq \pi - n\theta$  [J. Zhang, et al., arXiv:quant-ph/0308167].
- Theorem 2** A two-qubit gate with the nonlocal parameters  $[U] = [a_1, a_2, a_3]$  can generate the maximal entangled state from the product state iff

$$\frac{\pi}{2} \leq a_j + a_k \leq a_j + a_l + \frac{\pi}{2} \leq \pi,$$

or

$$\frac{3\pi}{2} \leq a_j + a_k \leq a_j + a_l + \frac{\pi}{2} \leq 2\pi.$$

Here  $(j, k, l)$  is the permutation of  $(1, 2, 3)$  with  $j \neq k \neq l$  [J. Zhang, et al., arXiv:quant-ph/0209120].

# Cross-resonance operation

- **Cross-resonance** (CR) is a technique used in the field of quantum computation to implement controlled operations between qubits in a superconducting qubit architecture.
- It relies on the interaction between two qubits, where one qubit is driven at the frequency corresponding to the **resonance frequency** of the other qubit.
- In a typical setup involving two qubits, known as the target qubit and control qubit, the control qubit is driven at its resonance frequency, causing it to undergo Rabi oscillations.

# Cross-resonance operation

- The most common application of cross-resonance is to implement a CNOT gate.
- Cross-resonance operations have been widely utilized in various experimental platforms for quantum computation, particularly in superconducting qubit architectures.
- The IBM superconducting transmon qubits have the cross-resonance operation with the effective Hamiltonian

$$H_{CR} = \frac{\sigma_z \otimes C}{2} + \frac{\mathbb{1}_2 \otimes D}{2},$$

with  $C = \omega_z \mathbb{1}_2 + \omega_{zx} \sigma_x + \omega_{zy} \sigma_y + \omega_{zz} \sigma_z$  and  $D = \omega_{1x} \sigma_x + \omega_{1y} \sigma_y + \omega_{1z} \sigma_z$  [T. Alexander, et al., [arXiv:2004.06755](https://arxiv.org/abs/2004.06755)].

- The Hamiltonian term  $\sigma_z \otimes \sigma_x$  is picked up and other unwanted terms are canceled by additional pulses. CNOT gate is local equivalent to  $e^{i\pi \sigma_z \otimes \sigma_x / 4}$ .



# Rescaled pulses for new two-qubit gates

- Finding the optimal parameters for the pulses is not an easy task. Classical machine learning may be helpful.
- To avoid the heavy calibrations, one can extract new gate from the calibrated pulses [P. Gokhale, et al., arXiv:2004.11205].
- The Gaussian square pulse has the area

$$\alpha = |A| \left( \omega + \sigma \sqrt{2\pi} \operatorname{erf}(n_\sigma) \right),$$

with the amplitude  $A$ , flat-top width  $\omega$ , standard deviation  $\sigma$  of the Gaussian tails, and the truncated number  $n_\sigma$ .

- One can **rescale the area** according to  $\alpha(\theta)/\alpha(\pi/2) = 2\theta/\pi$  in order to realize  $e^{i\theta\sigma_z \otimes \sigma_x/2}$ .

# Braid gates

- Two-qubit gates  $\check{B}$  satisfying

$$(\check{B} \otimes \mathbb{1})(\mathbb{1} \otimes \check{B})(\check{B} \otimes \mathbb{1}) = (\mathbb{1} \otimes \check{B})(\check{B} \otimes \mathbb{1})(\mathbb{1} \otimes \check{B}),$$

are called the **braid gates**.

- All  $4 \times 4$  solutions of braid matrices and braid gates are known [[J. Hietarinta, arXiv:hep-th/9210067](#); [H. A. Dye, arXiv:quant-ph/0211050](#)].
- Most of the solutions can generate entanglement therefore universal [[L. H. Kauffman, S. J. Lomonaco Jr., arXiv:quant-ph/0401090](#)].

