

# Data Analytics with Graph Algorithms

## –

# A Hands-on Tutorial with Neo4J

March 4<sup>th</sup> 2019

**Lecturer: Dr. Lena Wiese**

# Agenda

- Short CV of speaker
- Graph Theory Basics
- The Neo4J Database
- Graph Algorithms
  - Centralities
  - Path Finding
  - Community Detection

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  - Community Detection

# Short CV Dr. Lena Wiese

- **University of Göttingen**  
(Group Leader Knowledge Engineering)
- **University of Salzburg**  
(Guest Lecturer)
- **University of Hildesheim**  
(Visiting Professor for Databases)
- **National Institute of Informatics, Tokyo**  
(funded by DAAD)
- **Robert Bosch India Ltd.**, Bangalore, India
- **TU Dortmund** (Master/PhD)
- **Universidad Complutense**, Madrid
  
- Web: <http://wiese.free.fr/>



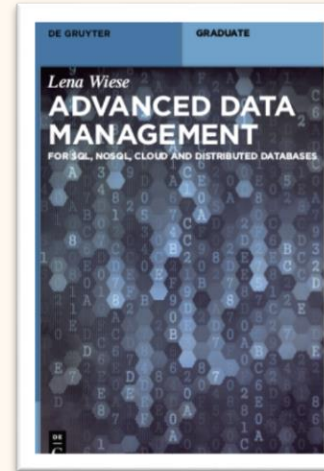
# Short CV Dr. Lena Wiese



# Short CV Dr. Lena Wiese

## Teaching and Research:

- **NoSQL Databases**  
(in particular Graph Databases)
- **Intelligent Data Management**  
(in particular, analytics of biomedical data for example, patient similarity analysis, disease prediction, etc.)
- **Intelligent Information Systems**  
(ontologies, recommender systems, etc.)



# Agenda

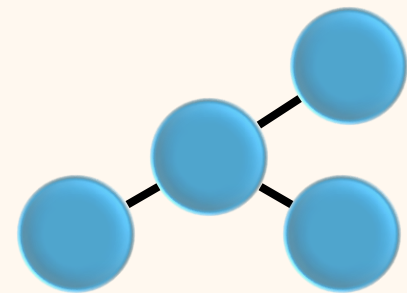
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# Graphs, graphs, graphs

There are lots of graphs in the real world:

- The Internet: a graph of web pages
- Social network: a graph of people
- Geographic Information System: a graph of locations
- Gene-Regulatory Network: a graph of genomic components

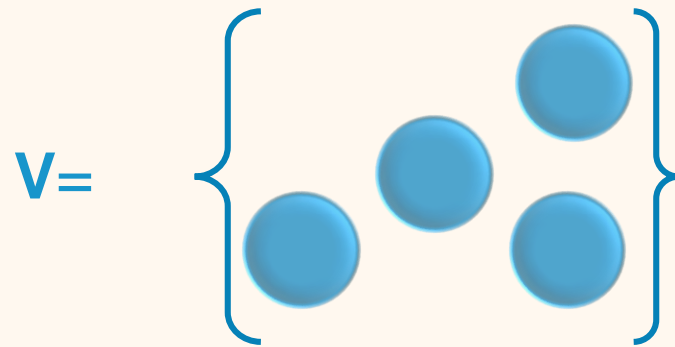




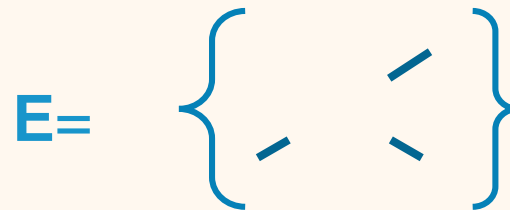
# Graphs in Mathematics

Mathematical definition of a graph:  $G=(V,E)$

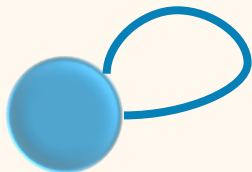
$V$  is a set of nodes  
(also called „vertices“)



$E$  is a set of edges  
(„links“/„relationships“)



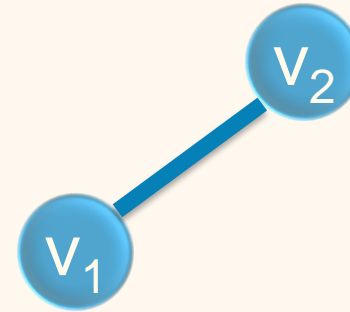
Special case:  
Self-loop



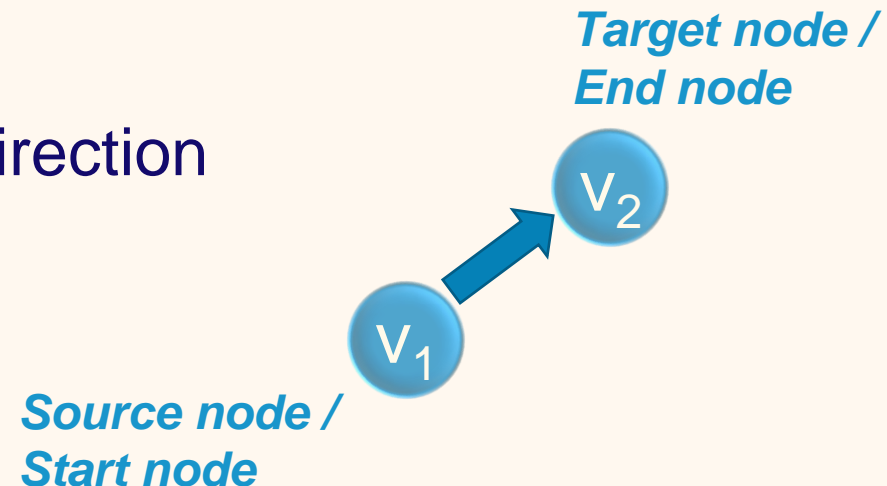
# Direction of edges

An edge can be

- **undirected:  $e=\{v_1, v_2\}$** 
  - can be traversed in both directions



- **directed:  $e=(v_1, v_2)$** 
  - can be traversed in one direction



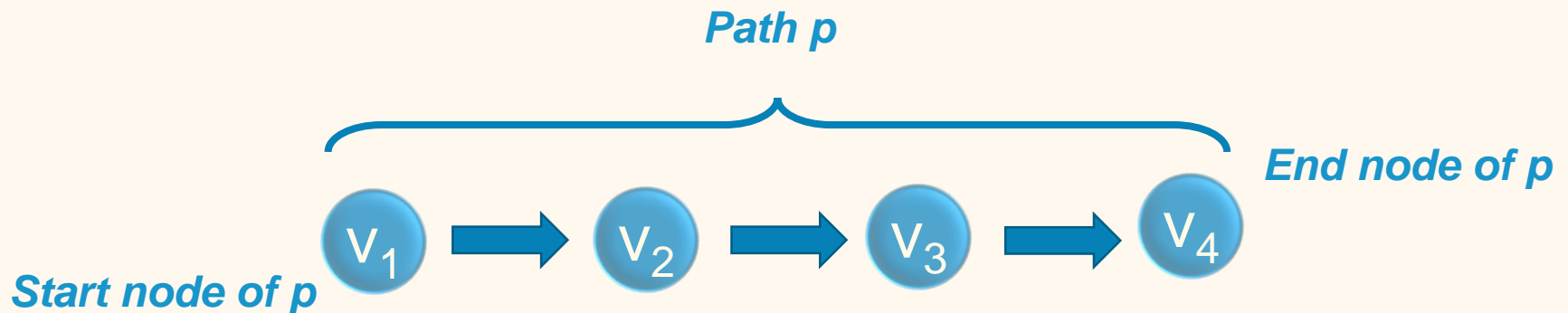
# Traversal and path

## Traversal:

Go from one node to another by following an edge

## Path:

A concatenation of nodes and edges that can be traversed in the graph

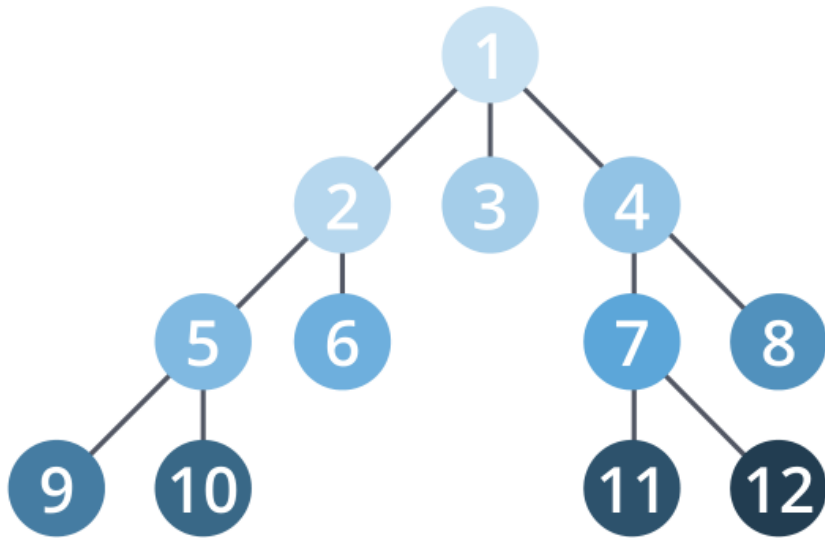


# Exhaustive Graph Traversals

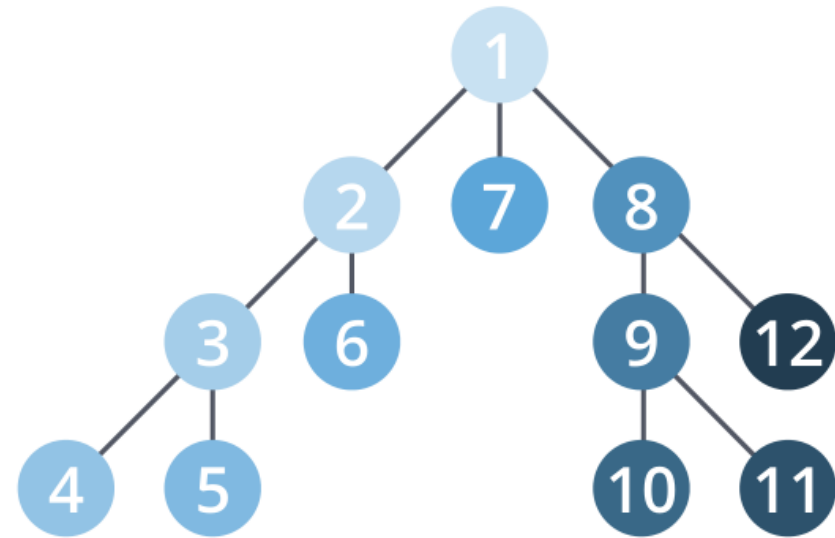
**BFS:** From the start node visit all direct neighbors first before visiting a neighbor's neighbor

**DFS:** From the start node choose one neighbor then visit the neighbor's neighbor

Picture source: [1, page 18]



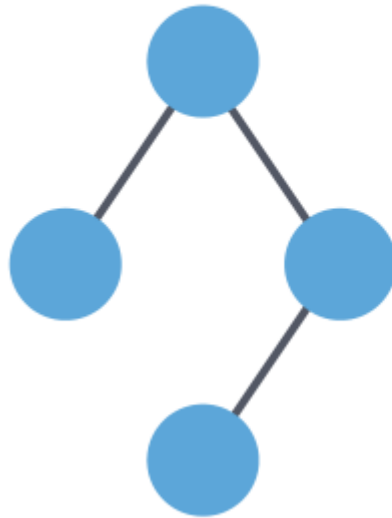
*Breadth-first search*



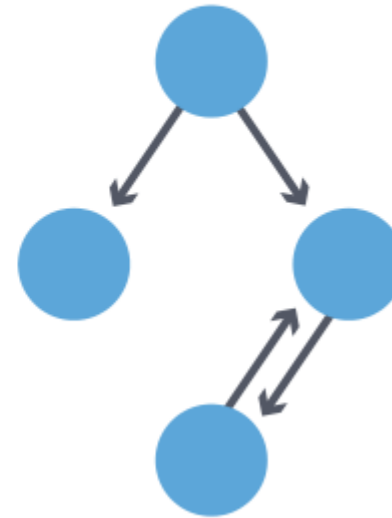
*Depth-first search*

# Direction in graphs

Undirected



Directed

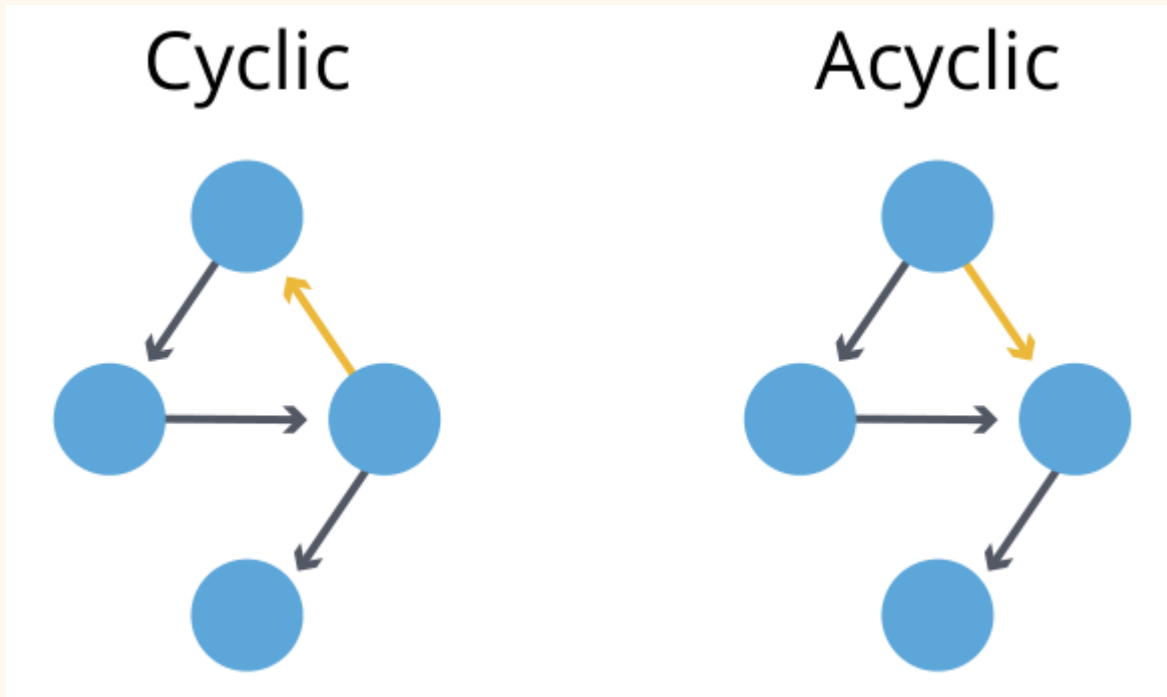


Picture source:  
[1, page 19]

