

Score and Rank Semi-Monotonicity for Closeness, Betweenness and Harmonic Centrality

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Abstract. In the study of the behavior of centrality measures with respect to network modifications, *score monotonicity* means that adding an arc increases the centrality score of the target of the arc; *rank monotonicity* means that adding an arc improves the importance of the target of the arc relative to the remaining nodes. It is known [8,7] that score and rank monotonicity hold in directed graphs for almost all the classical centrality measures. In undirected graphs one expects that the corresponding properties (where both endpoints of the new edge enjoy the increase in score/rank) hold when adding a new edge. However, recent results [6] have shown that in undirected networks this is not true: for many centrality measures, it is possible to find situations where adding an edge reduces the rank of one of its two endpoints. In this paper we introduce a weaker condition for undirected networks, *semi-monotonicity*, in which just one of the endpoints of a new edge is required to enjoy score or rank monotonicity. We show that this condition is satisfied by closeness and betweenness centrality, and that harmonic centrality satisfies it in an even stronger sense.

1 Introduction and Definitions

In this paper we discuss the behavior of centrality measures in undirected networks after the addition of a new edge. In particular, we are interested in the following question: if a new edge is added to a network, does the importance of *at least one* of its two endpoints increase? This question was left open in [6], where the authors proved that for many centrality measures it is possible to find situations where adding an edge reduces the rank of one of its two endpoints. Note that these results are in jarring contrast with the corresponding properties for directed networks, where it is known [8,7] that score and rank monotonicity hold for almost all centrality measures.

Formally, in this paper we introduce *semi-monotonicity*, a weaker condition than monotonicity for undirected networks in which we require that *at least one* endpoint of the new edge enjoys monotonicity. Score semi-monotonicity, in particular, means that adding a new edge increases the score of at least one of the two endpoints:

Definition 1 (Score semi-monotonicity). *Given an undirected graph G , a centrality c is said to be score semi-monotone on G iff for every pair of non-*

adjacent vertices x and y we have that

$$c_G(x) < c_{G'}(x) \quad \text{or} \quad c_G(y) < c_{G'}(y),$$

where G' is the graph obtained adding the edge $x-y$ to G . We say that c is score semi-monotone on a set of graphs iff it is score semi-monotone on all the graphs from the set.

As we already know from the directed case, a score increase does not imply that the rank relations between the two vertices involved in the new edge and the other vertices in the network remain unchanged. For this reason, rank monotonicity was introduced, where we require that every vertex that used to be dominated is still dominated after the addition of the new edge. Formally, the request for at least one of the two endpoints can be expressed as follows:

Definition 2 (Rank semi-monotonicity). *Given an undirected graph G , a centrality c is said to be rank semi-monotone on G iff for every pair of non-adjacent vertices x and y at least one of the following two statements holds:*

– for all vertices $z \neq x, y$:

$$\begin{aligned} c_G(z) < c_G(x) &\Rightarrow c_{G'}(z) < c_{G'}(x) \quad \text{and} \\ c_G(z) = c_G(x) &\Rightarrow c_{G'}(z) \leq c_{G'}(x), \end{aligned}$$

– for all vertices $z \neq x, y$:

$$\begin{aligned} c_G(z) < c_G(y) &\Rightarrow c_{G'}(z) < c_{G'}(y) \quad \text{and} \\ c_G(z) = c_G(y) &\Rightarrow c_{G'}(z) \leq c_{G'}(y), \end{aligned}$$

where G' is the graph obtained adding the edge $x-y$ to G . We say that c is rank semi-monotone on a set of graphs iff it is rank semi-monotone on all the graphs from the set.

In particular, we say that c is rank semi-monotone at x if the first statement holds, and rank semi-monotone at y if the second statement holds (if both statements hold, c is rank monotone).