# Braid gates

Consider the **eight-vertex-type** two-qubit gates, all solutions of braid gates are

$$\check{B}_I = \begin{pmatrix} e^{i\varphi_1} & 0 & 0 & 0 \\ 0 & 0 & e^{i\varphi_2} & 0 \\ 0 & e^{i\varphi_3} & 0 & 0 \\ 0 & 0 & 0 & e^{i\varphi_4} \end{pmatrix}, \quad \check{B}_{II} = \begin{pmatrix} 0 & 0 & 0 & e^{i\varphi_2} \\ 0 & e^{i\varphi_1} & 0 & 0 \\ 0 & 0 & e^{i\varphi_1} & 0 \\ e^{i\varphi_3} & 0 & 0 & 0 \end{pmatrix}$$

$$\check{B}_{III} = \begin{pmatrix} \cos \varphi_1 & 0 & 0 & \sin \varphi_1 e^{i\varphi_2} \\ 0 & -i \sin \varphi_1 & -\cos \varphi_1 & 0 \\ 0 & -\cos \varphi_1 & -i \sin \varphi_1 & 0 \\ -\sin \varphi_1 e^{-i\varphi_2} & 0 & 0 & \cos \varphi_1 \end{pmatrix},$$

$$\check{B}_{IV} = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 0 & 0 & e^{i\varphi_1} \\ 0 & 1 & 1 & 0 \\ 0 & -1 & 1 & 0 \\ -e^{-i\varphi_1} & 0 & 0 & 1 \end{pmatrix}$$

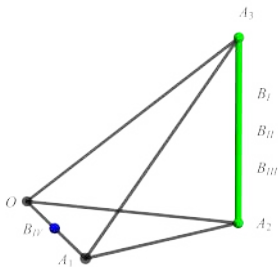
# Geometric representation

- The braid gates have the nonlocal parameters

$$\check{B}_I = \left[ \frac{\pi}{2}, \frac{\pi}{2}, \frac{\pi}{2} - \frac{1}{2}(\varphi_2 + \varphi_3 - \varphi_1 - \varphi_4) \right], \quad \check{B}_{III} = \left[ \frac{\pi}{2}, \frac{\pi}{2}, \frac{\pi}{2} - 2\varphi_1 \right],$$

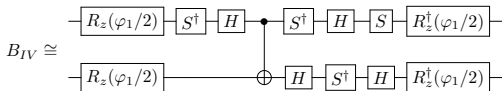
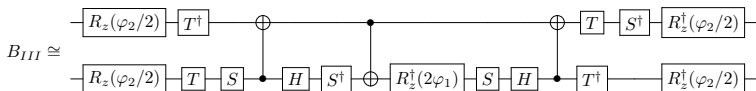
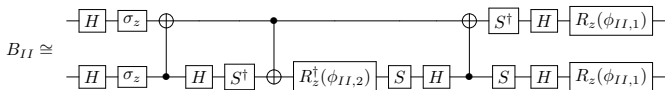
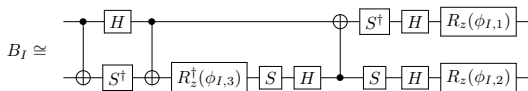
$$\check{B}_{II} = \left[ \frac{\pi}{2}, \frac{\pi}{2}, \frac{\pi}{2} - \frac{1}{2}(2\varphi_1 - \varphi_2 - \varphi_3) \right], \quad \check{B}_{IV} = \left[ \frac{\pi}{2}, 0, 0 \right].$$

- Braid gates  $\check{B}_I$ ,  $\check{B}_{II}$ , and  $\check{B}_{III}$  belong to the same group of two-qubit gates.  $\check{B}_{IV}$  is equivalent to CNOT up to single-qubit gates.



# Decompositions of braid gates

Here we define  $\phi_{I,1} = \frac{1}{2}(-\varphi_1 - \varphi_2 + \varphi_3 + \varphi_4)$ ,  $\phi_{I,2} = \frac{1}{2}(-\varphi_1 + \varphi_2 - \varphi_3 + \varphi_4)$ ,  $\phi_{I,3} = \frac{1}{2}(-\varphi_1 + \varphi_2 + \varphi_3 - \varphi_4)$ ,  $\phi_{II,1} = \frac{1}{2}(-\varphi_2 + \varphi_3)$ , and  $\phi_{II,2} = \frac{1}{2}(-\varphi_2 + 2\varphi_1 - \varphi_3)$ .