Example:  
Mutual friendship relation

Example:  
One-way streets

# Cycles in graphs



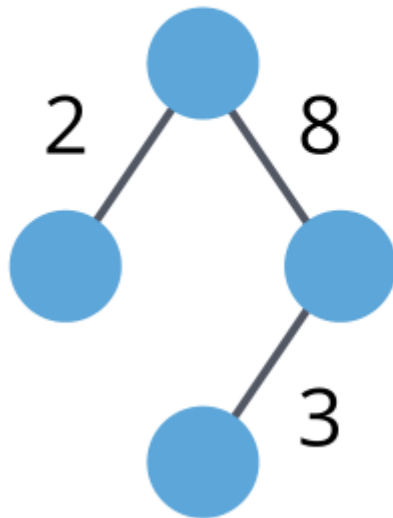
Picture source:  
[1, page 19]

**Cycle:** A path containing the same node as start and end node

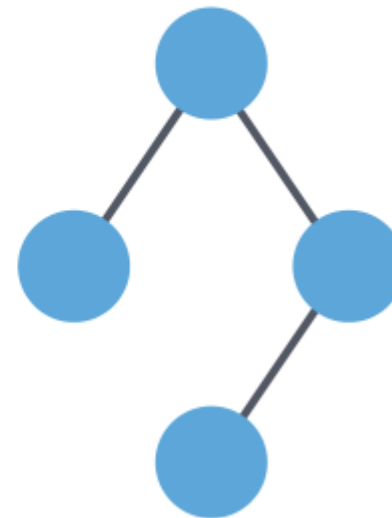
**Triangle:** A cycle with three nodes

# Weights in graphs

Weighted



Unweighted



Picture source:  
[1, page 19]

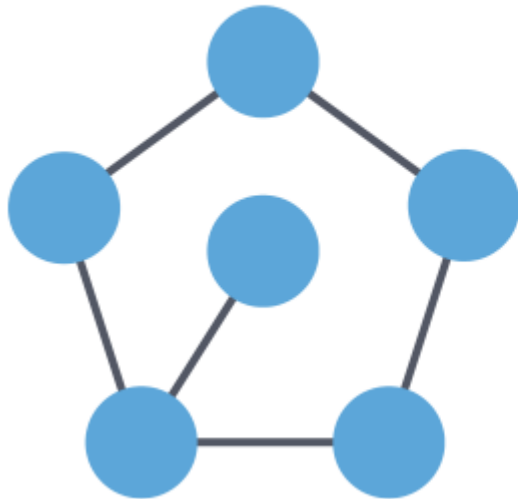
Example „shortest path“:  
Find a path between two nodes with minimum cost

Example „shortest path“:  
Find a path with minimum number of edges



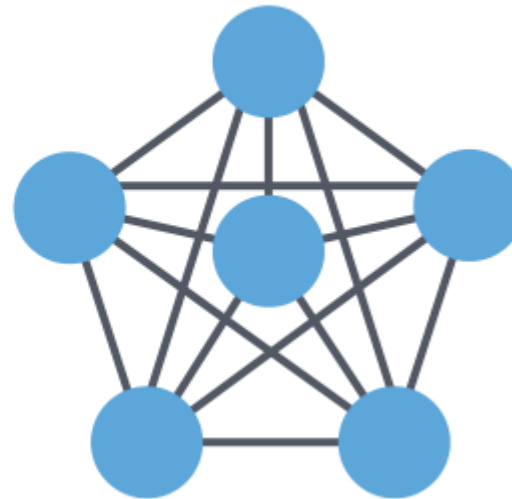
# Edge count in graphs

Sparse



Only few edges exist

Dense



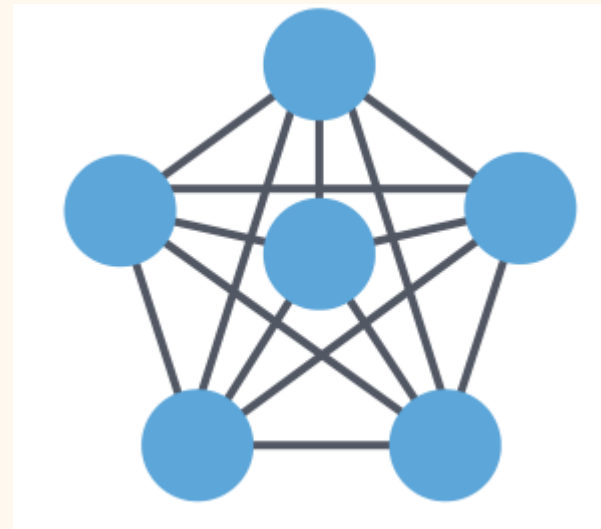
Many edges exist

**Complete graph:**  
All edges exist

Picture source:  
[1, page 19]

# Quiz

What is the **edge count** in a complete undirected graph without self-loops?



Picture source:  
[1, page 19]

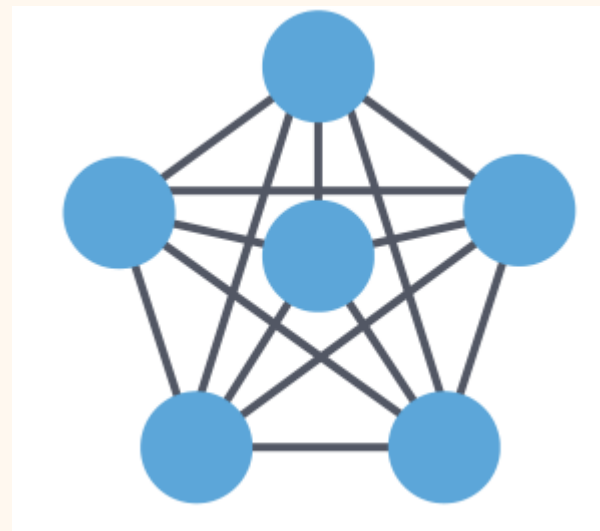
# Answer

What is the edge count in a complete undirected graph without self-loops?

$$\binom{|V|}{2} = \frac{|V|!}{2! (|V| - 2)!} = \frac{|V|(|V| - 1)}{2}$$

Example:  $|V|=6$

$$\binom{6}{2} = \frac{6!}{2!4!} = \frac{6 \cdot 5}{2} = 15$$



Picture source:  
[1, page 19]

# Agenda

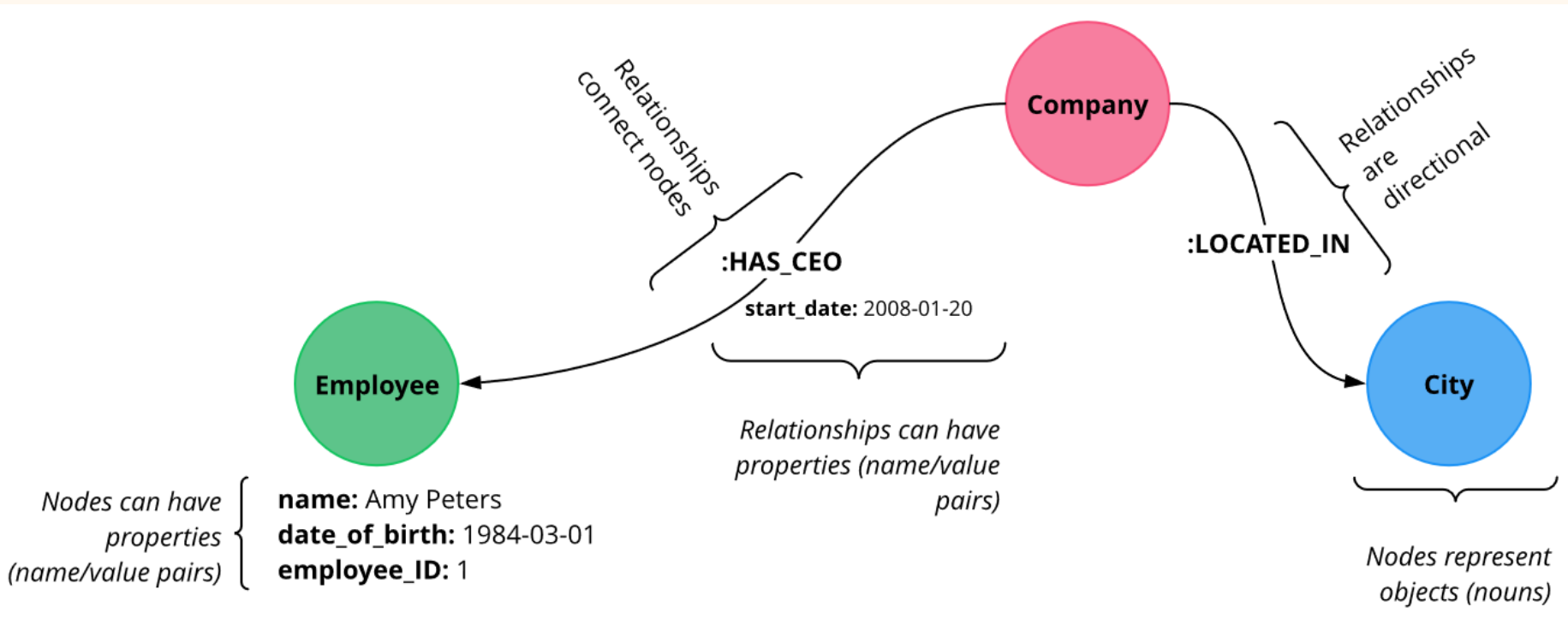
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# Property Graph Model

Nodes have **labels** (e.g. Employee)

Edges have **types** (e.g. :HAS\_CEO)

Information is stored in ***name:value*** pairs („properties“)



Picture source: <https://neo4j.com/developer/graph-database/>

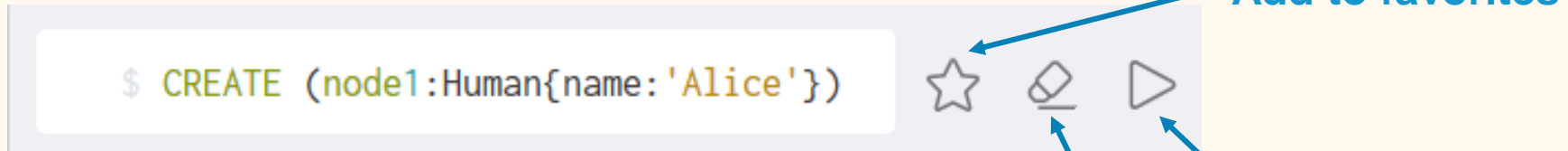
# Neo4J Graph Database

- Open source graph database written in Java
- Query language called Cypher
- Interface called Neo4J Browser

The screenshot shows the Neo4J Browser interface. The browser address bar displays 'localhost:7474/browser/'. The left sidebar contains 'Database Information', 'Node Labels' (with buttons for 'Enhancer', 'Gene', 'TFPair'), 'Relationship Types' (with buttons for 'EPI', 'binds'), and 'Property Keys' (with input fields for 'enhancerID', 'genename', 'knownBioGridPair', 'knownCompelPair', 'name', 'promotorID', 'pwm1', 'pwm2', 'zscore'). The main area shows a Cypher query: '\$ MATCH p=()->() RETURN p LIMIT 25'. Below the query, there are buttons for 'Graph', 'Table', 'Text', and 'Code'. The 'Graph' view is selected, showing a visualization of the query results with nodes and relationships. The nodes are labeled with IDs like 'VSCMY...', 'VSCET...', 'VSCRE...', 'VSGAT...', 'VSHNF...', 'VSGRE...', 'VSPAP1...', and 'VSGE1'. The relationships are labeled 'binds'.