Definition 3 (Strict rank semi-monotonicity). *Given an undirected graph G , a centrality c is said to be strictly rank semi-monotone on G iff for every pair of non-adjacent vertices x and y at least one of the following two statements holds:*

- for all vertices $z \neq x, y$: $c_G(z) \leq c_G(x) \Rightarrow c_{G'}(z) < c_{G'}(x)$,
- for all vertices $z \neq x, y$: $c_G(z) \leq c_G(y) \Rightarrow c_{G'}(z) < c_{G'}(y)$,

where G' is the graph obtained adding the edge $x-y$ to G . We say that c is strictly rank semi-monotone on a set of graphs iff it is strictly rank semi-monotone on all the graphs from the set.

Again, we say that it is strictly rank semi-monotone at x if the first statement holds, and strictly rank semi-monotone at y if the second statement holds (if both statements hold, c is strictly rank monotone).

In the rest of the paper, we assume that we are given an undirected connected graph G , and two non-adjacent vertices $x, y \in N_G$; G' will be the graph obtained by adding the edge $x - y$ to G . From now on, d_{uv} will refer to the distance (i.e., the length of a shortest path) between u and v in G (i.e., before the $x - y$ addition), and d'_{uv} will refer to the distance in G' instead. In general, we will use the prime symbol to refer to any property or function of G when translated to G' .

2 Distances and Basins

Geometric centrality measures [8] depend only on distances between vertices. In the next two sections we are going to prove new results about the semi-monotonicity of two geometric centrality measures—*closeness centrality* [2,3] and *harmonic centrality* [4,8]. To understand the semi-monotonic behavior of these centrality measures, we introduce a key notion:

Definition 4 (Basin). *Given an undirected graph G and two non-adjacent vertices x and y we define the basin of x (with respect to y) K_{xy} and the basin of y (with respect to x) K_{yx} as*

$$\begin{aligned} K_{xy} &:= \{u \in N_G \mid d_{ux} \leq d_{uy}\} \\ K_{yx} &:= \{u \in N_G \mid d_{uy} \leq d_{ux}\} \end{aligned}$$

That is, the basin of x contains those vertices that are not farther from x than from y : see Figure 1 for an example. Note that the vertices that are equidistant from x and y are included in both basins, and the score of such vertices cannot change in any geometric centrality when adding the edge $x - y$.

Let us consider the following property:

Definition 5 (Basin dominance). *A centrality c is said to be basin dominant on an undirected graph G iff for every pair of non-adjacent vertices x and y we have that*

$$\begin{aligned} c'(u) - c(u) &\leq c'(x) - c(x) && \text{for every } u \in K_{xy}, u \neq x \\ c'(v) - c(v) &\leq c'(y) - c(y) && \text{for every } v \in K_{yx}, v \neq y. \end{aligned} \tag{1}$$

It is strictly basin dominant iff the same conditions are satisfied, but inequalities (1) hold with the $<$ sign.

Intuitively, basin dominance means that the increase in score of x and y is at least as large as (or larger than, in the strict case) the increase in score of all other nodes in their respective basin.

The following theorems will be used throughout the paper:

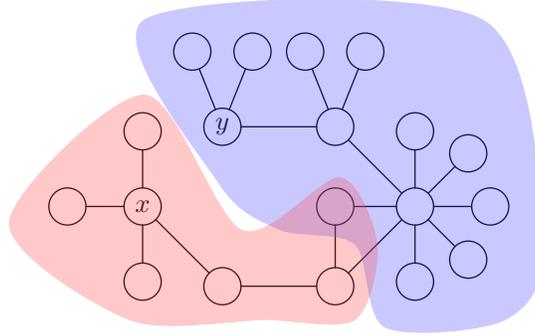


Fig. 1. An undirected graph G , with K_{xy} (the basin of x w.r.t. y) shown in red and K_{yx} (the basin of y w.r.t. x) in blue.