# Yang-Baxterization

- Yang-Baxterization is a mathematical procedure to **construct the solution of YBE** from solution of braid relation proposed by [M. Jimbo, *Lett. Math. Phys.* 10, 63-69 (1985)]. Later it is generalized by [Y. Cheng, M. L. Ge, K. Xue, *Comm. Math. Phys.*, 136, 195-208 (1991)].
- Suppose that the Braid operator has the spectral decomposition

$$\check{B} = \sum_{j=1}^N \lambda_j O_j,$$

with  $N$  distinct eigenvalues  $\lambda_j$ . The **ansatz** of  $\check{R}$  matrix is given by

$$\check{R}(x) = \sum_{j=1}^N \rho_j(x) O_j,$$

$$\rho_1(x) = \left(x + \frac{\lambda_1}{\lambda_2}\right) \left(x + \frac{\lambda_2}{\lambda_3}\right) \cdots \left(x + \frac{\lambda_{N-1}}{\lambda_N}\right);$$

$$\rho_2(x) = \left(1 + x \frac{\lambda_1}{\lambda_2}\right) \left(x + \frac{\lambda_2}{\lambda_3}\right) \cdots \left(x + \frac{\lambda_{N-1}}{\lambda_N}\right);$$

...

$$\rho_N(x) = \left(1 + x \frac{\lambda_1}{\lambda_2}\right) \left(1 + x \frac{\lambda_2}{\lambda_3}\right) \cdots \left(1 + x \frac{\lambda_{N-1}}{\lambda_N}\right).$$

# Yang-Baxterization

- If the  $\check{B}$  matrix has two distinct eigenvalues, such as  $\check{B}_{IV}$ , then we have

$$\check{R}(x) = \frac{1}{\lambda_2} \left( \check{B} + x\lambda_1\lambda_2\check{B}^{-1} \right).$$

- If the  $\check{B}$  matrix has three distinct eigenvalues, such as  $\check{B}_I(\varphi_1 = \varphi_4)$ ,  $\check{B}_{II}$ , and  $\check{B}_{III}$ , then we have

$$\check{R}(x) = \alpha(x)\check{B} + \beta(x)\mathbb{1} + \gamma(x)\check{B}^{-1},$$

The coefficients  $\alpha(x)$ ,  $\beta(x)$ , and  $\gamma(x)$  are given by

$$\alpha(x) = -\frac{1}{\lambda_3}(x-1), \quad \beta(x) = \left( 1 + \frac{\lambda_1}{\lambda_2} + \frac{\lambda_1}{\lambda_3} + \frac{\lambda_2}{\lambda_3} \right) x, \quad \gamma(x) = \lambda_1 x(x-1).$$

It may not give the right  $R$  matrix. The permutations  $(\lambda_1, \lambda_2, \lambda_3) \rightarrow (\lambda_2, \lambda_1, \lambda_3)$  and  $(\lambda_1, \lambda_2, \lambda_3) \rightarrow (\lambda_1, \lambda_3, \lambda_2)$  may give different  $R$  matrices.

# Yang-Baxter gates from $\check{B}_I$ including XXZ

- Define the parameters

$$\phi = \frac{1}{2}(\varphi_2 + \varphi_3) - \varphi_1, \quad \alpha = \frac{1}{2}(\varphi_2 - \varphi_3), \quad x = e^\eta.$$

- The braid gate  $\check{B}_I$  with  $\varphi_1 = \varphi_4$  gives the Yang-Baxter gates

$$\check{R}_{I,1}(\eta) \cong \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & \frac{\sin \phi}{\sin(\phi - i\eta)} & \frac{-ie^{i\alpha} \sinh \eta}{\sin(\phi - i\eta)} & 0 \\ 0 & \frac{-ie^{-i\alpha} \sinh \eta}{\sin(\phi - i\eta)} & \frac{\sin \phi}{\sin(\phi - i\eta)} & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}.$$