- Visualization of query results
- Support for graph algorithms

# Create Nodes with Cypher



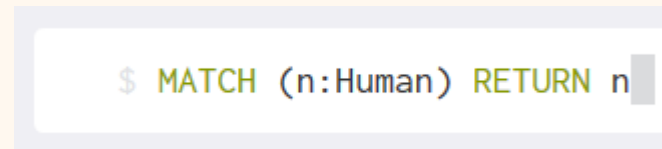
Create a new node called *node1*  
with label *Human*  
and a *name*-property set to *Alice*

## Task 1:

Create a new node called *node2*  
with label *Human*  
and a *name*-property set to *Bob*

## Task 2:

Find all nodes matching the label *Human*  
and return them





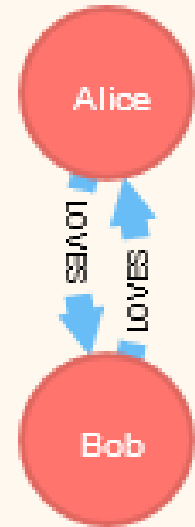
# Create Edges with Cypher

```
1 MATCH (node1:Human{name:'Alice'}) MATCH (node2:Human{name:'Bob'})
2 CREATE (node1)-[:LOVES{until:'forever'}]->(node2)
```

Find two nodes **matching** the label *Human* and **create** a new edge with type *:LOVES* and an *until*-property set to *forever*

## Task 3:

Create a new edge pointing from *Bob* to *Alice*



## Task 4:

Find all paths *p* with edges of type *:LOVES* and return them

```
$ MATCH p=()-[r:LOVES]->() RETURN p
```

# MERGE

## Avoid duplicates:

If node / edge does **not** exist: CREATE it as new

If node / edge exists: MATCH and return the node / edge

```
$ MERGE (node1:Human{name:'Alice'})
```

# Deletion with Cypher

Delete edge between two nodes with certain **properties**

```
1 MATCH (n1:Human)-[e:LOVES]->(n2:Human) WHERE n1.name='Alice' AND n2.name='Bob'  
2 DELETE e
```

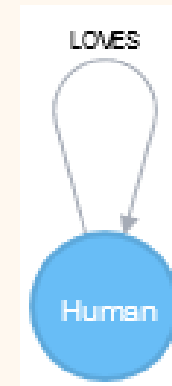
Or delete edge by **ID**

```
$ MATCH ()-[e:LOVES]->() WHERE id(e)=14 DELETE e
```

# Graph schema with Cypher

Show schema information (all **types** and **labels**)

```
$ CALL db.schema
```



# Agenda

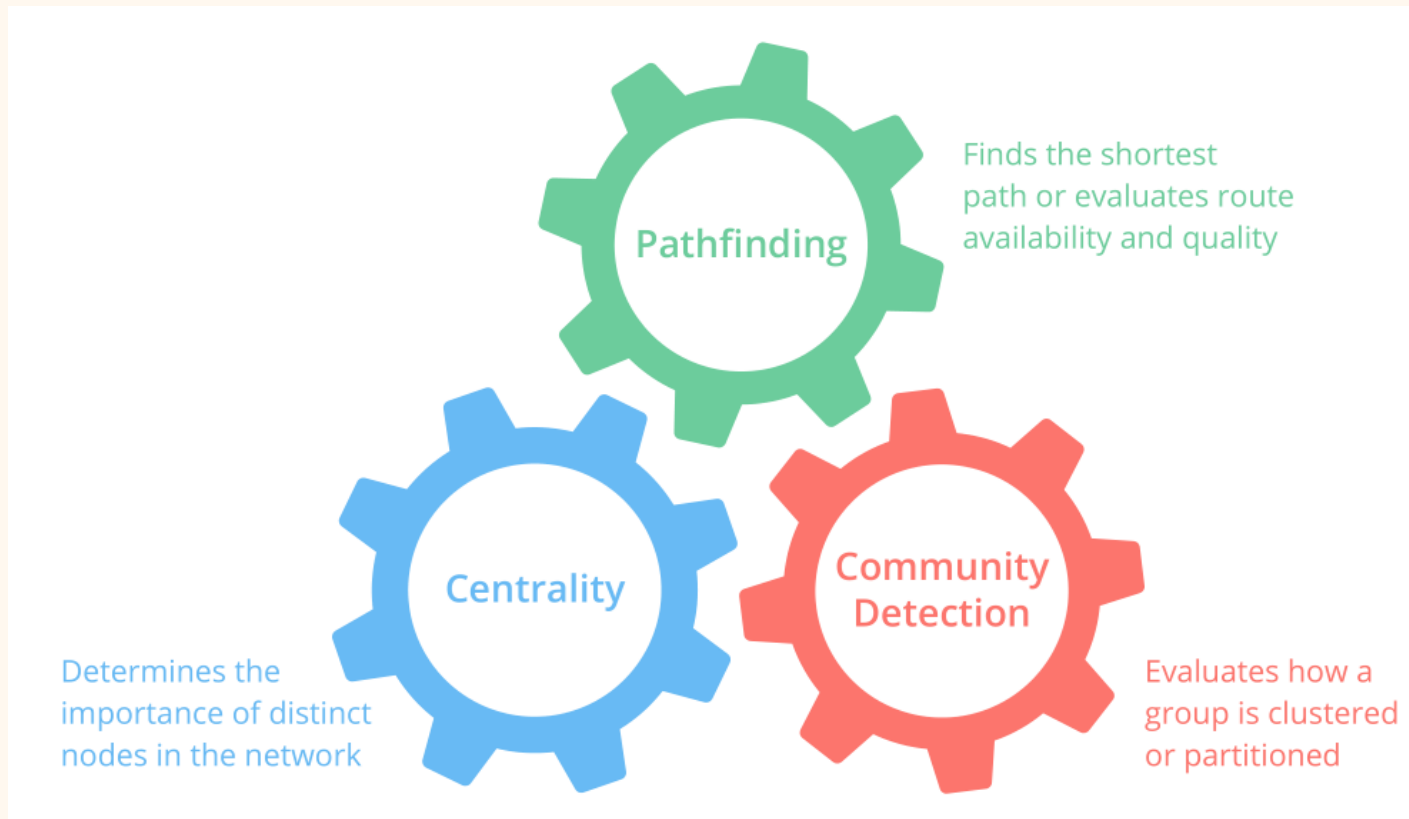
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# Graph Algorithms

**Centrality:** Find one or more nodes with an optimal score

**Pathfinding:** Find one or more optimal paths in a graph

**Community Detection:**  
Find densely connected subgraphs



Picture source: [1, page 15]

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# Centrality: PageRank

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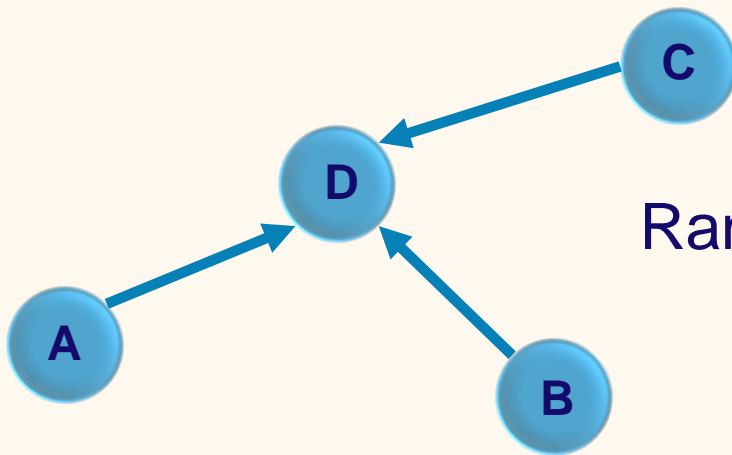
Google's PageRank score represents the importance of web pages in the Internet

- Rank of a page depends on the **in-links** to a page
- PageRank is transitive
  - PageRank of a node is influenced by the **neighbors'** PageRank
- In each iteration a node distributes its rank to its neighbors along its **out-links**

Reference: Page, L., Brin, S., Motwani, R., & Winograd, T. (1999). The PageRank citation ranking: Bringing order to the web. Stanford InfoLab

# Centrality: PageRank

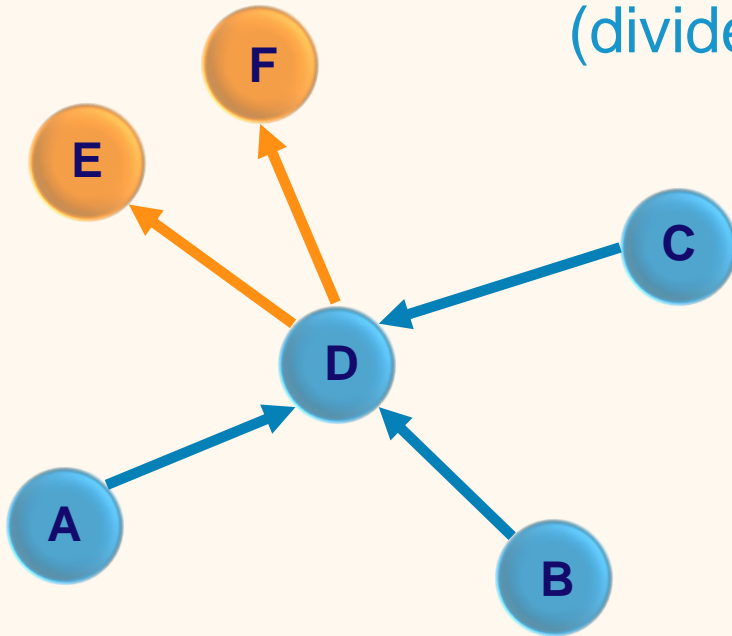
Step 1: Sum up the ranks of your in-links



$$\text{Rank}(\mathbf{D}) = \text{Rank}(\mathbf{A}) + \text{Rank}(\mathbf{B}) + \text{Rank}(\mathbf{C})$$

# Centrality: PageRank

**Step 2:** Distribute your rank among your out-links  
(divide by amount of out-links)



$$\text{Rank}(\mathbf{E}) = \frac{1}{2} \text{Rank}(\mathbf{D})$$

$$\text{Rank}(\mathbf{F}) = \frac{1}{2} \text{Rank}(\mathbf{D})$$

# Centrality: PageRank

## Random surfer model:

- A web surfer randomly follows **out-links** of pages
- **Uniform** probability distribution: if a page has **m** out-links, each link is followed with probability  $1/m$
- The higher the PageRank, the higher the likelihood that a random surfer will be at this page at an arbitrary point of time

**Note:** A **few** in-links from very **important** pages raise your PageRank more than **many** in-links from **unimportant** pages.

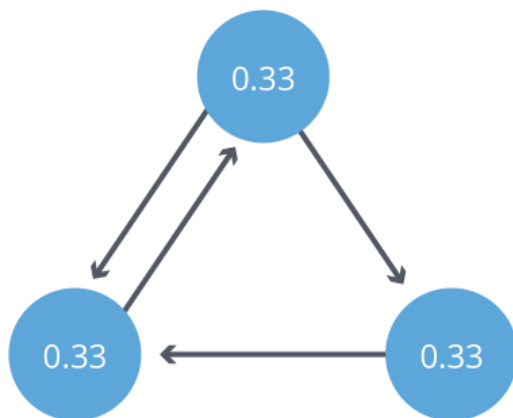
# Centrality: PageRank

**Example:** start with equal PageRank for all pages, then iterate for a certain number of rounds

Pass 0

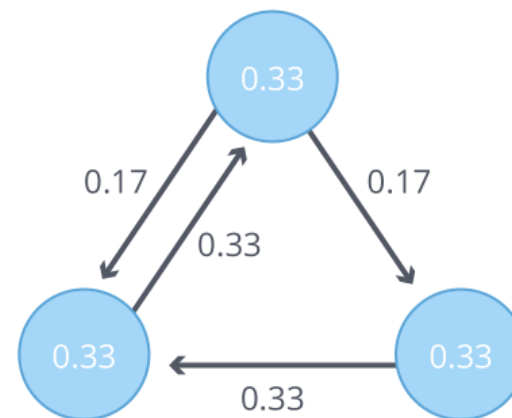
Step 1

Node Value =  $1/n$  ( $n$  = Total # of Nodes)



Step 2

Link Value = Node Value / # of Its Out-Links



Picture source:  
[1, page 35]

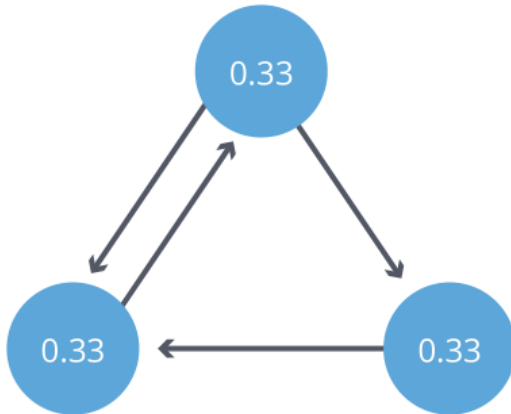
# Centrality: PageRank

Picture source:  
[1, page 35]