Theorem 1. *If a centrality measure is strictly basin dominant on a graph then it is strictly rank semi-monotone on the same graph.*

Proof. Let c be strictly basin dominant, and let us assume by contradiction that c is not strictly rank semi-monotone. This implies that we should be able to find u, v such that:

$$\begin{cases} c(x) \geq c(v) \\ c(y) \geq c(u) \\ c'(v) \geq c'(x) \\ c'(u) \geq c'(y). \end{cases} \quad (2)$$

As a consequence of (2), $c'(v) - c(v) \geq c'(x) - c(x)$ and $c'(u) - c(u) \geq c'(y) - c(y)$, which by the assumption of strict basin dominance imply $v \notin K_{xy}$ and $u \notin K_{yx}$. Therefore $v \in K_{yx}$ and $u \in K_{xy}$. But then again using strict basin dominance and (2) we have $c'(y) - c(y) > c'(v) - c(v) \geq c'(x) - c(x) > c'(u) - c(u) \geq c'(y) - c(y)$, a contradiction. \square

For the non-strict case, a similar result holds:

Theorem 2. *If a centrality measure is basin dominant on a graph then it is rank semi-monotone on the same graph.*

Proof. Let c be basin dominant, and let us assume by contradiction that c is not rank semi-monotone. This implies that we should be able to find u, v satisfying one of the following four sets of inequalities:

$$\begin{cases} c(x) > c(v) \\ c(y) > c(u) \\ c'(v) \geq c'(x) \\ c'(u) \geq c'(y) \end{cases} \quad \begin{cases} c(x) = c(v) \\ c(y) > c(u) \\ c'(v) > c'(x) \\ c'(u) \geq c'(y) \end{cases} \quad \begin{cases} c(x) > c(v) \\ c(y) = c(u) \\ c'(v) \geq c'(x) \\ c'(u) > c'(y) \end{cases} \quad \begin{cases} c(x) = c(v) \\ c(y) = c(u) \\ c'(v) > c'(x) \\ c'(u) > c'(y) \end{cases} \quad (3)$$

The proof is similar to the strict case. In all sets of inequalities we have $c'(v) - c(v) > c'(x) - c(x)$ and $c'(u) - c(u) > c'(y) - c(y)$, which by the assumption of basin dominance imply $v \notin K_{xy}$ and $u \notin K_{yx}$. Therefore $v \in K_{yx}$ and $u \in K_{xy}$. Using basin dominance and (3), in all cases we have $c'(y) - c(y) \geq c'(v) - c(v) > c'(x) - c(x) \geq c'(u) - c(u) > c'(y) - c(y)$, a contradiction. \square

3 Closeness Centrality

Closeness centrality [2,3] was shown to be score monotone but not rank monotone on connected undirected networks [6]. In the latter paper, it was left open the problem of whether closeness was (in our terminology) rank semi-monotone or not. In the rest of this section, we will solve this open problem by showing that closeness is in fact rank semi-monotone, but not in strict form.

Recall that, given an undirected connected graph $G = (N_G, A_G)$, the *peripherality* of a vertex $u \in N_G$ is the sum of the distances from u to all the other vertices of G :

$$p(u) = \sum_{v \in N_G} d_{uv}.$$

The *closeness centrality* of u is just the reciprocal of its peripherality:

$$c(u) = \frac{1}{p(u)}.$$

Consistently with the previous notation, we will use $p(u)$ for the peripherality of u in G , and $p'(u)$ for the peripherality of u in G' .

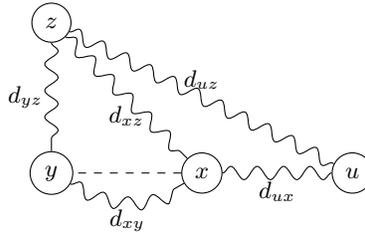


Fig. 2. Path labels represent the distance between the two endpoints. The dashed edge is the $x - y$ edge that we add to G , obtaining G' .