$$\check{R}_{I,2}(\eta) \cong \begin{pmatrix} \sinh\left(\frac{1}{2}(\eta + i\phi)\right) & 0 & 0 & 0 \\ 0 & 0 & e^{i\alpha} \sinh\left(\frac{1}{2}(\eta - i\phi)\right) & 0 \\ 0 & e^{-i\alpha} \sinh\left(\frac{1}{2}(\eta - i\phi)\right) & 0 & 0 \\ 0 & 0 & 0 & \sinh\left(\frac{1}{2}(\eta + i\phi)\right) \end{pmatrix},$$

$$\check{R}_{I,3}(\eta) \cong \begin{pmatrix} \cosh\left(\frac{1}{2}(\eta + i\phi)\right) & 0 & 0 & 0 \\ 0 & 0 & e^{i\alpha} \cosh\left(\frac{1}{2}(\eta - i\phi)\right) & 0 \\ 0 & e^{-i\alpha} \cosh\left(\frac{1}{2}(\eta - i\phi)\right) & 0 & 0 \\ 0 & 0 & 0 & \cosh\left(\frac{1}{2}(\eta + i\phi)\right) \end{pmatrix}.$$



# Yang-Baxter gates from $\check{B}_{II}$

- Define the parameters

$$\phi = \frac{1}{2}(\varphi_2 + \varphi_3) - \varphi_1, \quad \alpha = \frac{1}{2}(\varphi_2 - \varphi_3), \quad x = e^\eta.$$

- The braid gate  $\check{B}_{II}$  gives the Yang-Baxter gates

$$\check{R}_{II,1}(\eta) \cong \begin{pmatrix} \frac{\sin \phi}{\sin(\phi - i\eta)} & 0 & 0 & \frac{-ie^{i\alpha} \sinh \eta}{\sin(\phi - i\eta)} \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ \frac{-ie^{-i\alpha} \sinh \eta}{\sin(\phi - i\eta)} & 0 & 0 & \frac{\sin \phi}{\sin(\phi - i\eta)} \end{pmatrix}$$

$$\check{R}_{II,2}(\eta) \cong \begin{pmatrix} 0 & 0 & 0 & e^{i\alpha} \sinh\left(\frac{1}{2}(\eta - i\phi)\right) \\ 0 & \sinh\left(\frac{1}{2}(\eta + i\phi)\right) & 0 & 0 \\ 0 & 0 & \sinh\left(\frac{1}{2}(\eta + i\phi)\right) & 0 \\ e^{-i\alpha} \sinh\left(\frac{1}{2}(\eta - i\phi)\right) & 0 & 0 & 0 \end{pmatrix}$$

$$\check{R}_{II,3}(\eta) \cong \begin{pmatrix} 0 & 0 & 0 & e^{i\alpha} \cosh\left(\frac{1}{2}(\eta - i\phi)\right) \\ 0 & \cosh\left(\frac{1}{2}(\eta + i\phi)\right) & 0 & 0 \\ 0 & 0 & \cosh\left(\frac{1}{2}(\eta + i\phi)\right) & 0 \\ e^{-i\alpha} \cosh\left(\frac{1}{2}(\eta - i\phi)\right) & 0 & 0 & 0 \end{pmatrix}$$

# Yang-Baxter gates from $\check{B}_{III}$

The braid gate  $\check{B}_{III}$  gives the Yang-Baxter gates

$$\check{R}_{III,1}(\eta) \cong \begin{pmatrix} \frac{\cosh \eta \cos \varphi_1}{\cosh(\eta + i\varphi_1)} & 0 & 0 & -e^{i\varphi_2} \frac{\sinh \eta \sin \varphi_1}{\cosh(\eta + i\varphi_1)} \\ 0 & i \frac{\cosh \eta \sin \varphi_1}{\sinh(\eta + i\varphi_1)} & -\frac{\sinh \eta \cos \varphi_1}{\sinh(\eta + i\varphi_1)} & 0 \\ 0 & -\frac{\sinh \eta \cos \varphi_1}{\sinh(\eta + i\varphi_1)} & i \frac{\cosh \eta \sin \varphi_1}{\sinh(\eta + i\varphi_1)} & 0 \\ e^{-i\varphi_2} \frac{\sinh \eta \sin \varphi_1}{\cosh(\eta + i\varphi_1)} & 0 & 0 & \frac{\cosh \eta \cos \varphi_1}{\cosh(\eta + i\varphi_1)} \end{pmatrix}$$