Pass 0

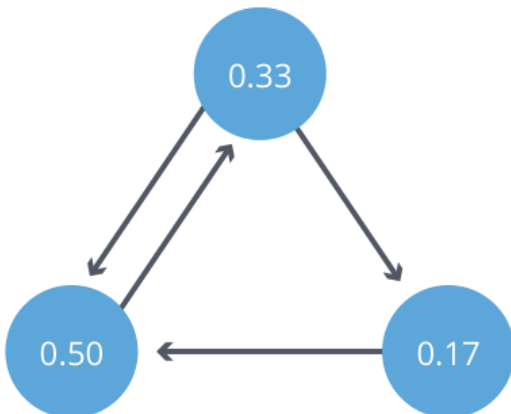
## Step 1

Node Value =  $1/n$  ( $n$  = Total # of Nodes)



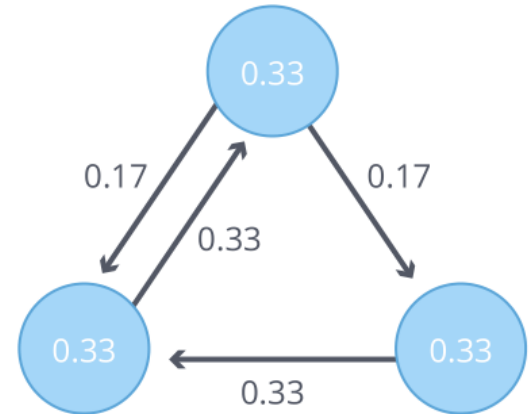
Node Value = Sum of Prior In-Link Values

Pass 1

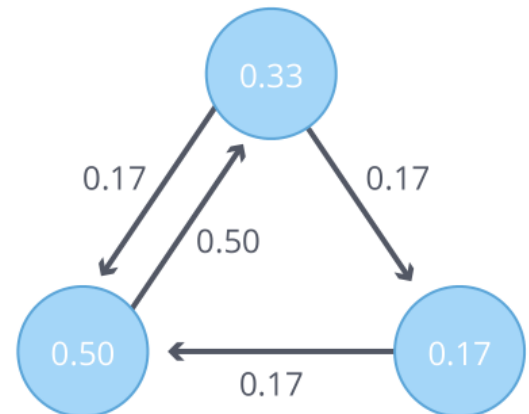


## Step 2

Link Value = Node Value / # of Its Out-Links



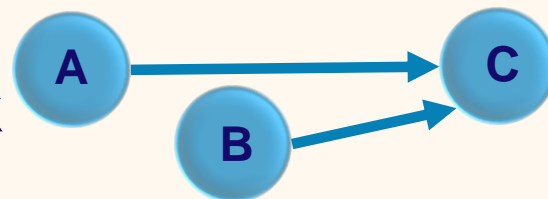
Link Value = Node Value / # of Its Out-Links



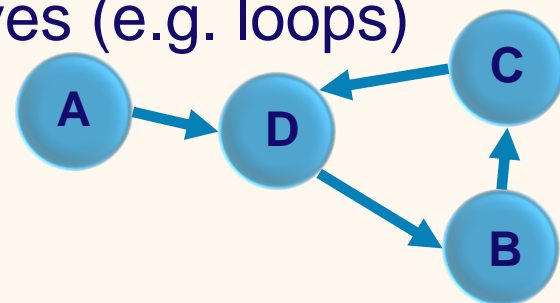
# Centrality: PageRank

## Problem Cases:

- **Dead end:** a node that has no out-link



- **Rank sink:** a group of nodes that have in-links from other nodes but out-links just among themselves (e.g. loops)



- **Disconnected subgraphs:** subgraphs without links between each other

These problem cases disturb the **distribution** of PageRank

## Solution:

- add a constant **rank source** to each page
- corresponds to random **jumps** (instead of following links)



# Centrality: PageRank

The final formula:

$$\text{Rank}(u) = d \cdot \sum_{v \in \text{In}(u)} \frac{\text{Rank}(v)}{|\text{Out}(v)|} + (1 - d)$$

**Damping factor**  
(reduces the likelihood of following a link by random surfer)

$v$  has an **in-link** to  $u$

Rank of  $v$  divided by **amount of out-links**  
(for random surfer)

**rank source** for each page (for random jumps)

Possibly: **Normalization** of ranks so that all ranks sum up to 1

# Graph Algorithms in Cypher

- Two options:

1. **Streaming** of result

```
CALL algo.pageRank.stream
```

```
CALL algo.<name>.stream
```

- Immediately outputs the result

2. **Writing** result into a property

```
CALL algo.pageRank
```

```
CALL algo.<name>
```

- Properties can be queried in a second query

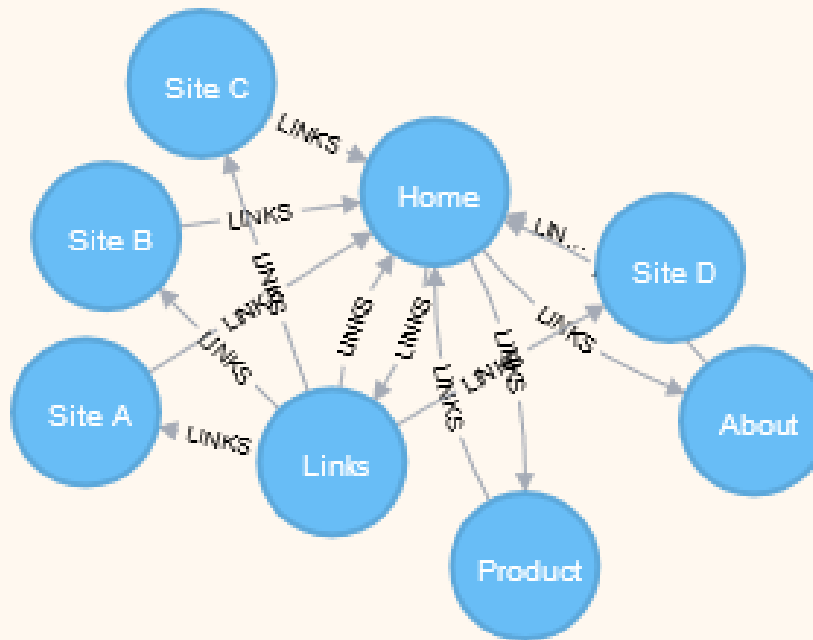
```
write: true,writeProperty:"pagerank":
```

- Clause **YIELD** defines which statistics to print out

```
YIELD nodeId, score
```

# PageRank in Cypher

**Task 5:** Create sample graph of 8 web pages (see <https://neo4j.com/docs/graph-algorithms/current/> ➔ PageRank)

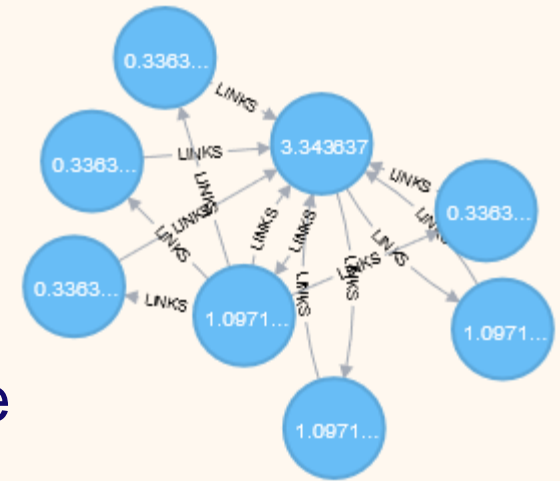


# PageRank in Cypher

## Task 6:

- Run PageRank with 20 iterations
- Increase the amount of iterations and observe the effect on the PageRank
- Which page is the most important one?
- **Optional:** Write the PageRank into a node property called *pagerank* and display the page rank for each node

```
1 CALL algo.pageRank.stream("Page", "LINKS",
2 {iterations:20})
3 YIELD nodeId, score
```



```
1 CALL algo.pageRank('Page', 'LINKS',
2 {iterations:20, dampingFactor:0.85, write: true,writeProperty:"pagerank"})
```

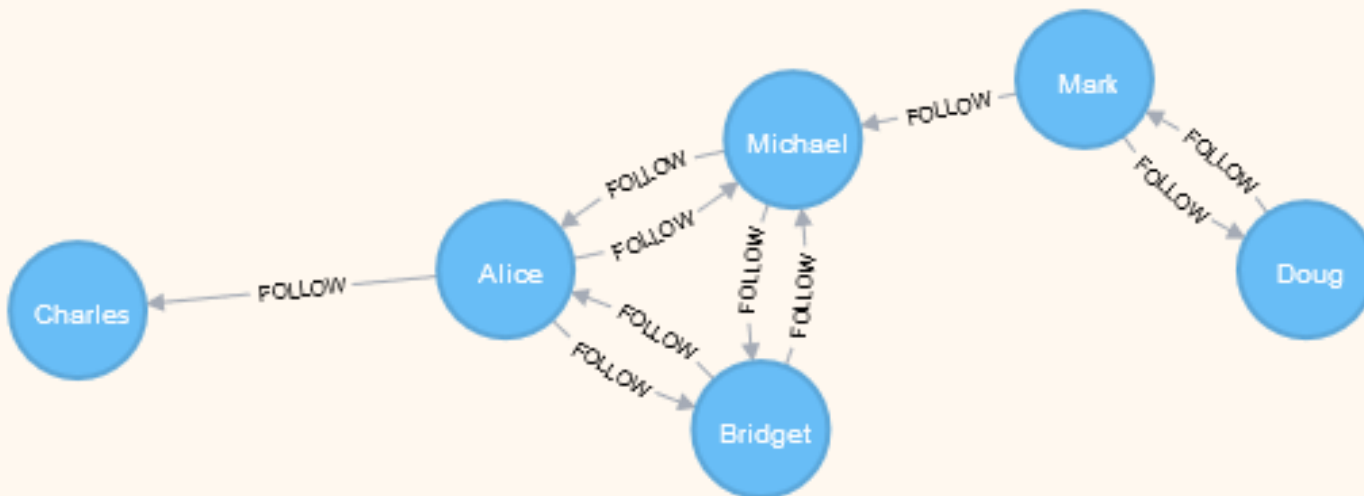
# Centrality: Degree

# Centrality: Degree

Degree centrality is the amount of in-links and out-links of each node

## Task 7:

- Create sample graph of 6 users with a **:FOLLOW** relationship (see <https://neo4j.com/docs/graph-algorithms/current/> → The Strongly Connected Components algorithm)



# Centrality: Degree

Degree centrality is the amount of in-links and out-links of each node

## Task 8:

- For each user return the amount of out-links and the amount of in-links

```
1 MATCH (u:User)
2 RETURN u.id AS name,
3 size((u)-[:FOLLOW]->()) AS follows,
4 size((u)<-[:FOLLOW]-()) AS followers
```

- Which user has the highest degree centrality?
- **Optional:** Return the total amount of links of each node

# Centrality: Betweenness



# Centrality: Betweenness

Betweenness centrality is the fraction of shortest paths going through a node

$$C_B(v) = \sum_{s \neq v \neq t \in V} \frac{\sigma_{st}(v)}{\sigma_{st}}$$

The amount of shortest paths between s and t going through v

The total amount of shortest paths between s and t

Sum over all node pairs s,t (different from v)

Betweenness centrality for node v

Nodes with high betweenness centrality ensure crucial connections in the graph: „Bridge“ between different subgraphs

# Centrality: Betweenness

$$C_B(v) = \sum_{s \neq v \neq t \in V} \frac{\sigma_{st}(v)}{\sigma_{st}}$$

Let  $v = \text{Alice}$

Shortest paths via Alice:

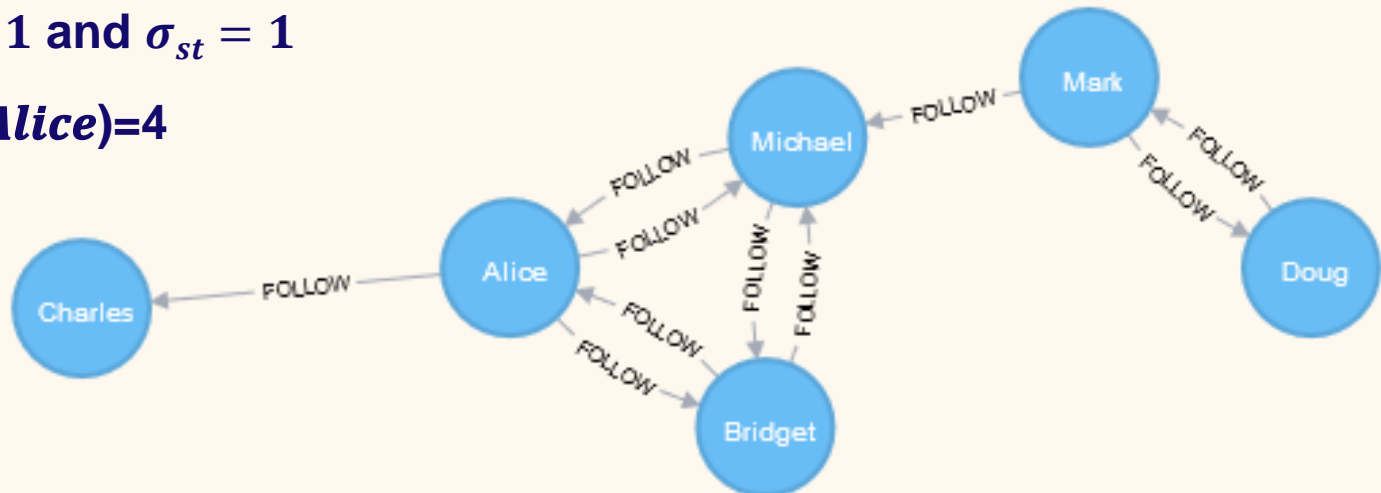
1. Michael-Charles
2. Mark-Charles
3. Doug-Charles
4. Bridget-Charles