We have that:

Lemma 1. *Closeness centrality is basin dominant on connected undirected graphs.*

Proof. We first show that for every vertex $u \in K_{xy}$ and for every $z \neq u, x$:

$$d_{uz} - d'_{uz} \leq d_{xz} - d'_{xz}. \tag{4}$$

Then, the result follows by adding both sides for all z (as for $z = u$ or $z = x$ the inequality trivializes) and using the fact that closeness is reciprocal to peripherality. To prove (4) we consider two cases (see Figure 2):

- if $d'_{uz} = d_{uz}$ then the inequality trivially holds (because $d_{xz} \geq d'_{xz}$ always);
- if $d'_{uz} < d_{uz}$ then $d'_{uz} = d_{ux} + 1 + d_{yz} < d_{uz}$ and $d'_{xz} = d_{yz} + 1$. Using $d_{uz} \leq d_{ux} + d_{xz}$ from the triangle inequality we have

$$\begin{aligned} d_{uz} - d'_{uz} &= d_{uz} - (d_{ux} + 1 + d_{yz}) \leq \\ &(d_{ux} + d_{xz}) - (d_{ux} + 1 + d_{yz}) = d_{xz} - (d_{yz} + 1) = d_{xz} - d'_{xz}. \end{aligned} \quad \square$$

This lemma is the undirected version of Lemma 2 in [7] (provided that in the latter you sum the inequality in the statement over all nodes w): it is interesting to observe that in the directed case the inequality holds *for all* u 's, while here it holds only within the basin.

Applying Lemma 1 and using Theorem 2, we obtain that:

Theorem 3. *Closeness centrality is rank semi-monotone on connected undirected graphs.*

Interestingly, and regardless of their initial score, we can always tell which of the two endpoints of the edge $x - y$ will have smaller peripherality (i.e., higher centrality) in G' . In fact:

Lemma 2. *The following property holds:*

$$p'(x) - p'(y) = |K_{yx}| - |K_{xy}|.$$

Proof. We can write the peripherality of x and y in G' as

$$\begin{aligned} p'(x) &= \sum_{u \in K_{xy}} d_{ux} + \sum_{u \in K_{yx}} (1 + d_{uy}) - \sum_{u \in K_{xy} \cap K_{yx}} (1 + d_{uy}) \\ p'(y) &= \sum_{u \in K_{xy}} (1 + d_{ux}) + \sum_{u \in K_{yx}} d_{uy} - \sum_{u \in K_{xy} \cap K_{yx}} (1 + d_{ux}). \end{aligned}$$

Note that for each $u \in K_{xy} \cap K_{yx}$ we have $d_{ux} = d_{uy}$. Computing the difference between the two expressions gives the result. \square

It is not hard to build a graph G where x has a smaller basin than y but a greater score: Lemma 2 tells us that y becomes more central than x in G' , due to having a greater basin.

We conclude this section by showing that:

Theorem 4. *Closeness centrality is not strictly rank semi-monotone on (an infinite family of) connected undirected graphs.*

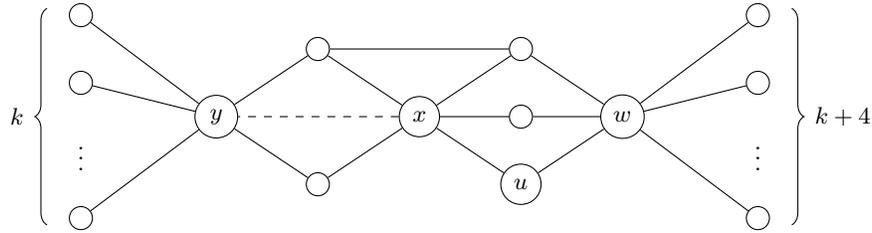


Fig. 3. A counterexample to strict rank semi-monotonicity for closeness centrality. For all $k \geq 10$, u and x have the same score before and after the addition of the edge $x - y$. Moreover, u has the same score of y (or smaller) before the addition, but a higher score after the addition, breaking strict rank semi-monotonicity.