$$\check{R}_{III,2}(\eta) \cong \begin{pmatrix} \sinh \eta \cos \varphi_1 & 0 & 0 & e^{i\varphi_2} \cosh \eta \sin \varphi_1 \\ 0 & i \cosh \eta \sin \varphi_1 & -\sinh \eta \cos \varphi_1 & 0 \\ 0 & -\sinh \eta \cos \varphi_1 & i \cosh \eta \sin \varphi_1 & 0 \\ e^{-i\varphi_2} \cosh \eta \sin \varphi_1 & 0 & 0 & \sinh \eta \cos \varphi_1 \end{pmatrix}$$

$$\check{R}_{III,3}(\eta) \cong \begin{pmatrix} \cosh \eta \cos \varphi_1 & 0 & 0 & -e^{i\varphi_2} \sinh \eta \sin \varphi_1 \\ 0 & i \sinh \eta \sin \varphi_1 & -\cosh \eta \cos \varphi_1 & 0 \\ 0 & -\cosh \eta \cos \varphi_1 & i \sinh \eta \sin \varphi_1 & 0 \\ e^{-i\varphi_2} \sinh \eta \sin \varphi_1 & 0 & 0 & \cosh \eta \cos \varphi_1 \end{pmatrix}$$

# Yang-Baxter gates from $\check{B}_{IV}$

- Define the parameter  $\theta$

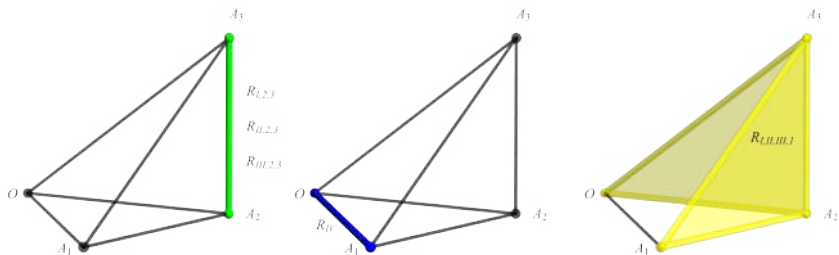
$$\tan(\theta - \pi/4) = x.$$

- The braid gate  $\check{B}_{IV}$  gives the unique Yang-Baxter gate

$$\check{R}_{IV}(\theta) = \begin{pmatrix} \cos \theta & 0 & 0 & e^{i\varphi_1} \sin \theta \\ 0 & \cos \theta & \sin \theta & 0 \\ 0 & -\sin \theta & \cos \theta & 0 \\ -e^{-i\varphi_1} \sin \theta & 0 & 0 & \cos \theta \end{pmatrix}.$$

# Geometric representation

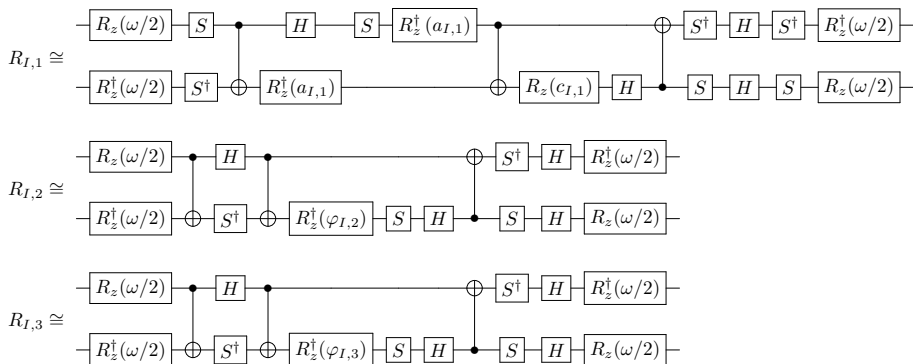
- The obtained Yang-Baxter gates  $\check{R}_{I,1,2,3}$ ,  $\check{R}_{II,1,2,3}$ ,  $\check{R}_{III,1,2,3}$ , and  $\check{R}_{IV}$  can be categorized as three families.