For all these paths:

Only one shortest path (and only through Alice)

$\sigma_{st}(\text{Alice}) = 1$  and  $\sigma_{st} = 1$

Hence  $C_B(\text{Alice})=4$



# Centrality: Betweenness

## Task 9:

- For each user return the betweenness centrality along the **:FOLLOW** relationship

```
1 CALL algo.betweenness.stream('User','FOLLOW',{direction:'out'})
2 YIELD nodeId, centrality
3 MATCH (user:User) WHERE id(user) = nodeId
4 RETURN user.id AS user,centrality
5 ORDER BY centrality DESC;
```

- **Optional:** remove the edges between **Michael** and **Bridget** and observe the effect on the betweenness centrality

# Agenda

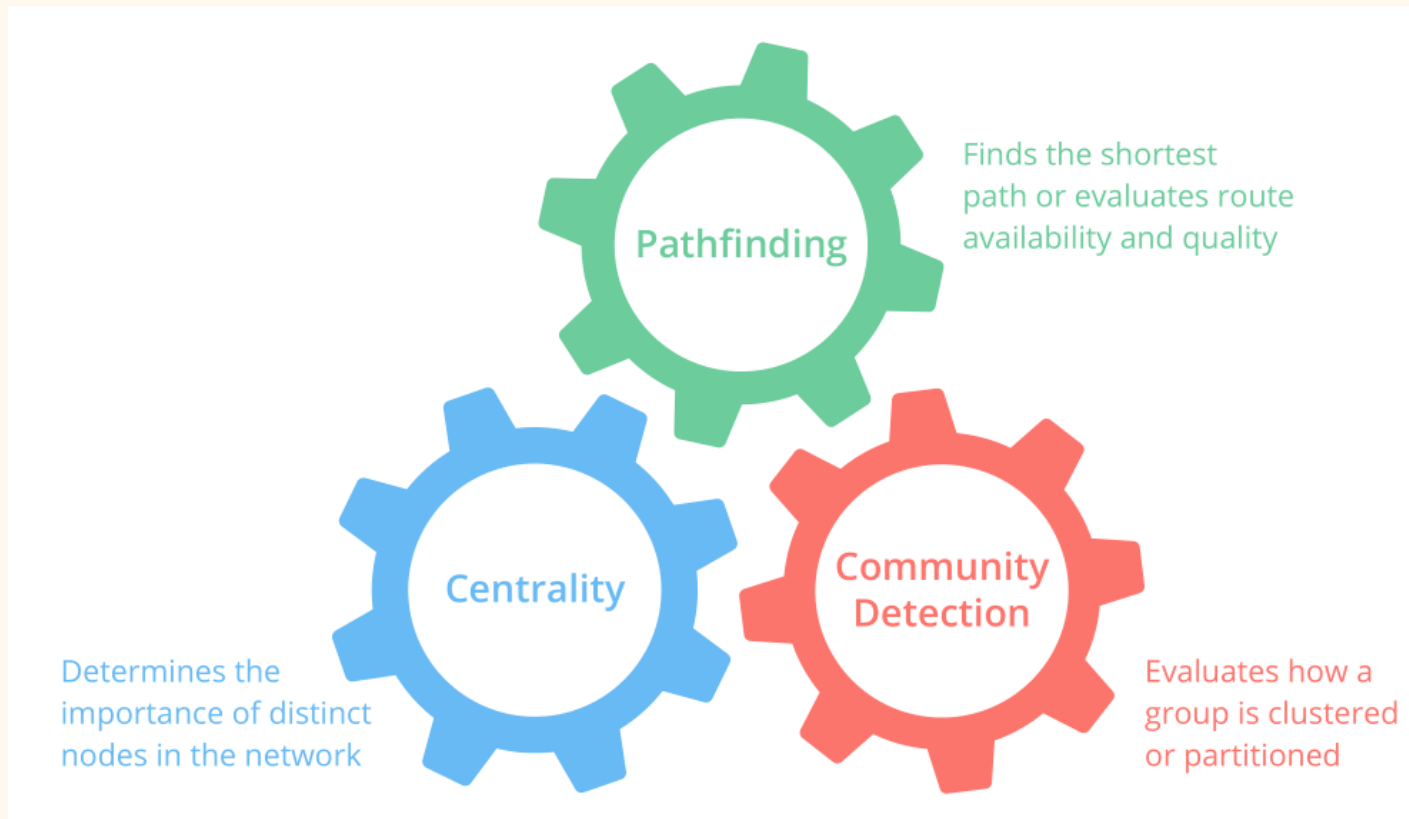
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# Graph Algorithms

**Centrality:** Find one or more nodes with an optimal score

**Pathfinding:** Find one or more optimal paths in a graph

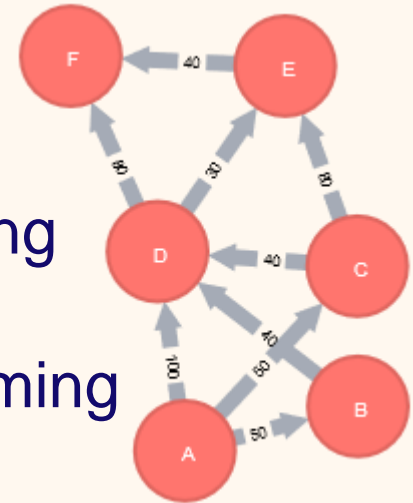
**Community Detection:**  
Find densely connected subgraphs



Picture source: [1, page 15]

# Pathfinding: Minimum Weight Spanning Tree

# Pathfinding: Minimum Weight Spanning Tree



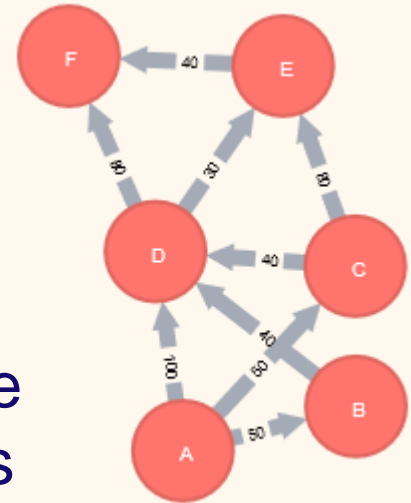
- **Minimum Spanning Tree:**
  - If a graph has  $|V|$  nodes, then a spanning tree has  $|V|-1$  edges
  - The tree has a **root** node (with no incoming edges)
  - From the root node we can reach **all other** nodes in the graph with minimal **total** cost by using the edges of the spanning tree

## Use case:

Distribute information from the root node to all other nodes for example in communication networks or social networks

# Pathfinding: Minimum Weight Spanning Tree

- **Prim's algorithm**
  - Maintain a list of unvisited nodes
  - Start with the **root** node
  - From the unvisited nodes select the one that can be connected to the tree nodes with **minimal** cost
  - **Repeat** until all nodes are visited

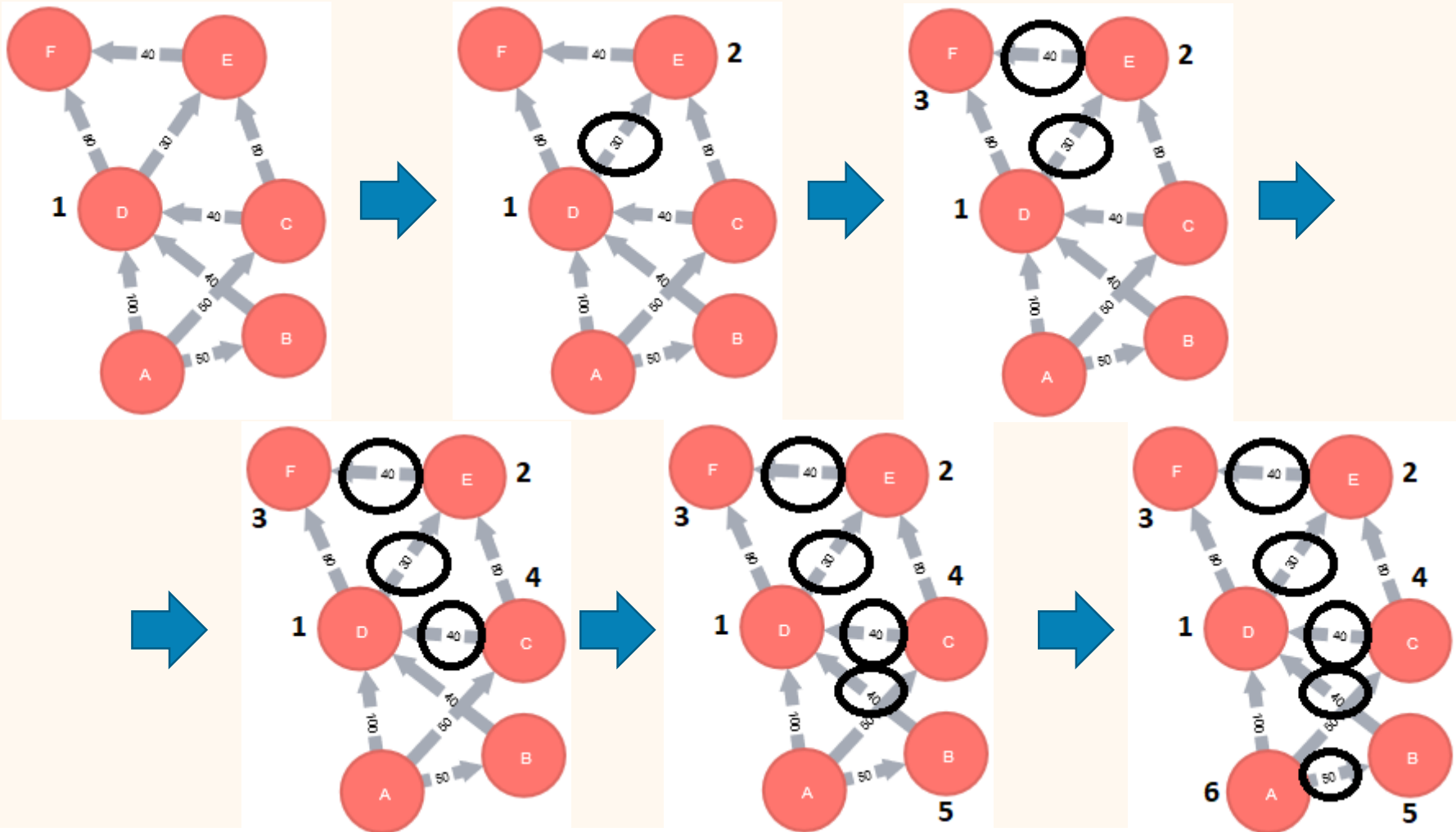


**For undirected graphs:**

**We ignore direction of edges (traversal in both directions)**



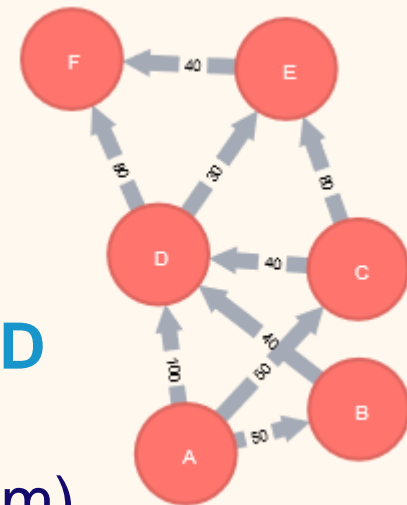
# Pathfinding: Minimum Weight Spanning Tree



# Pathfinding: Minimum Weight Spanning Tree

## Task 10:

- Create sample graph of 6 locations with a **:ROAD** relationship (see <https://neo4j.com/docs/graph-algorithms/current/> → The Shortest Path algorithm)
- **Optional:** Display the cost of each road on the edge



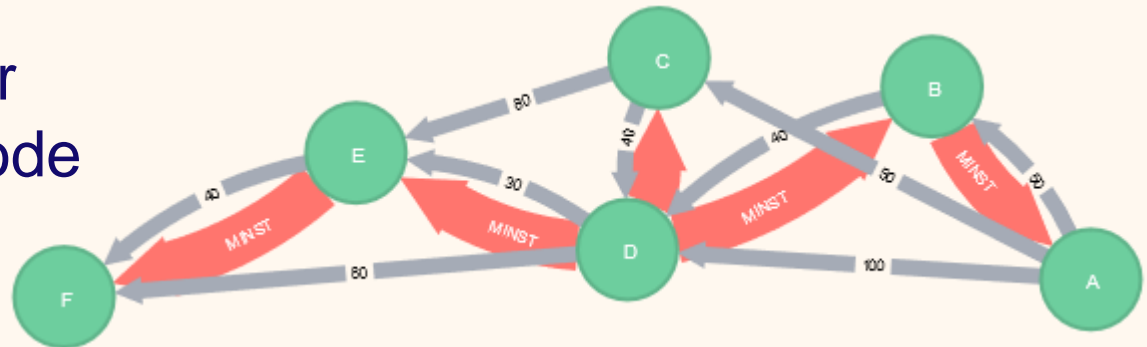
# Pathfinding: Minimum Weight Spanning Tree