Proof. Consider the graphs in Figure 3, where $u \in K_{xy}$. This is an infinite family of graphs with a parameter k which controls the sizes of the two stars around vertices w and y . Computing the peripheralities of u , x and y before and after the addition of $x - y$, we obtain

$$\begin{aligned}
 p(u) &= 2 \cdot (k + 4) + 4 \cdot k + 13 & p'(u) &= 2 \cdot (k + 4) + 3 \cdot k + 12 \\
 p(x) &= 3 \cdot (k + 4) + 3 \cdot k + 9 & p'(x) &= 3 \cdot (k + 4) + 2 \cdot k + 8 \\
 p(y) &= 4 \cdot (k + 4) + k + 15 & p'(y) &= 4 \cdot (k + 4) + k + 12.
 \end{aligned}$$

For all $k \geq 10$, we have that

$$p(x) = p(u), \quad p'(x) = p'(u), \quad p(y) \leq p(u), \quad p'(y) > p'(u),$$

showing that closeness is not semi-monotone at y (because y used to be at least as central as u , but it is less central after the addition of the edge) and not strictly rank semi-monotone at x (it is always as central as x , before and after adding the edge). \square

4 Harmonic Centrality

Harmonic centrality [4] solves the issue of unreachable vertices in closeness centrality. In particular, if we assume $\infty^{-1} = 0$, we can define it as

$$h(u) = \sum_{v \in N_G \setminus \{u\}} \frac{1}{d_{uv}},$$

so that unreachable vertices have a null impact on the summation and, thus, on the final centrality score of the node. Being a geometric measure, it is trivially score monotone but not rank monotone, as shown in [6], where the same counterexample disproving rank monotonicity for closeness centrality also shows that harmonic centrality fails at satisfying this axiom.

Lemma 3. *Harmonic centrality is strictly basin dominant on connected undirected graphs.*

Proof. We first show that for every vertex $u \in K_{xy}$ and for every $z \neq u, x$:

$$\frac{1}{d'_{uz}} - \frac{1}{d_{uz}} \leq \frac{1}{d'_{xz}} - \frac{1}{d_{xz}}. \quad (5)$$

Then, the result follows by adding both sides for all z . Note that the unique term in $h(u)$ and $h'(u)$ (i.e., when $z = x$) is equal to the unique term in $h(x)$ and $h'(x)$ (i.e., when $z = u$), and they are both equal to 0. Also, we remark that (5) holds with a strict inequality at least in one case, i.e., when $z = y$, since $d'_{xy} < d_{xy}$ always.

We consider two cases:

- if $d'_{uz} = d_{uz}$ then the inequality trivially holds (because $d'_{xz} \leq d_{xz}$ always);
- if $d'_{uz} < d_{uz}$ then $d'_{uz} = d_{ux} + 1 + d_{yz} < d_{uz}$ and $d'_{xz} = d_{yz} + 1$. Using $d_{uz} \leq d_{ux} + d_{xz}$ from the triangle inequality we have

$$\begin{aligned} \frac{1}{d'_{uz}} - \frac{1}{d_{uz}} - \left(\frac{1}{d'_{xz}} - \frac{1}{d_{xz}} \right) &= \frac{1}{d_{ux} + 1 + d_{yz}} - \frac{1}{d_{uz}} - \left(\frac{1}{d_{yz} + 1} - \frac{1}{d_{xz}} \right) \\ &= \frac{(d_{yz} + 1) - (d_{ux} + 1 + d_{yz})}{(d_{ux} + 1 + d_{yz})(d_{yz} + 1)} + \frac{d_{uz} - d_{xz}}{d_{uz}d_{xz}} \\ &= -\frac{d_{ux}}{(d_{ux} + 1 + d_{yz})(d_{yz} + 1)} + \frac{d_{uz} - d_{xz}}{d_{uz}d_{xz}} \\ &< -\frac{d_{ux}}{d_{uz}d_{xz}} + \frac{(d_{ux} + d_{xz}) - d_{xz}}{d_{uz}d_{xz}} = 0, \end{aligned}$$

which proves the inequality. \square

Using Lemma 3 and Theorem 1, we obtain that:

Theorem 5. *Harmonic centrality is strictly rank semi-monotone on connected undirected graphs.*

The stronger result we can give for harmonic centrality should be compared to the fact that on strongly connected graphs harmonic centrality is strictly rank monotone, whereas closeness centrality is just rank monotone [6].

5 Betweenness Centrality

Betweenness centrality [1,9] focuses not only on the length of shortest paths, but also on how many of them involve a given node, trying to estimate the amount of flow passing through nodes in a network. Note that betweenness is not a geometric measure.