- Based on Theorem 1, the geometric representation can give the optimal gate decompositions.

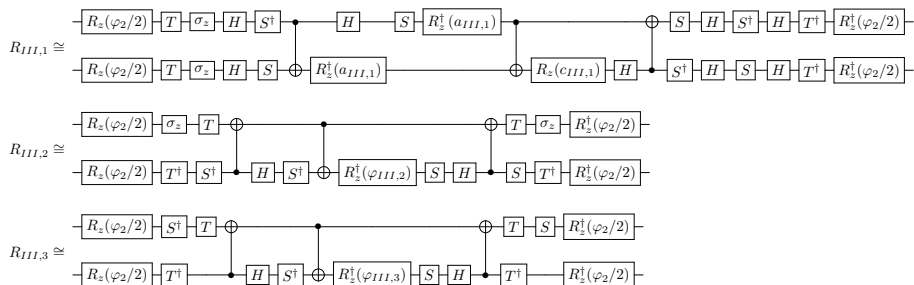
# Decompositions of Yang-Baxter gates

Here we define  $\omega = \frac{1}{2}(\varphi_2 - \varphi_3)$ . The parameters  $a_{I,1}$ ,  $c_{I,1}$ ,  $\varphi_{I,2}$ , and  $\varphi_{I,3}$  are nonlocal parameters (tetrahedron parameters).



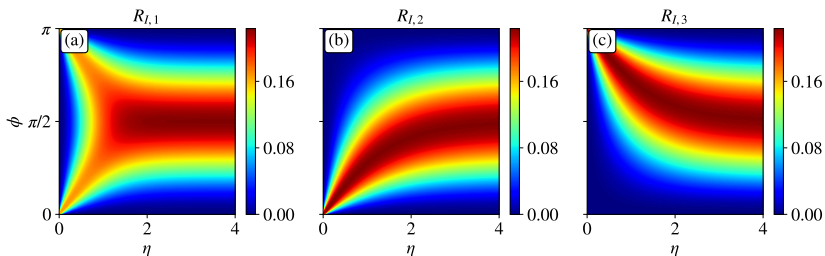
# Decompositions of Yang-Baxter gates

Here the parameters  $a_{III,1}$ ,  $c_{III,1}$ ,  $\varphi_{III,2}$ , and  $\varphi_{III,3}$  are nonlocal parameters (tetrahedron parameters).



# Entangling power

- The braid gate parameters  $\varphi_j$  and the spectral parameter  $\eta$  both determine the entangling power of the YBGs.



- The entangling power of  $\check{R}_{II}$  and  $\check{R}_{III}$  are similar with  $\check{R}_I$ .

## Two examples

- Consider the YBG  $\check{R}_{IV}$  obtained from the braid gate  $\check{B}_{IV}$ .

$$\check{R}_{IV}(\theta) = \begin{pmatrix} \cos \theta & 0 & 0 & e^{i\varphi_1} \sin \theta \\ 0 & \cos \theta & \sin \theta & 0 \\ 0 & -\sin \theta & \cos \theta & 0 \\ -e^{-i\varphi_1} \sin \theta & 0 & 0 & \cos \theta \end{pmatrix}.$$

- Consider the YBG  $\check{R}_{XXX}(\phi)$  (belonging to  $\check{R}_{I,1}$ ) from the Trotterization of XXX model with  $\tan \phi = \mu$

$$\check{R}_{XXX}(\phi) = \frac{1 + i\mu P}{1 + i\mu} \cong \begin{pmatrix} e^{i\phi} & 0 & 0 & 0 \\ 0 & \cos \phi & i \sin \phi & 0 \\ 0 & i \sin \phi & \cos \phi & 0 \\ 0 & 0 & 0 & e^{i\phi} \end{pmatrix}.$$