## Task 11:

- Choose location **D** as the root node and compute a minimum spanning tree by creating new edges of type **:MINST**

```
1 MATCH (n:Loc {name:"D"})
2 CALL algo.spanningTree.minimum('Loc', 'ROAD', 'cost', id(n),
3   {write:true, writeProperty:"MINST"})
4 YIELD loadMillis, computeMillis, writeMillis, effectiveNodeCount
5 RETURN loadMillis, computeMillis, writeMillis, effectiveNodeCount
```

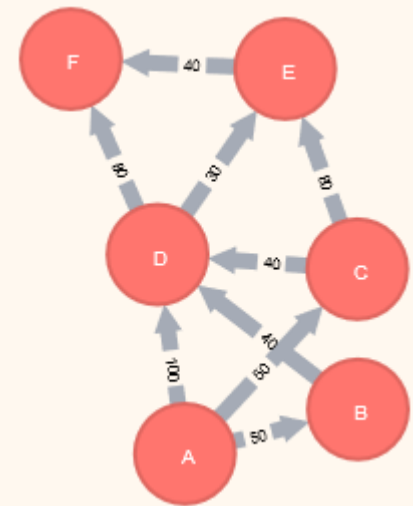
- Optional:** use another location as the root node



# Pathfinding: Shortest Path

# Pathfinding: Shortest Path

- **Shortest Path with minimum cost:**
  - Provide a start and an end node
  - Find the minimum-cost path between these two nodes
  - In unweighted graph: edge cost 1



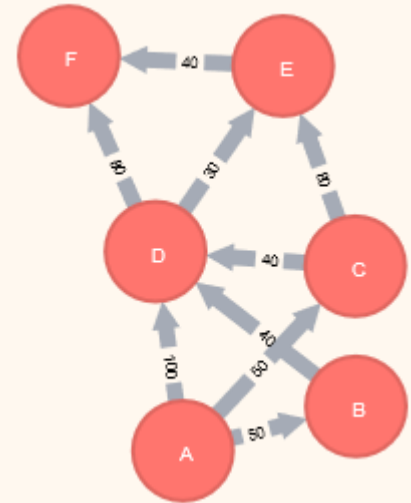
- **Example:**
  - Start node **A** and end node **F**
  - One minimum-cost shortest path:

$$\mathbf{A - C - D - E - F}$$
$$50 + 40 + 30 + 40 = 160$$

**Use case:**  
**Travel route planning**

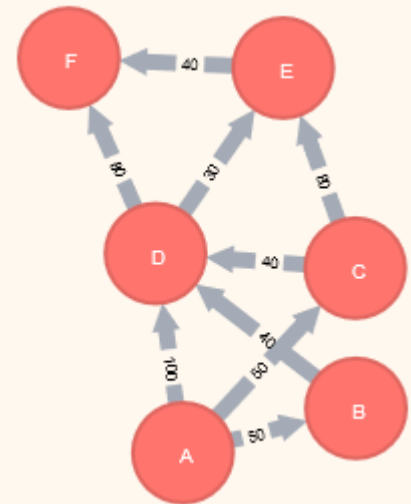
# Pathfinding: Shortest Path

- **Dijkstra's algorithm:**
  - Maintain a list of unvisited nodes
  - For each node maintain its distance to the start node (initially:  $\infty$ )
  - Start with the **start** node and set its distance to 0
  - For **all** unvisited **neighbor** nodes: set their distance to be the edge cost to the start node
  - Remove the start node from the list of unvisited nodes

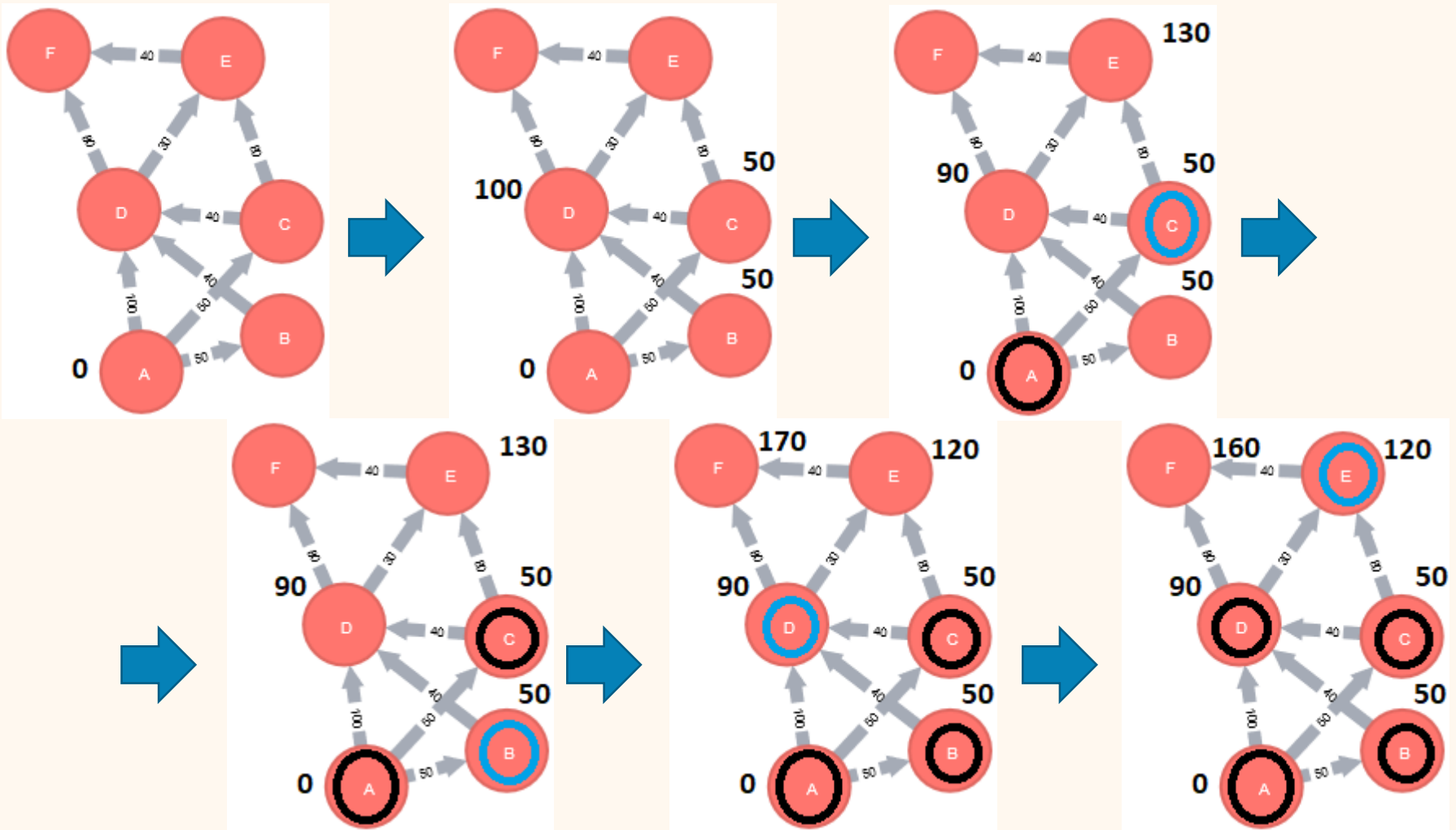


# Pathfinding: Shortest Path

- **Dijkstra's algorithm (continued):**
  - Select one unvisited node with currently **smallest** distance
    - Make it the **current node**
  - For all **unvisited** neighbors of the current node:
    - Sum up the edge cost to the neighbor and the distance of the current node
    - If smaller than current distance of the neighbor: update distance (shorter path found)
  - **Remove** current node from list of unvisited nodes
  - Repeat until **end node** is visited



# Pathfinding: Shortest Path





# Pathfinding: Shortest Path

## Task 12:

- Choose location **A** as the start node and location **F** as the end node and compute the shortest path

```
⚠ 1 MATCH (start:Loc{name:'A'}), (end:Loc{name:'F'})
   2 CALL algo.shortestPath.stream(start, end, 'cost')
   3 YIELD nodeId, cost
   4 RETURN algo.getNodeById(nodeId).name AS name, cost
```

- **Optional:** use other locations as start and end nodes

# Agenda

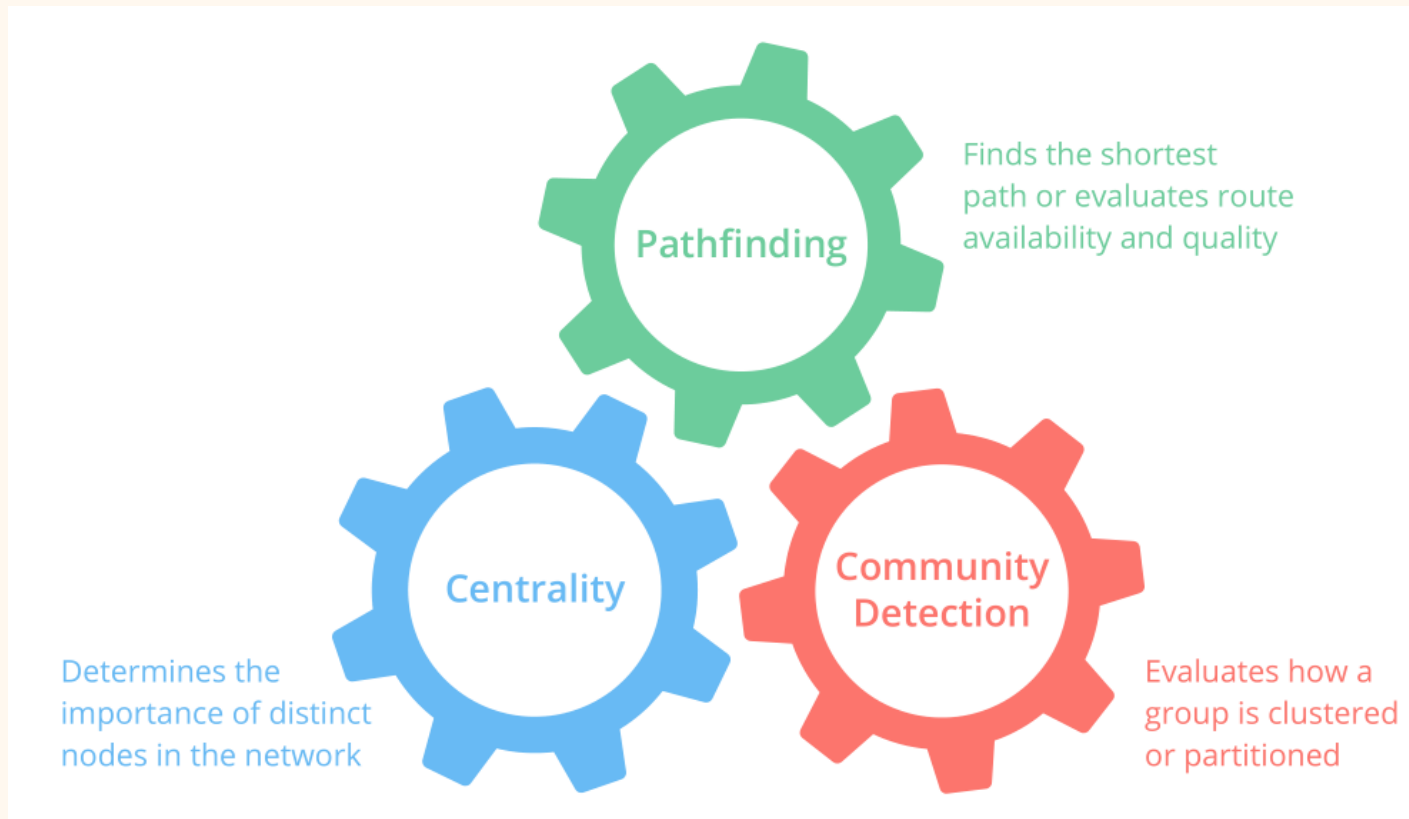
- Short CV of speaker
- Graph Theory Basics
- The Neo4J Database
- Graph Algorithms
  - Centralities
  - Path Finding
  - **Community Detection**