Formally, if we call σ_{vw} the number of shortest paths between two vertices v and w and $\sigma_{vw}(u)$ the number of such paths passing through u , then we can define the betweenness centrality of a vertex $u \in N_G$ as

$$b(u) = \sum_{\substack{i,j \neq u, \\ \sigma_{ij} > 0}} \frac{\sigma_{ij}(u)}{\sigma_{ij}}.$$

In the following, we let $N_G(u)$ denote the set of neighbors of u in G , and $G[u]$ the subgraph of G induced by $N_G(u)$ (sometimes called the *ego network* of u). As in the previous section, we denote with σ and σ' the number of shortest paths before and after the addition of an edge $x-y$, and with b and b' the betweenness centrality before and after the addition of the edge.

We know from [6] that this centrality measure is neither rank nor score monotone. Nonetheless, we can show that the betweenness centrality of two vertices can never decrease after we link them with a new edge. In fact:

Lemma 4. *The following properties hold:*

$$\frac{\sigma'_{ij}(x)}{\sigma'_{ij}} - \frac{\sigma_{ij}(x)}{\sigma_{ij}} \geq 0 \quad \text{for all } i, j \neq x.$$

As a consequence, $b'(x) \geq b(x)$.

Proof. For all $i, j \in N_G$ such that $i, j \neq x$, let us call p_x ($p_{\bar{x}}$, respectively) the number of shortest paths between i and j passing (not passing, resp.) through x . We have to show that, for each such pair $i, j \neq x$, the following holds:

$$\frac{\sigma'_{ij}(x)}{\sigma'_{ij}} - \frac{\sigma_{ij}(x)}{\sigma_{ij}} = \frac{p'_x}{p'_x + p'_{\bar{x}}} - \frac{p_x}{p_x + p_{\bar{x}}} \geq 0. \quad (6)$$

Summing over all $i, j \neq x$ proves the second part of the statement.

To show that (6) is indeed true we consider two cases:

- $d'_{ij} < d_{ij}$, meaning that in G' all the shortest paths $i \sim j$ pass through the edge $x-y$ (in particular through x). Thus, we obtain:

$$1 - \frac{p_x}{p_x + p_{\bar{x}}} \geq 0,$$

which is clearly true, since the second term of the left side of the inequality is a value between 0 and 1.

- $d'_{ij} = d_{ij}$, which implies that $p'_{\bar{x}} = p_{\bar{x}}$ and $p'_x \geq p_x$. We express p'_x as $p_x + \alpha$ with $\alpha \geq 0$, obtaining:

$$\begin{aligned} \frac{p_x + \alpha}{p_x + \alpha + p_{\bar{x}}} - \frac{p_x}{p_x + p_{\bar{x}}} &= \frac{p_x^2 + p_x p_{\bar{x}} + \alpha p_x + \alpha p_{\bar{x}} - (p_x^2 + \alpha p_x + p_x p_{\bar{x}})}{(p_x + \alpha + p_{\bar{x}})(p_x + p_{\bar{x}})} = \\ &= \frac{\alpha p_{\bar{x}}}{(p_x + \alpha + p_{\bar{x}})(p_x + p_{\bar{x}})} \geq 0, \end{aligned}$$

which is again true, concluding the proof. \square

Incidentally, we can always tell if the betweenness centrality of a vertex is zero without actually computing it. In fact,

Lemma 5. *Let G be a connected undirected graph and $u \in N_G$. Then:*

$$b(u) = 0 \iff G[u] \text{ is a clique.}$$

Proof. If $b(u) = 0$ no shortest paths are passing through u : but then any two neighbors of u must be adjacent (or otherwise they would have distance 2, and the path through u has length 2). Conversely, suppose that $G[u]$ is a clique and let $i, j \neq u$. A path from i to j cannot involve u , otherwise it would touch two neighbors of u , say i', j' , and we might shorten it by skipping u and taking the $i' - j'$ edge instead of $i' - u - j'$. \square

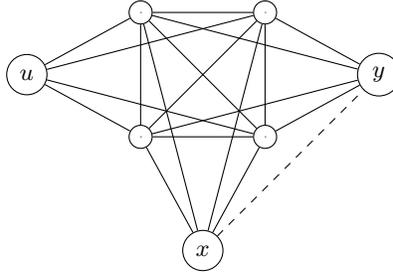


Fig. 4. Simple counterexample for score semi-monotonicity and strict rank semi-monotonicity for betweenness centrality. The dashed edge is the $x - y$ edge that we add to G , obtaining G' . The betweenness score of vertices x , y and u is 0 both in G and G' .