# Entangling properties of YBGs

- The entangling power of  $\check{R}_{IV}(\theta)$  is

$$e_p(\check{R}_{IV}(\theta)) = \frac{2}{9} \sin^2(2\theta).$$

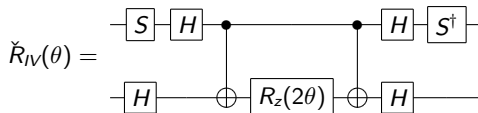
- The entangling power of  $\check{R}_{XXX}(\phi)$  is

$$e_p(\check{R}_{XXX}(\phi)) = \frac{1}{6} \sin^2(2\phi).$$

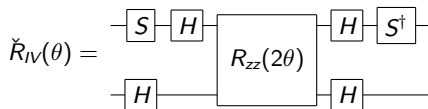
- When  $\theta = \pi/4$  or  $\theta = 3\pi/4$ ,  $\check{R}_{IV}(\theta)$  is the perfect entangler.
- When  $\phi = \pi/4$  or  $\phi = 3\pi/4$ ,  $\check{R}_{XXX}(\phi)$  is the perfect entangler.

# Gate realization of $\check{R}_{IV}(\theta)$

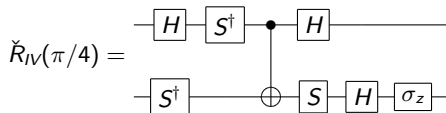
- The YBG  $\check{R}_{IV}(\theta)$  requires minimal two CNOTs (superconducting quantum computers)



- The YBG  $\check{R}_{IV}(\theta)$  requires minimal one  $R_{zz}$  (trapped-ion quantum computers)

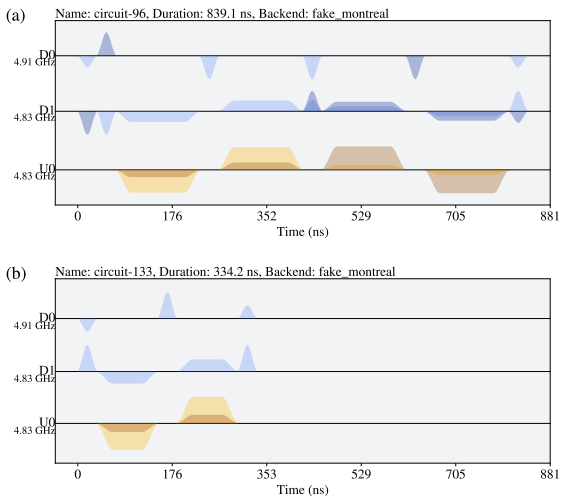


- When  $\theta = \pi/4$ , we have



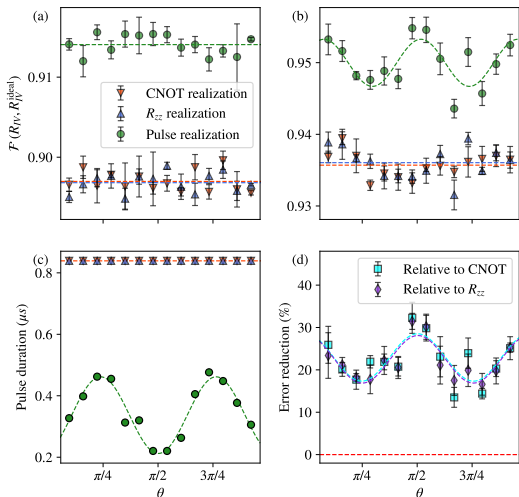
# Pulse realization of $\check{R}_{IV}(\theta)$

Suppose that  $\theta = \pi/3$ . We have the pulses of (a) CNOT realization and (b) direct pulse construction.



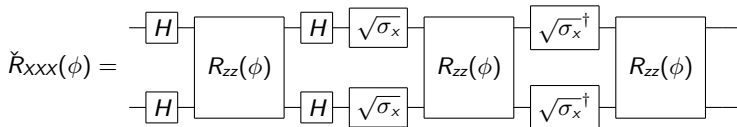
# Gate fidelity of $\check{R}_{IV}(\theta)$

Gate fidelity from (a) simulator or (b) ibmq\_montreal machine. (c) is the total pulse duration. (d) is the error reductions of pulse realization relative to other realizations.