# Graph Algorithms

**Centrality:** Find one or more nodes with an optimal score

**Pathfinding:** Find one or more optimal paths in a graph

**Community Detection:**  
Find densely connected subgraphs

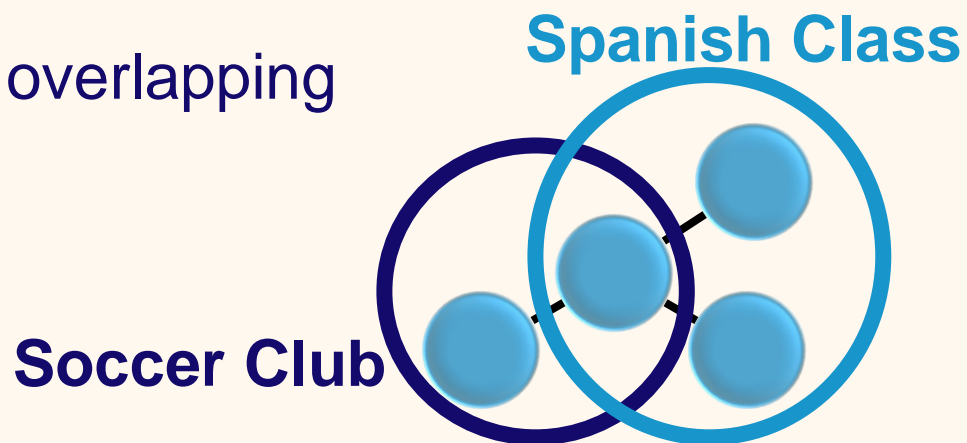


Picture source: [1, page 15]

# Community Detection

**Community: Densely connected subgraph**

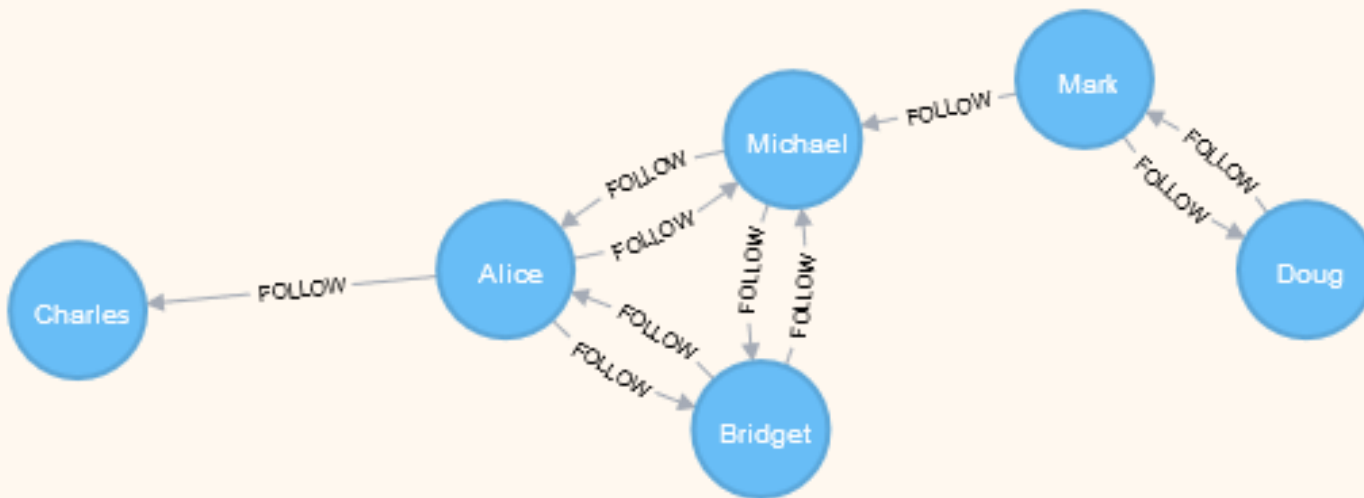
- More communication between the nodes of a community than to other nodes of the graph
- Communities may be overlapping



# Community Detection: Strongly Connected Components

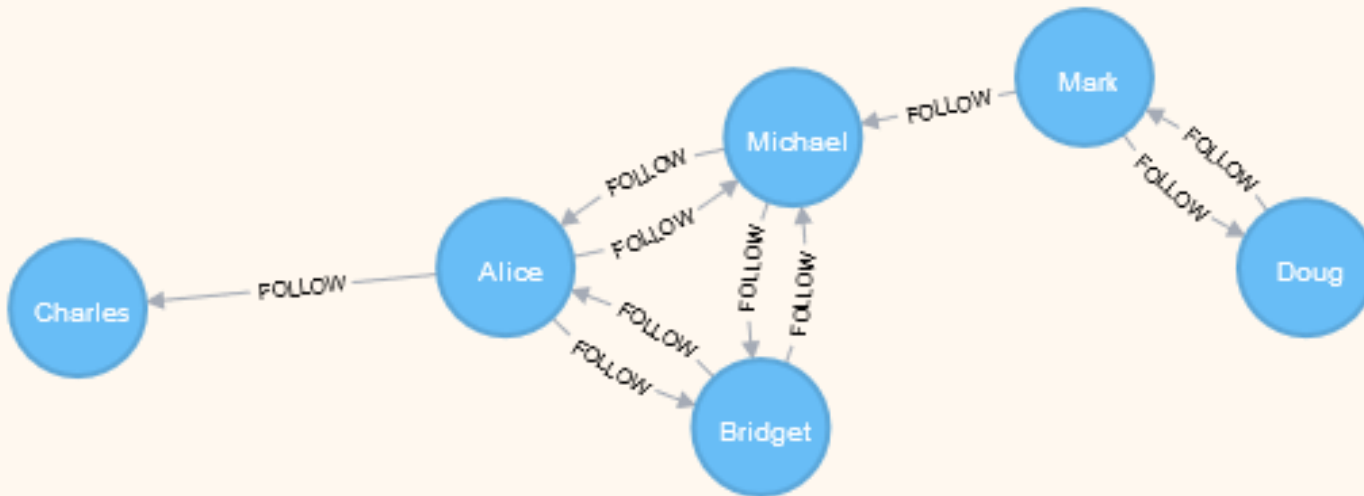
# Community Detection: Strongly Connected Components

- **Strongly Connected Components:**
  - Apply to directed graphs
  - Two nodes **A**, **B** are in the same strongly connected component if there is a path from **A** to **B** and a path from **B** to **A**



# Community Detection: Strongly Connected Components

- **Example:**
  - **Charles** has no out-links and is a component on his own
  - **Michael, Alice** and **Bridget** can all reach each other
  - **Mark** and **Doug** can reach each other

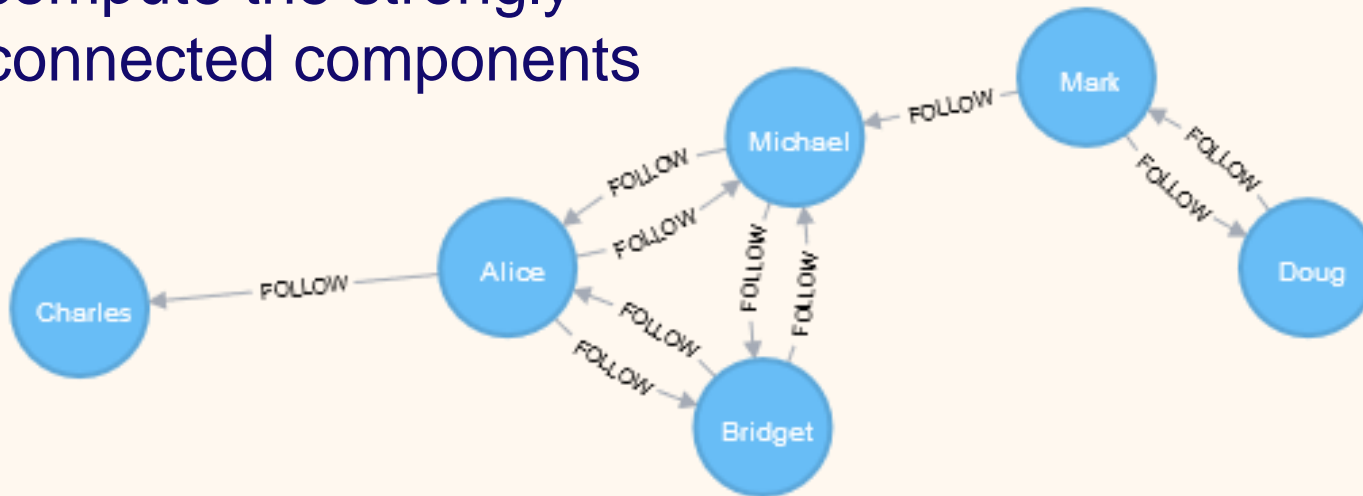


# Community Detection: Strongly Connected Components

## Task 13:

- For **User** nodes and the **:FOLLOW** relationship compute the strongly connected components

```
1 CALL algo.scc.stream("User","FOLLOW")
2 YIELD nodeId, partition
3 MATCH (u:User) WHERE id(u) = nodeId
4 RETURN u.id AS name, partition
```



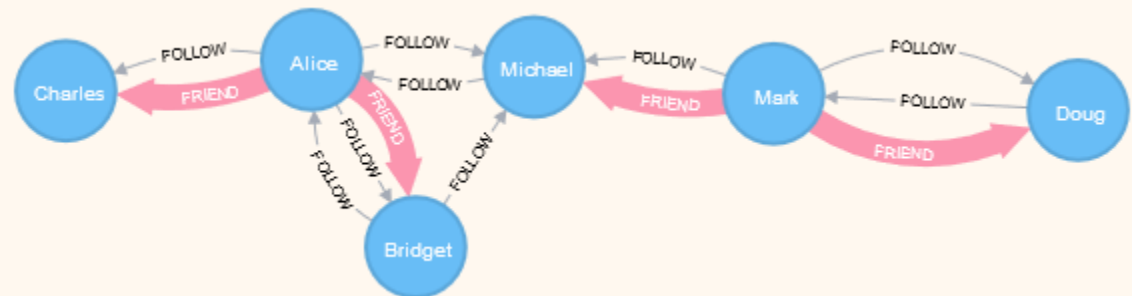
- Optional:** add a **:FOLLOW** edge from **Charles** to **Alice** and observe the effect



# Community Detection: Weakly Connected Components

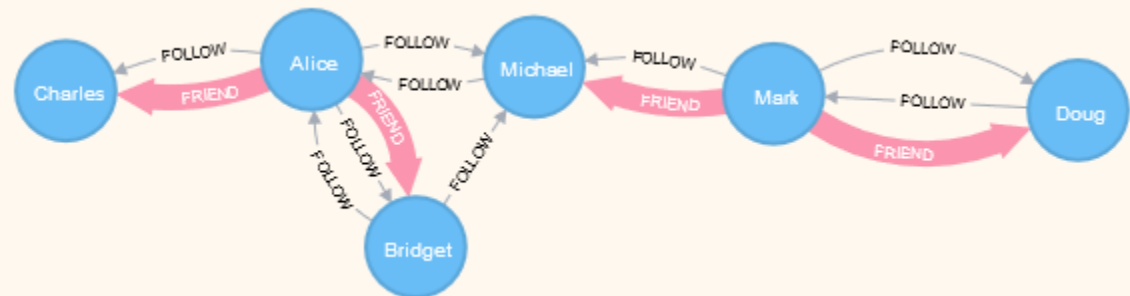
# Community Detection: Weakly Connected Components

- **Weakly Connected Components:**
  - We interpret the graph as undirected
  - Two nodes **A**, **B** are in the same weakly connected component if there is a path from **A** to **B** or a path from **B** to **A**



# Community Detection: Weakly Connected Components

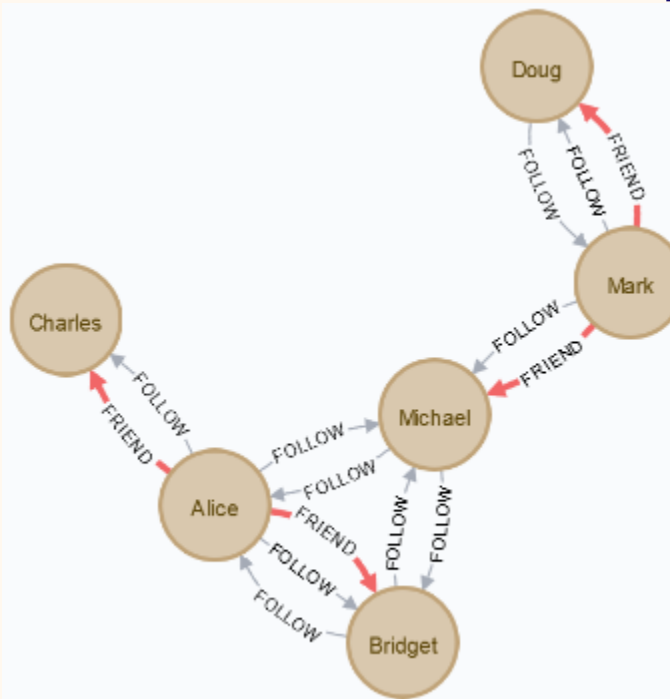
- **Example**
  - New **:FRIEND** relationship
  - **Charles, Alice and Bridget** are weakly connected
  - **Michael, Mark and Doug** are weakly connected



# Community Detection: Weakly Connected Components

## Task 14:

- Add the **:FRIEND** relationship to the graph (see <https://neo4j.com/docs/graph-algorithms/current/> → The Connected Components algorithm)



```
1 MERGE (nAlice:User {id:'Alice'})
2 MERGE (nBridget:User {id:'Bridget'})
3 MERGE (nCharles:User {id:'Charles'})
4 MERGE (nDoug:User {id:'Doug'})
5 MERGE (nMark:User {id:'Mark'})
6 MERGE (nMichael:User {id:'Michael'})
7
8 MERGE (nAlice)-[:FRIEND]->(nBridget)
9 MERGE (nAlice)-[:FRIEND]->(nCharles)
10 MERGE (nMark)-[:FRIEND]->(nDoug)
11 MERGE (nMark)-[:FRIEND]->(nMichael);
```