As a consequence, and differently from geometric measures, we can show that

Theorem 6. *Betweenness centrality is not score semi-monotone on (an infinite family of) connected undirected graphs.*

Proof. Consider a graph G such that $G[x]$ and $G[y]$ are cliques and moreover, $N_G(x) = N_G(y)$, that is, x and y have the same neighborhood in G (see Figure 4 for an example). Then, by Lemma 5 we know that $b(x) = b(y) = 0$. It is easy to observe that $G'[x]$ and $G'[y]$ are still cliques, hence we can use the same lemma and say that $b'(x) = b'(y) = 0$, meaning that the addition of the $x - y$ leaves the score of both the endpoints unchanged. \square

Moreover, we can say that

Theorem 7. *Betweenness centrality is not strictly rank semi-monotone on (an infinite family of) connected undirected graphs.*

Proof. Consider a graph G with three vertices x, y and u adjacent to a clique, but with x, y , and u not adjacent to each other, as in Figure 4. Then, by Lemma 5 we know that $b(x) = b(y) = b(u) = 0$; if we add the edge $x - y$ to G , obtaining G' , we also have $b'(x) = b'(y) = b'(u) = 0$ by the same lemma. \square

We are now going to show that, somehow unexpectedly, betweenness centrality is rank semi-monotone on connected undirected graphs. In fact, we show that it enjoys the same dominance property of closeness centrality, in spite of being a non-geometric measure:

Lemma 6. *Betweenness centrality is basin dominant on connected undirected graphs.*

Proof. Let us call $\Delta_z = b'(z) - b(z)$ the score difference for a vertex $z \in N_G$, and for every pair of nodes i, j (with $i \neq j$) let also

$$\Delta_z(i, j) = \frac{\sigma'_{ij}(z)}{\sigma'_{ij}} - \frac{\sigma_{ij}(z)}{\sigma_{ij}}.$$

Obviously

$$\Delta_z = \sum_{i, j \neq z} \Delta_z(i, j).$$

We want to show that $\Delta_u \leq \Delta_x$ for every $u \in K_{xy}$. We do this by analyzing the summands $\Delta_u(i, j)$ and $\Delta_x(i, j)$ separately. For reasons of space, the remaining part of the proof appears only in the full version available on arXiv [5]. \square

Hence, applying Theorem 2 with Lemma 6 we have that:

Theorem 8. *Betweenness centrality is rank semi-monotone on connected undirected graphs.*

6 Conclusions and Future Work

Table 1 summarizes the results of this paper along with those of [7,6,8]. For all the negative results, we have an infinite family of counterexamples (for instance, there are infinitely many graphs on which closeness is shown to be not strictly semi-monotone).

The notion of basin dominance turned out to be the key idea in all proofs of semi-monotonicity. It would be interesting to investigate whether basin dominance applies to other geometric measures, or even other centrality measures based on shortest paths, as in that case one gets immediately rank semi-monotonicity.

Proving or disproving score and (strict) rank semi-monotonicity for other measures (in particular, for the spectral ones) remains an open problem.

	undirected		directed [7,8]	
	score	rank	score	rank
Closeness	monotone [6]	semi-monotone	monotone	monotone
Harmonic centrality	monotone [6]	strictly semi-mon.	monotone	strictly monotone
Betweenness	not semi-monotone	semi-monotone	not monotone	not monotone

Table 1. Summary of the results about monotonicity obtained in this paper (in bold-face) and in [7,6,8]. All results are about (strongly) connected graphs.

Acknowledgments

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