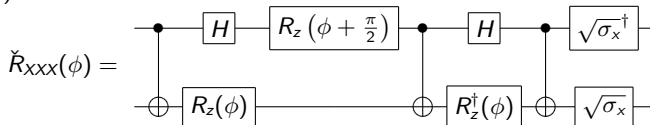


# Gate realization of $\check{R}_{XXX}(\phi)$

- The YBG  $\check{R}_{XXX}(\phi)$  requires minimal three  $R_{zz}$  (trapped-ion quantum computers)

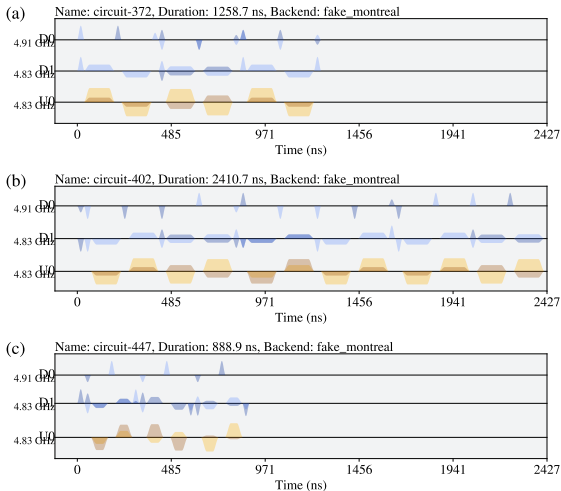


- The YBG  $\check{R}_{XXX}(\phi)$  requires minimal three CNOTs (superconducting quantum computers)



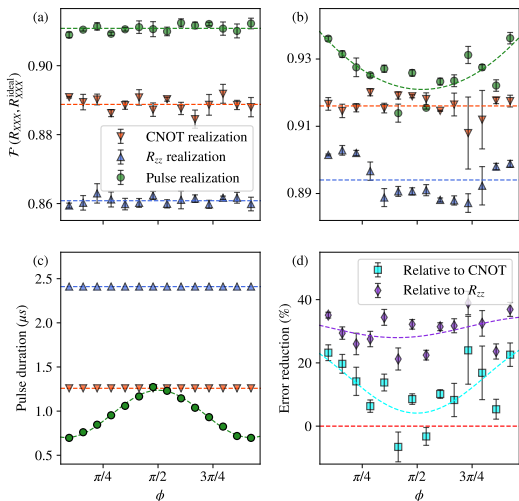
# Pulse realization of $\check{R}_{XXX}(\phi)$

Suppose that  $\phi = \pi/3$ . We have the pulses of (a) CNOT realization, (b)  $R_{ZZ}$  realization, and (c) direct pulse construction.



# Gate fidelity of $\check{R}_{XXX}(\phi)$

Gate fidelity from (a) simulator or (b) `ibmq_montreal` machine. (c) is the total pulse duration. (d) is the error reductions of pulse realization relative to other realizations.



# Conclusion

- The YBGs are the **building blocks** of the integrable circuits. It is a new method to study the far from equilibrium many-body quantum systems on quantum computers.
- The **optimal constructions** of YBGs can be obtained from their two-qubit geometric representations.
- YBGs realized from the **pulse realization** has the highest gate fidelities.
- In experiments, if two-qubit gates on the edge  $OA_1$  and faces  $OA_2A_3$  and  $A_1A_2A_3$  can be constructed with high fidelity, then high-fidelity YBGs can be constructed (with some additional single-qubit gates).



# Outlook

The quantum circuit model may be viewed as a generalized **factorisable scattering model**. It can also be viewed as a generalization of **topological quantum computation** [Y. Zhang, [arXiv:1111.3940](https://arxiv.org/abs/1111.3940)].

- TCIS: Trotter circuit of integrable system.
- IC: integrable circuit, which is composed of Yang-Baxter gates.
- LISC: logarithmic information spreading circuit, in which the information spreads logarithmic after a quench.
- QSC: quantum simulation circuit, which simulates the dynamics of physical system.
- Completely integrable many-qubit gates.

Unitary solutions of the tetrahedron equations as gates

<https://arxiv.org/pdf/2407.10731>

will help us to compress circuits even more. The tetrahedron equations are higher dimensional version of Yang-Baxter equations.

# Outlook