# Community Detection: Weakly Connected Components

## Task 15:

- Find the weakly connected components according to the **:FRIEND** relationship

```
1 CALL algo.unionFind.stream('User', 'FRIEND', {})  
2 YIELD nodeId, setId  
3 RETURN algo.getNodeById(nodeId).id AS user, setId
```

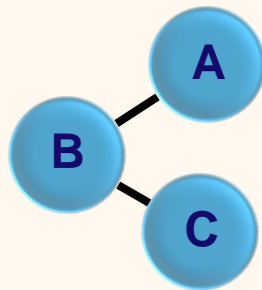
- **Optional:** add a **:FRIEND** edge from **Michael** to **Alice** and observe the effect

# Community Detection: Triangle Counting

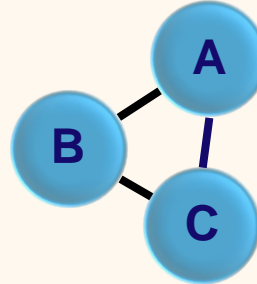
# Community Detection: Triangle Counting

- **Triplet:** 3 connected nodes

- **Open triplet**



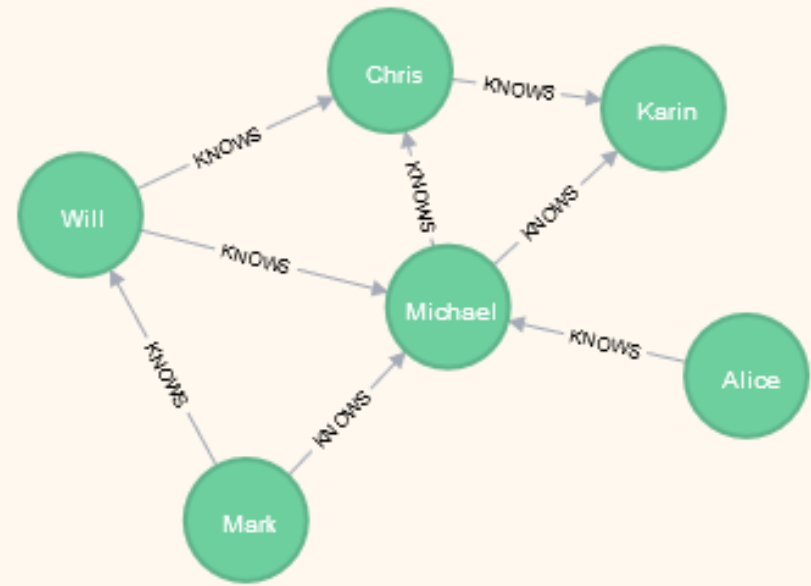
- **Closed triplet = triangle**



- The **older** a community, the **more** triangles are present:
  - if **A** and **B** are friends and **B** and **C** are friends
  - high probability that **A** and **C** become friends, too

# Community Detection: Triangle Counting

- **Example:**
  - **Michael** participates in 3 triangles
  - **Alice** participates in no triangle

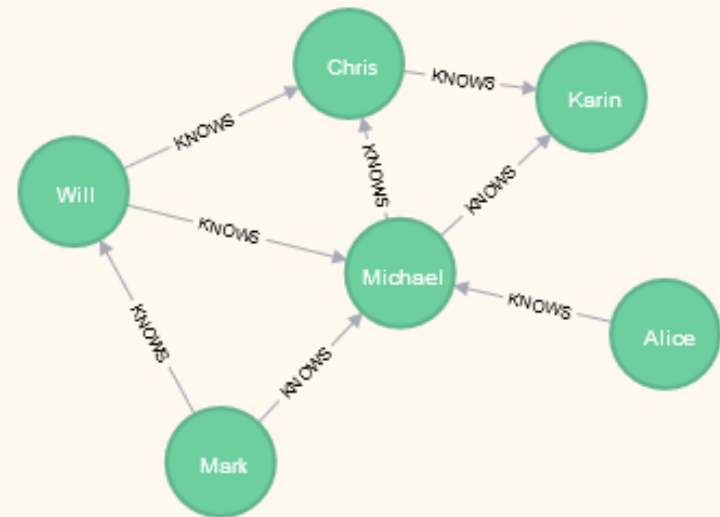




# Community Detection: Triangle Counting

## Task 16:

- Create a new graph with 6 persons and a **:KNOWS** relationship  
(see <https://neo4j.com/docs/graph-algorithms/current/>  
➔ The Triangle Counting / Clustering Coefficient algorithm)

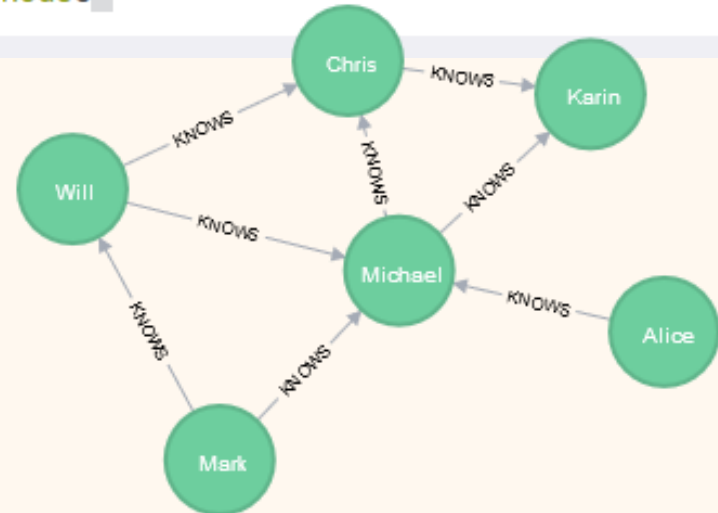


# Community Detection: Triangle Counting

## Task 17:

- Return all triangles in the graph

```
1 CALL algo.triangle.stream('Person','KNOWS')
2 YIELD nodeA,nodeB,nodeC
3
4 RETURN algo.getNodeById(nodeA).id AS nodeA, algo.getNodeById(nodeB).id AS nodeB,
   algo.getNodeById(nodeC).id AS nodeC
```



# Community Detection: Triangle Counting

The local clustering coefficient of a node determines how well connected the node's neighbors are

## Local clustering coefficient:

- For a node  $v$ , count the edges among its neighbors
- Divide by the amount of edges in a complete graph among its neighbors
- **Example:** a person in a social network who is good at connecting his/her friends has a high coefficient

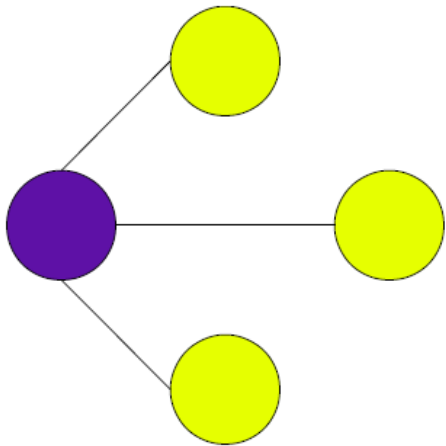
# Community Detection: Triangle Counting

The local clustering coefficient of a node determines how well connected the node's neighbors are

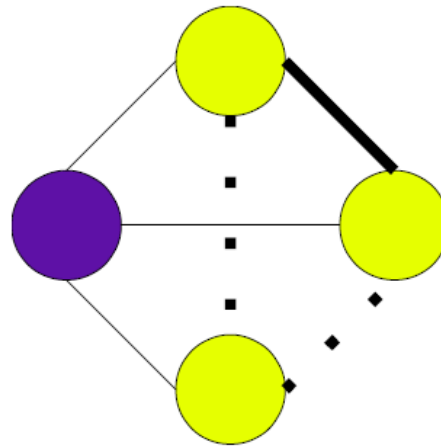
## Local clustering coefficient for undirected graph:

- Recall our quiz from the beginning (edge count in complete undirected graph)
- If a node  $v$  has  $k$  neighbors, the **complete** graph among the neighbors has  $c = (k \cdot (k-1))/2$  edges
- Let the **actual** edge count among the neighbors be  $r$
- $LCC(v) = r/c$

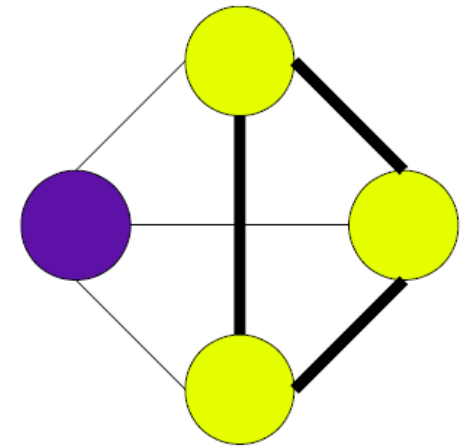
# Community Detection: Triangle Counting



(a) No pairs formed among neighbors:  $C = 0$



(b) One pair formed among neighbors:  $C = 1 / 3$



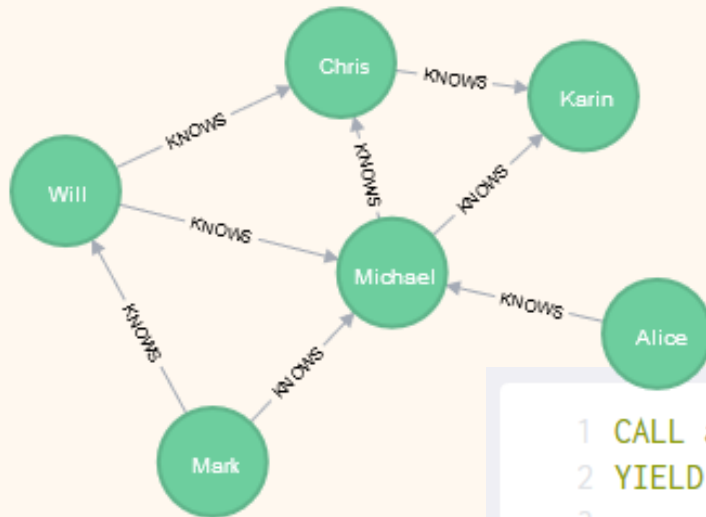
(c) Three pairs formed among neighbors:  $C = 3 / 3$

Picture source: Luis Casillas Santillán/Alonso Castillo Pérez  
<http://www.revistascientificas.udg.mx/index.php/REC/article/viewFile/5091/4754/16111>

# Community Detection: Triangle Counting

## Task 18:

- Compute the amount of triangles and the local clustering coefficient for each node



```
1 CALL algo.triangleCount.stream('Person', 'KNOWS')
2 YIELD nodeId, triangles, coefficient
3
4 RETURN algo.getNodeById(nodeId).id AS name, triangles, coefficient
5 ORDER BY coefficient DESC
```

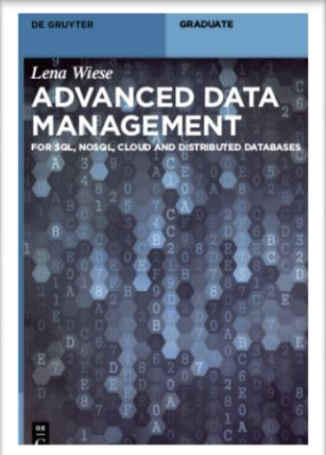
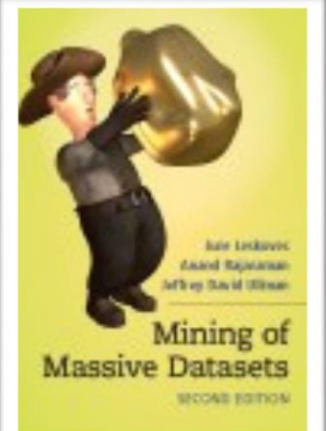
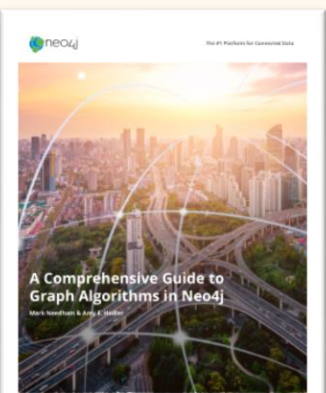
- Which person has the best-connected friends in the graph?

# Agenda

- Short CV of speaker
- Graph Theory Basics
- The Neo4J Database
- Graph Algorithms
  - Centralities
  - Path Finding
  - Community Detection

# References

1. Mark Needham & Amy E. Hodler: A Comprehensive Guide to Graph Algorithms in Neo4j. Neo4j.com, 2018.
2. Anand Rajaraman, Jure Leskovec & Jeffrey D. Ullman: Mining of Massive Datasets. Mmds.org, 2014.
3. Lena Wiese: Advanced Data Management for SQL, NoSQL, Cloud and Distributed Databases. De Gruyter Graduate, 2015.





# Extra exercises

„In this guide we'll learn how to use the Neo4j Graph Algorithms package using a Game of Thrones dataset.”

```
$ :play https://guides.neo4j.com/sandbox/graph-algorithms/
```